



<https://hao-ai-lab.github.io/dsc204a-w24/>

DSC 204A: Scalable Data Systems

Winter 2024

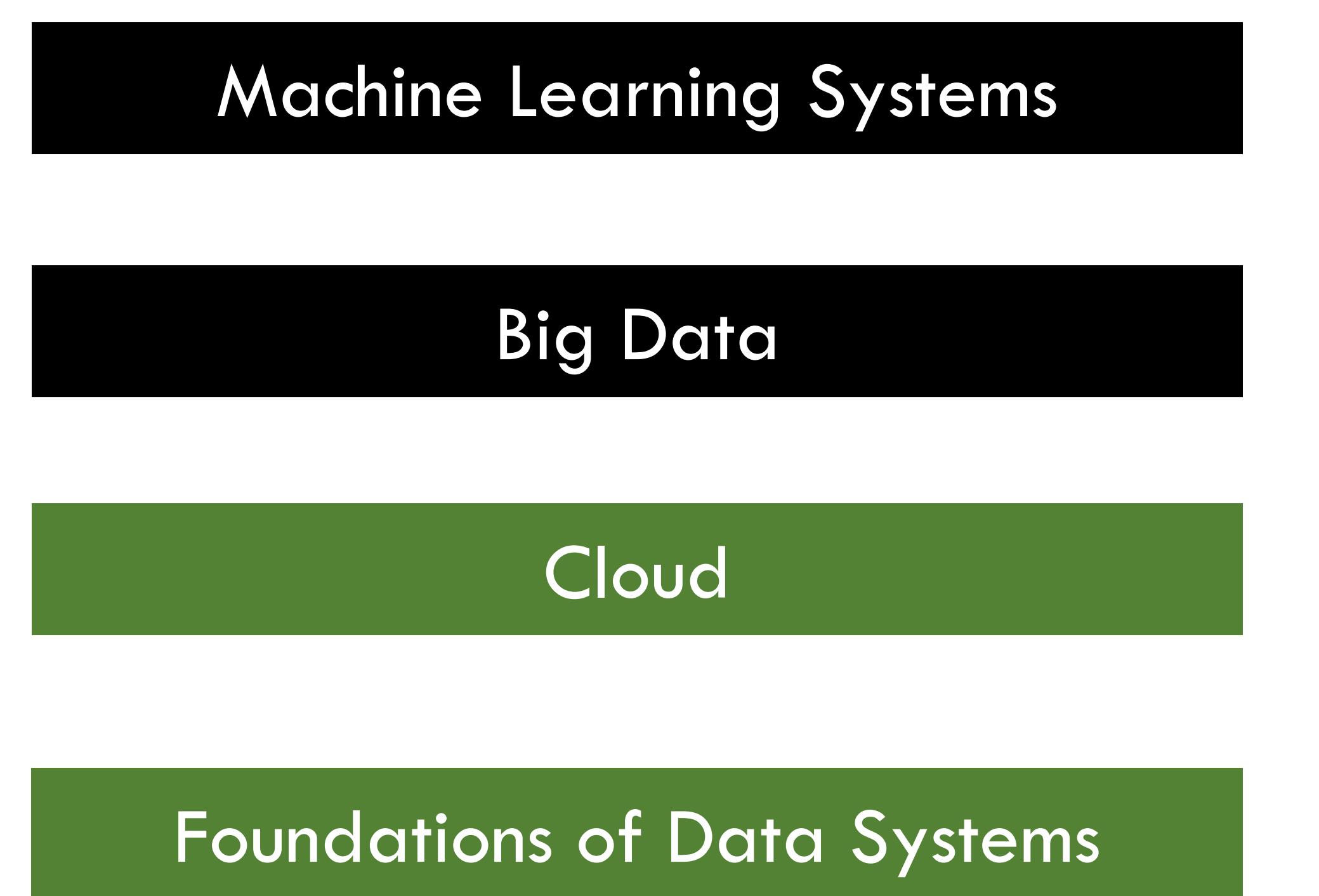
Machine Learning Systems

Big Data

Cloud

Foundations of Data Systems

Where We Are



Recap: Networking

- Q1: True _ False **X** Protocols specify the implementation
- Q2: True _ False **X** Congestion control takes care of the sender not overflowing the receiver
- Q3: True **X** False _ A random access protocol is efficient at low utilization
- Q4: True _ False **X** At the data link layer, hosts are identified by IP addresses
- Q5: True **X** False _ The physical layer is concerned with sending and receiving bits

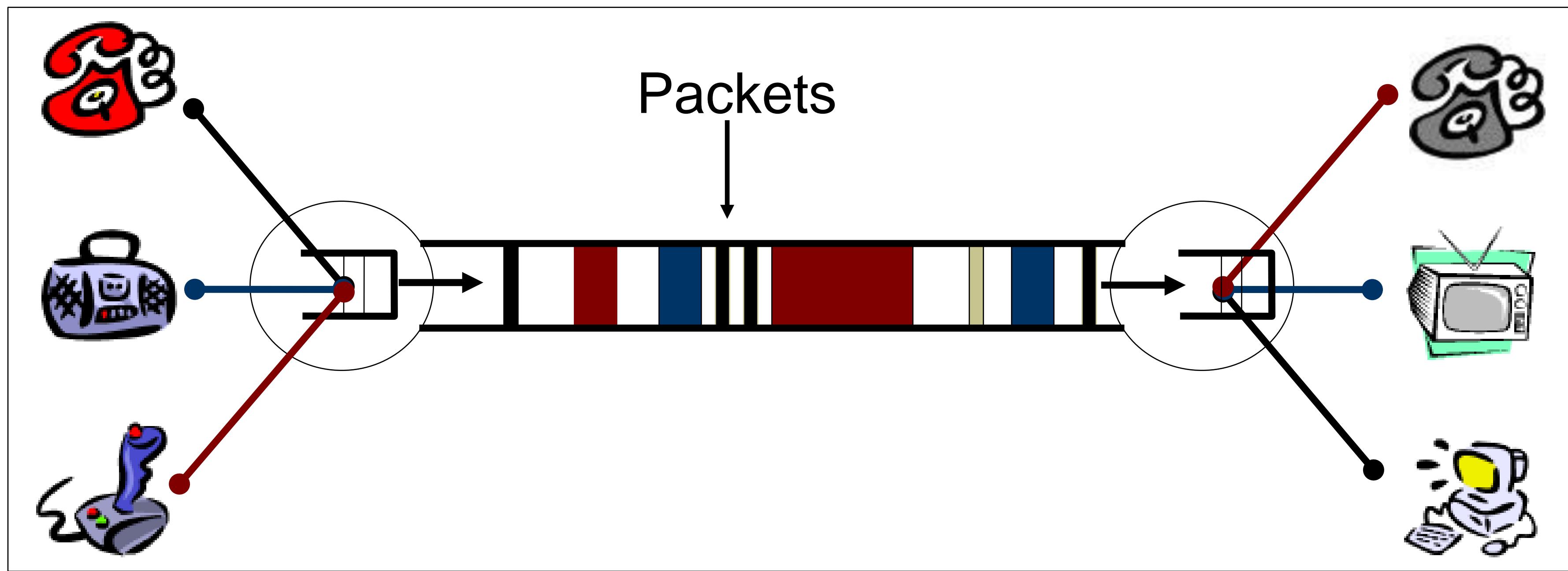
Recap: Networking

- Q1: True _ False X Layering improves application performance
- Q2: True X False _ Routers forward a packet based on its destination address
- Q3: True _ False X “Best Effort” packet delivery ensures that packets are delivered in order
- Q4: True _ False X Port numbers belong to network layer

Today's topic

- Network Basics
- Layering and protocols
- **Collective communication**

Communication

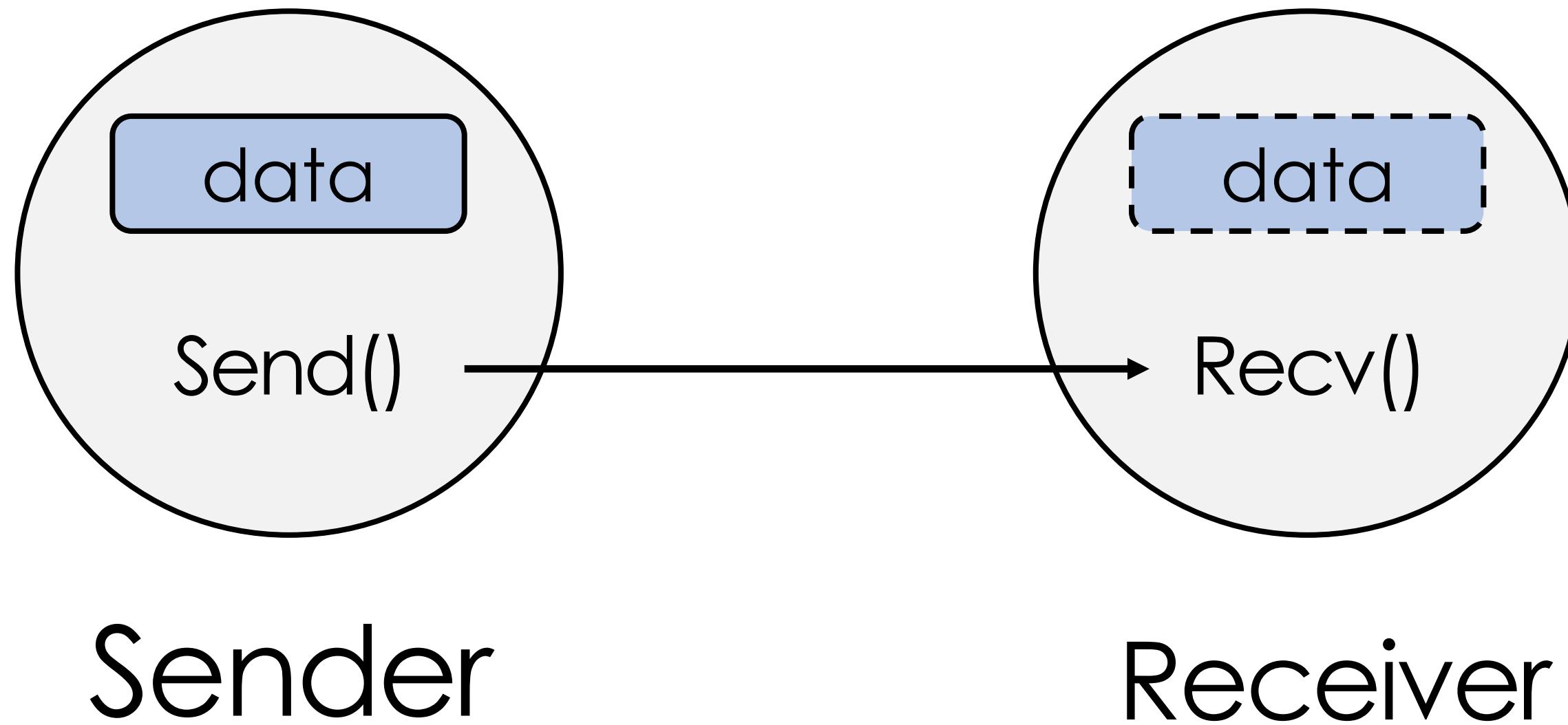


Communication: Point-to-point communication

1. Establish TCP connection
2. Application sends data
3. Data goes down 5 layers, through the network, and arrives in receiver



Program P2P communication: Very Simple in Ray



```
def send(array: np.array):  
    # Running on sender process  
    ref = ray.put(array)  
    return ref  
  
def receive(ref):  
    # Running on receiver process  
    array = ray.get(ref)  
    print(array)
```

Case study: Gradient update in DL

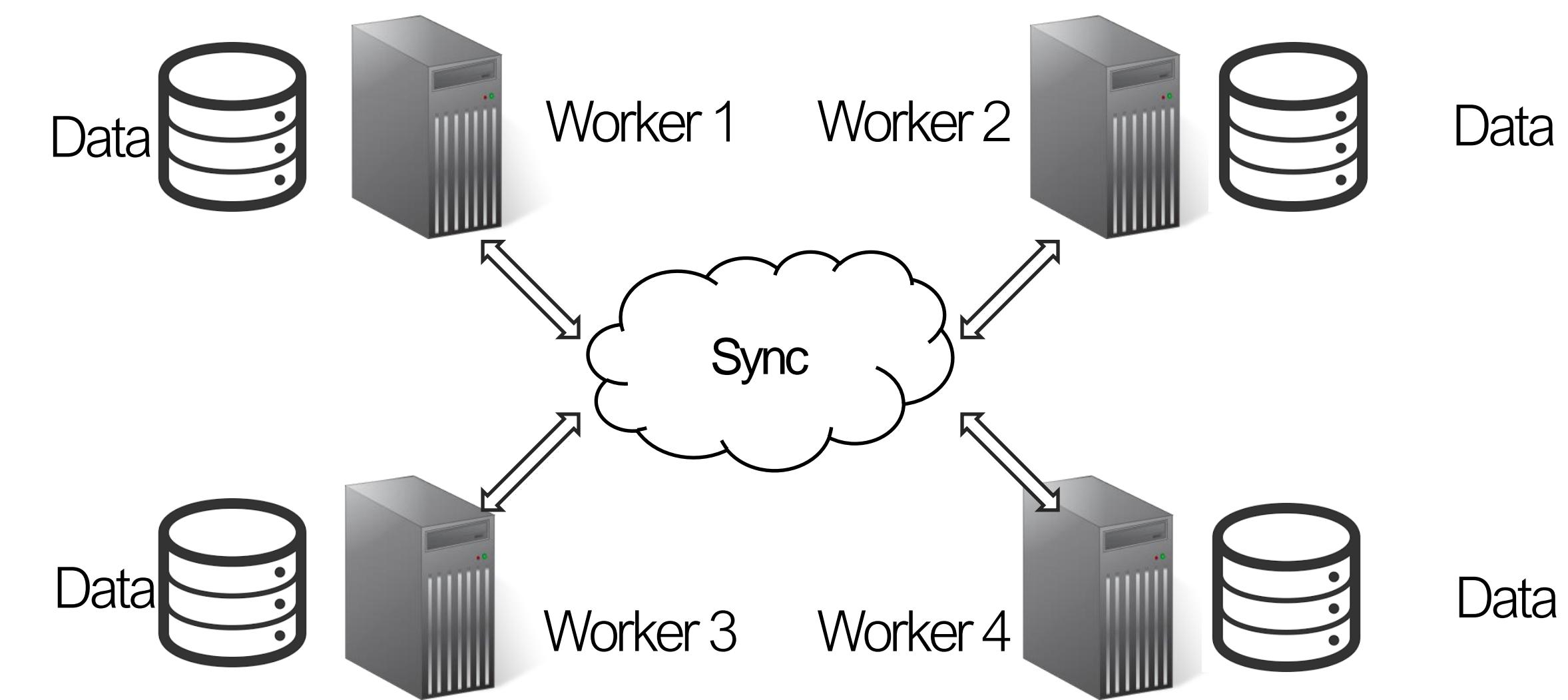
Gradient / backward computation

$$\theta^{(t)} = \theta^{(t-1)} + \boxed{\varepsilon \cdot \nabla_{\mathcal{L}}(\theta^{(t-1)}, D^{(t)})}$$

The diagram shows the gradient update equation $\theta^{(t)} = \theta^{(t-1)} + \varepsilon \cdot \nabla_{\mathcal{L}}(\theta^{(t-1)}, D^{(t)})$. A rectangular box encloses the term $\varepsilon \cdot \nabla_{\mathcal{L}}(\theta^{(t-1)}, D^{(t)})$. An arrow points down to the right side of the box, and two arrows point up from below to the left and right sides of the box, labeled "objective" and "data" respectively.

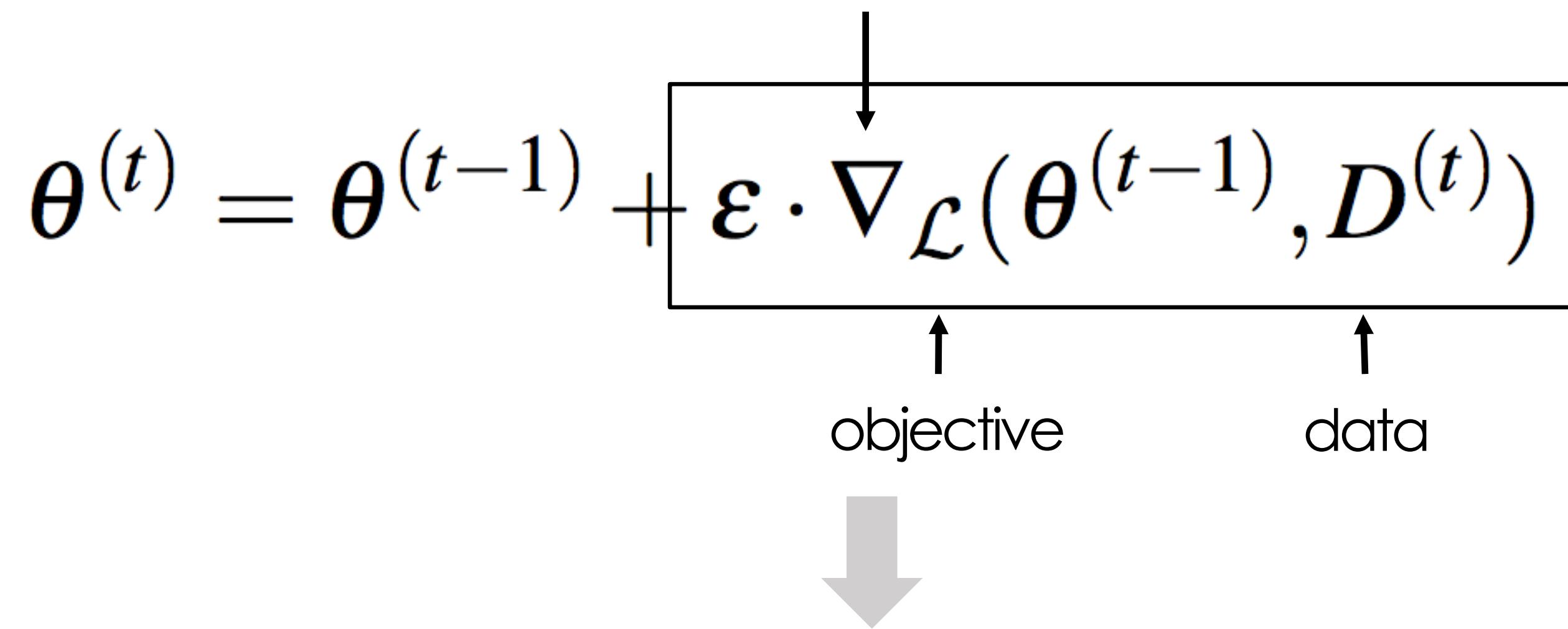
What If Data is super Big?

What if Data is Super Big?



Case study: Gradient update

Gradient / backward computation

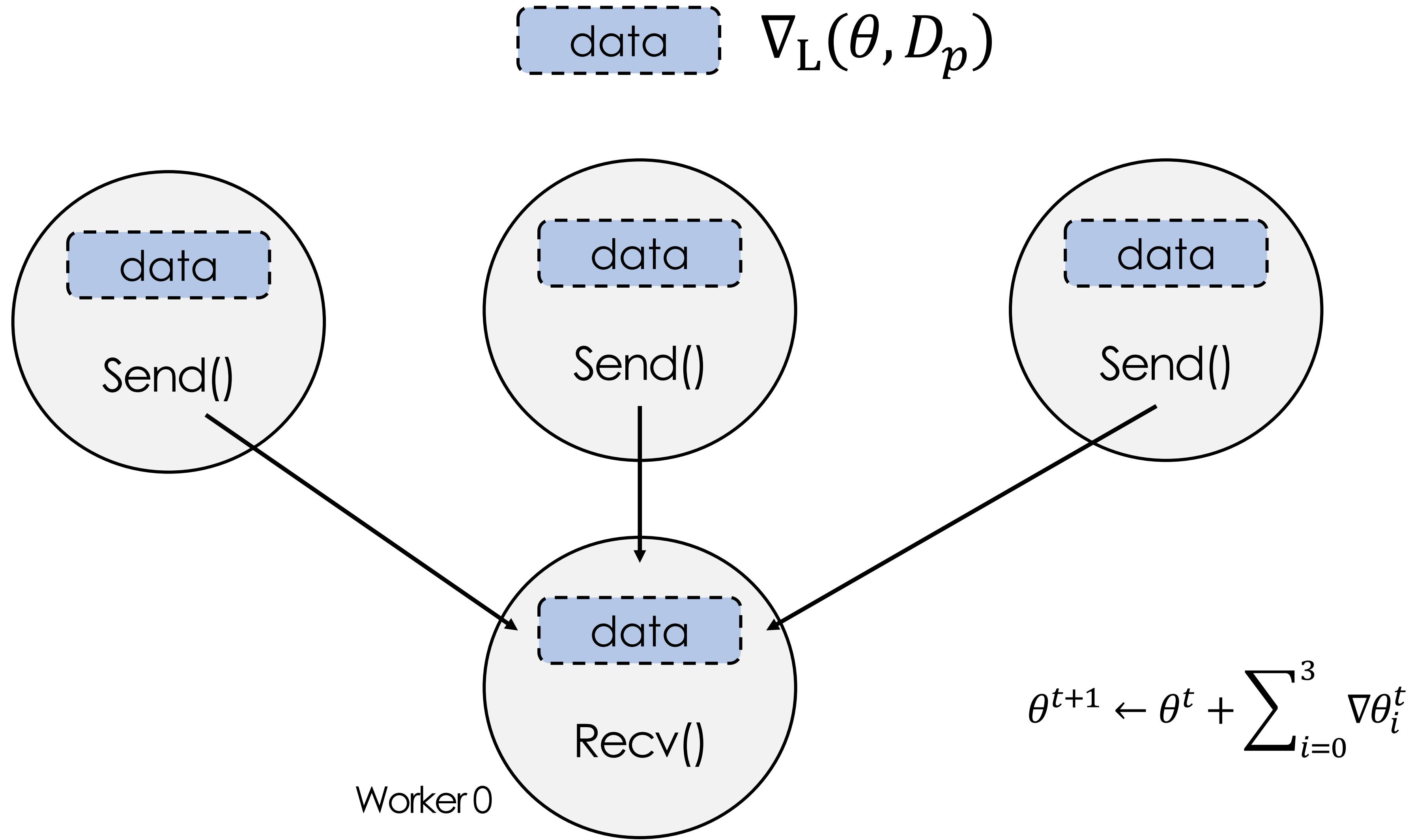


$$\theta^{(t+1)} = \theta^{(t)} + \epsilon \sum_{p=1}^P \nabla_{\mathcal{L}}(\theta^{(t)}, D_p^{(t)})$$

A large downward arrow originates from the right side of the equation, pointing towards the bottom. A diagonal arrow points from the text "How to perform this sum?" to the summation symbol \sum .

How to perform this sum?

Collective Primitive: Reduce



Program This?

```
@ray.remote(num_gpus=1)
class GPUWorker:
    def __init__(self):
        self.gradients = cupy.ones((10,), dtype=cupy.float32)

    def put_gradients(self):
        return ray.put(self.gradients)

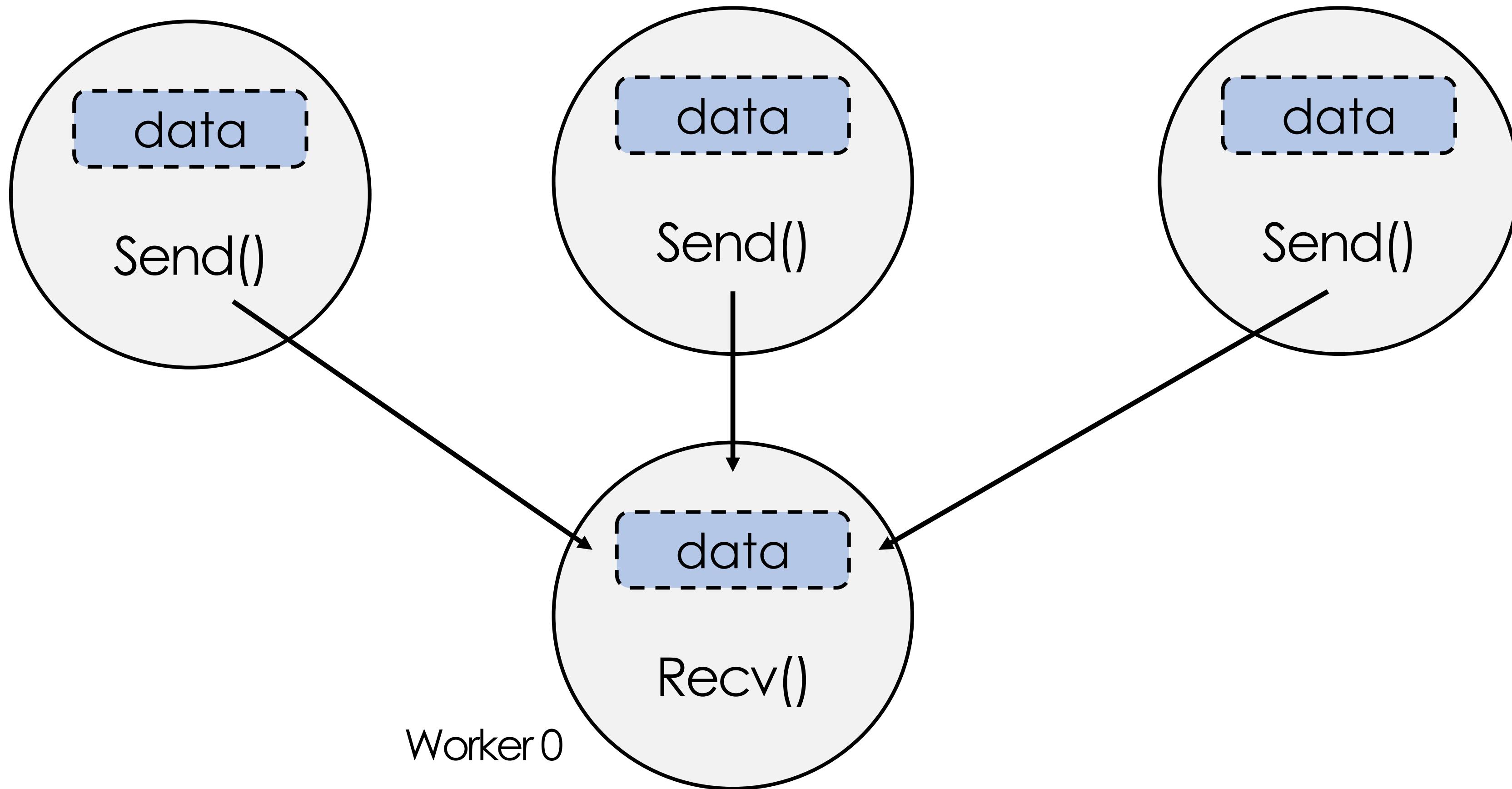
    def reduce_gradients(self, grad_id_refs):
        grad_ids = ray.get(grad_id_refs)
        reduced_result = cupy.ones((10,), dtype=float32)
        for grad_id in grad_ids:
            array = ray.get(grad_id)
            reduced_result += array
        result_id = ray.put(reduced_result)
        return result_id
```

```
# Allreduce the gradients using Ray APIs
# Let all workers to put their gradients into the Ray object store.
gradient_ids = [worker.put_gradients.remote() for worker in workers]
ray.wait(gradient_ids, num_returns=len(gradient_ids, timeout=None))

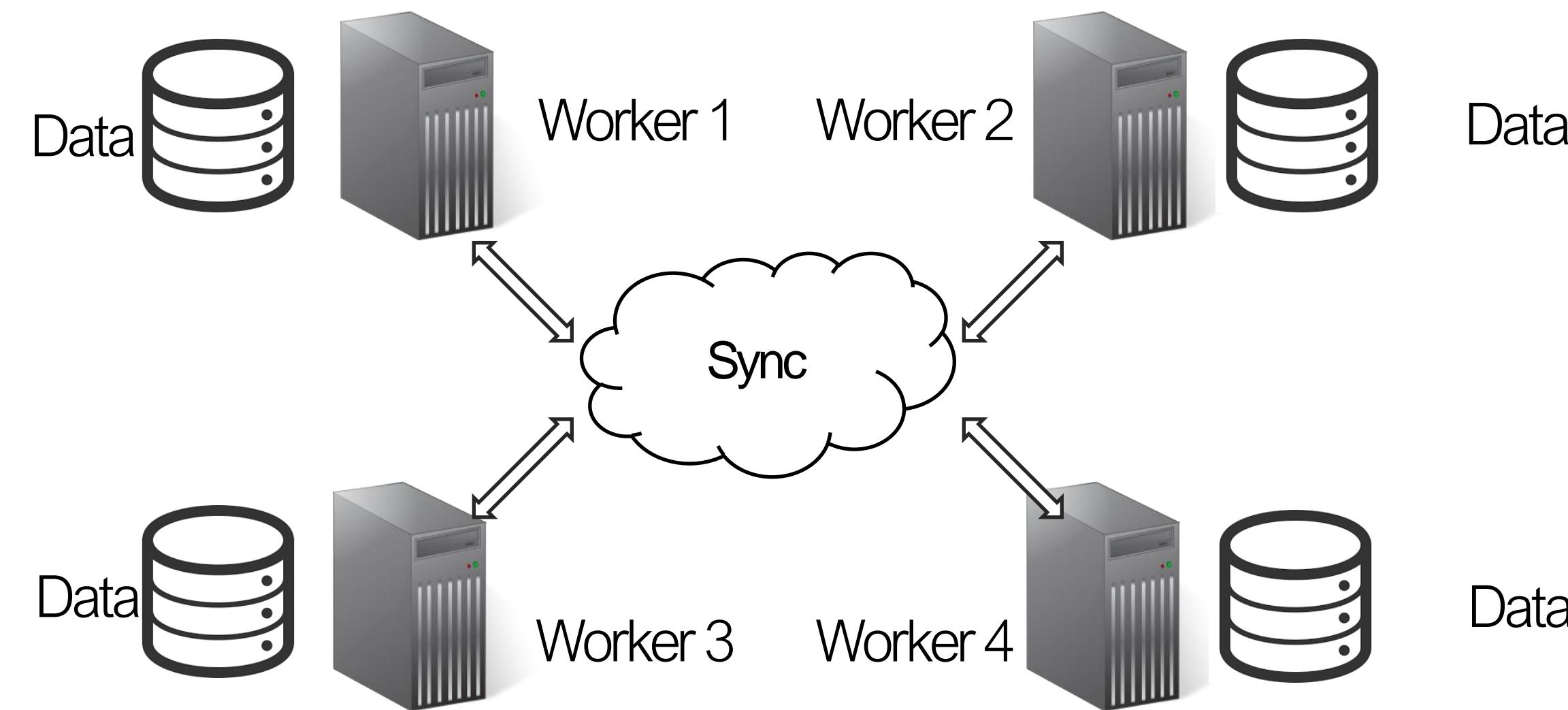
# Let worker 0 reduce the gradients
reduced_id_ref = workers[0].reduce_gradients.remote(gradient_ids)
```

Analyze Performance

- Message over networks: $3 \times N$.
- Can we do better?



Not Yet Finished: Synchronization

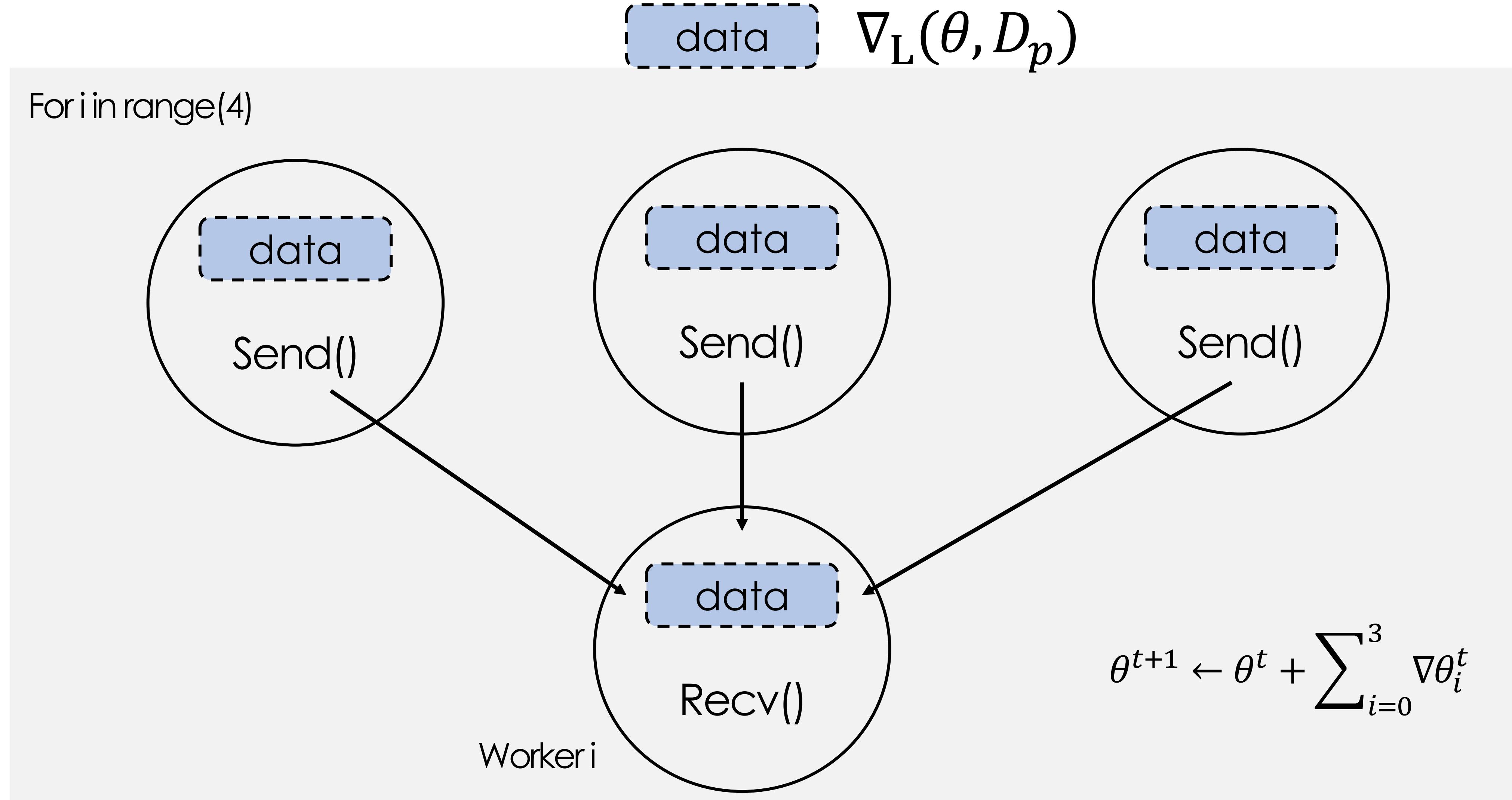


For each worker

$$\theta^{(t+1)} = \theta^{(t)} + \epsilon \sum_{p=1}^P \nabla_{\mathcal{L}}(\theta^{(t)}, D_p^{(t)})$$

How to perform this sum?

Problem: We need All-Reduce



Program This?

```
@ray.remote(num_gpus=1)
class GPUworker:
    def __init__(self):
        self.gradients = cupy.ones((10,), dtype=cupy.float32)

    def put_gradients(self):
        return ray.put(self.gradients)

    def reduce_gradients(self, grad_id_refs):
        grad_ids = ray.get(grad_id_refs)
        reduced_result = cupy.ones((10,), dtype=float32)
        for grad_id in grad_ids:
            array = ray.get(grad_id)
            reduced_result += array
        result_id = ray.put(reduced_result)
        return result_id

    def get_reduced_gradient(self, reduced_gradient_id_ref):
        reduced_gradient_id = ray.get(reduced_gradient_id_ref)
        reduced_gradient = ray.get(reduced_gradient_id)
        # do whatever with the reduced gradients
        return True

# Allreduce the gradients using Ray APIs
# Let all workers to put their gradients into the Ray object store.
gradient_ids = [worker.put_gradients.remote() for worker in workers]
ray.wait(gradient_ids, num_returns=len(gradient_ids, timeout=None))

# Let worker 0 reduce the gradients
reduced_id_ref = workers[0].reduce_gradients.remote(gradient_ids)

# All others workers get the reduced gradients
results = []
for i, worker in enumerate(workers):
    results.append(worker.get_reduced_gradient.remote([reduced_id_ref]))
ray.get(results)
```

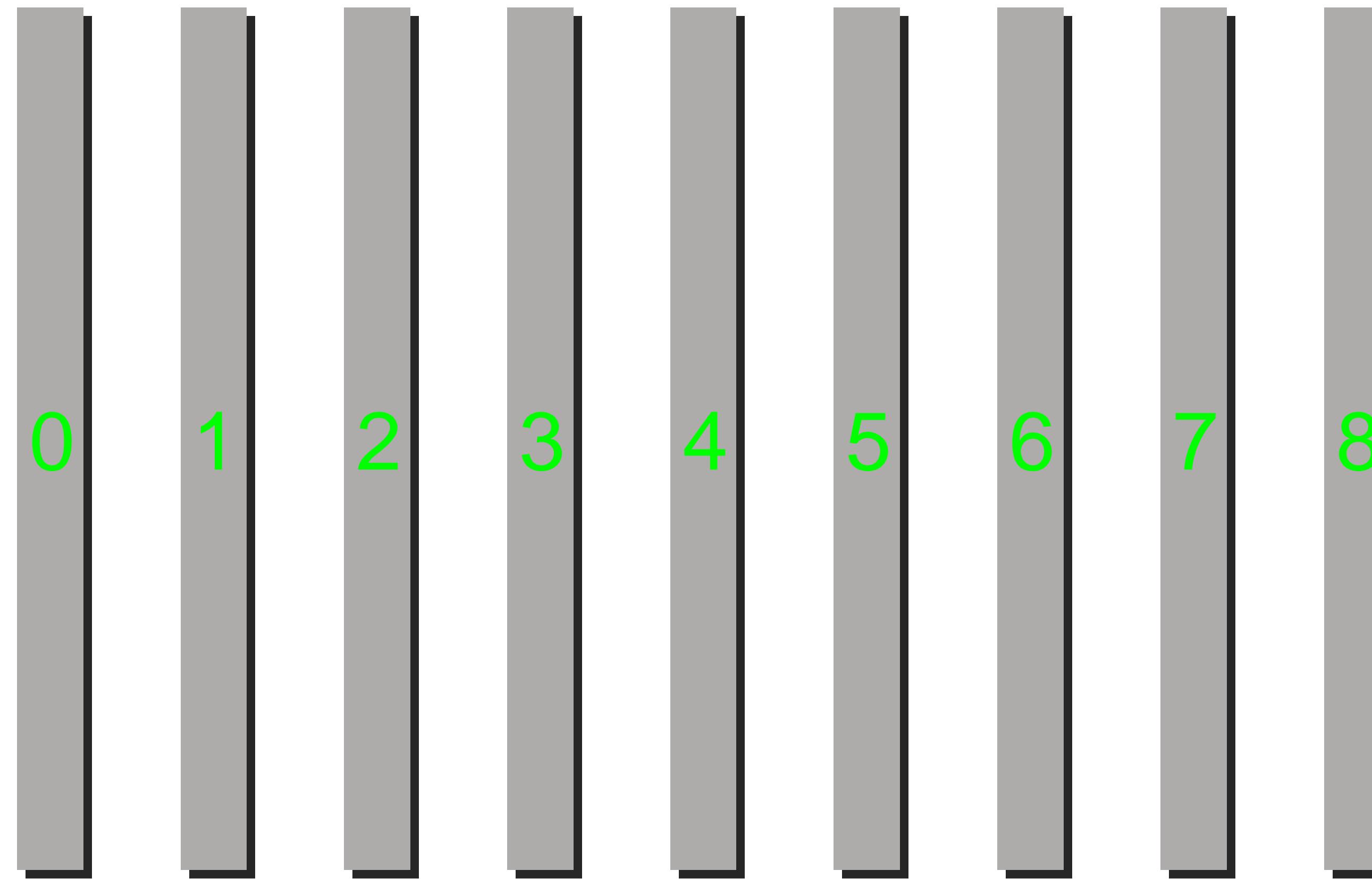
Performance

- Message size over networks:
 - Sum: $3N$
 - Send Sum back: $3N$
 - $= 6N$
- Can we do better?
 - Hint: we cannot do better than $3N$

Why Collective Communication?

- Programming Convenience
 - Use a set of well-defined communication primitives to express complex communication patterns
- Performance
 - Since they are well defined and well structured, we can optimize them to the extreme
- ML Systems ❤️ Collective communication

Make it Formal



- A 1D **Mesh** of workers (or devices, or nodes)

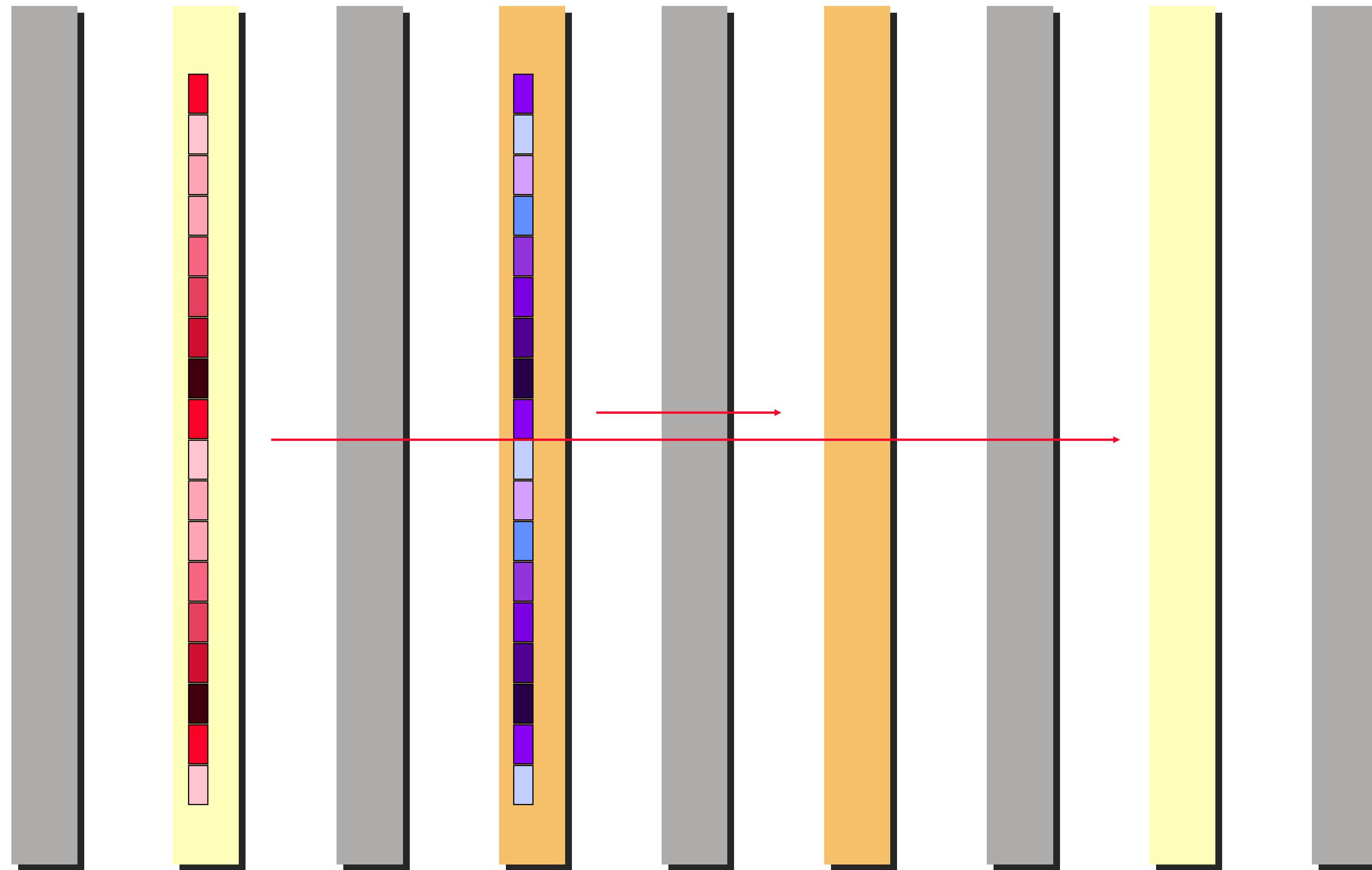
Model of Parallel Computation

- a node can send directly to any other node (maybe not true)
- a node can simultaneously receive and send
- cost of communication
 - sending a message of length n between any two nodes

$$\alpha + n \beta$$

- if a message encounters a link that simultaneously accommodates M messages, the cost becomes

$$\alpha + Mn \beta$$

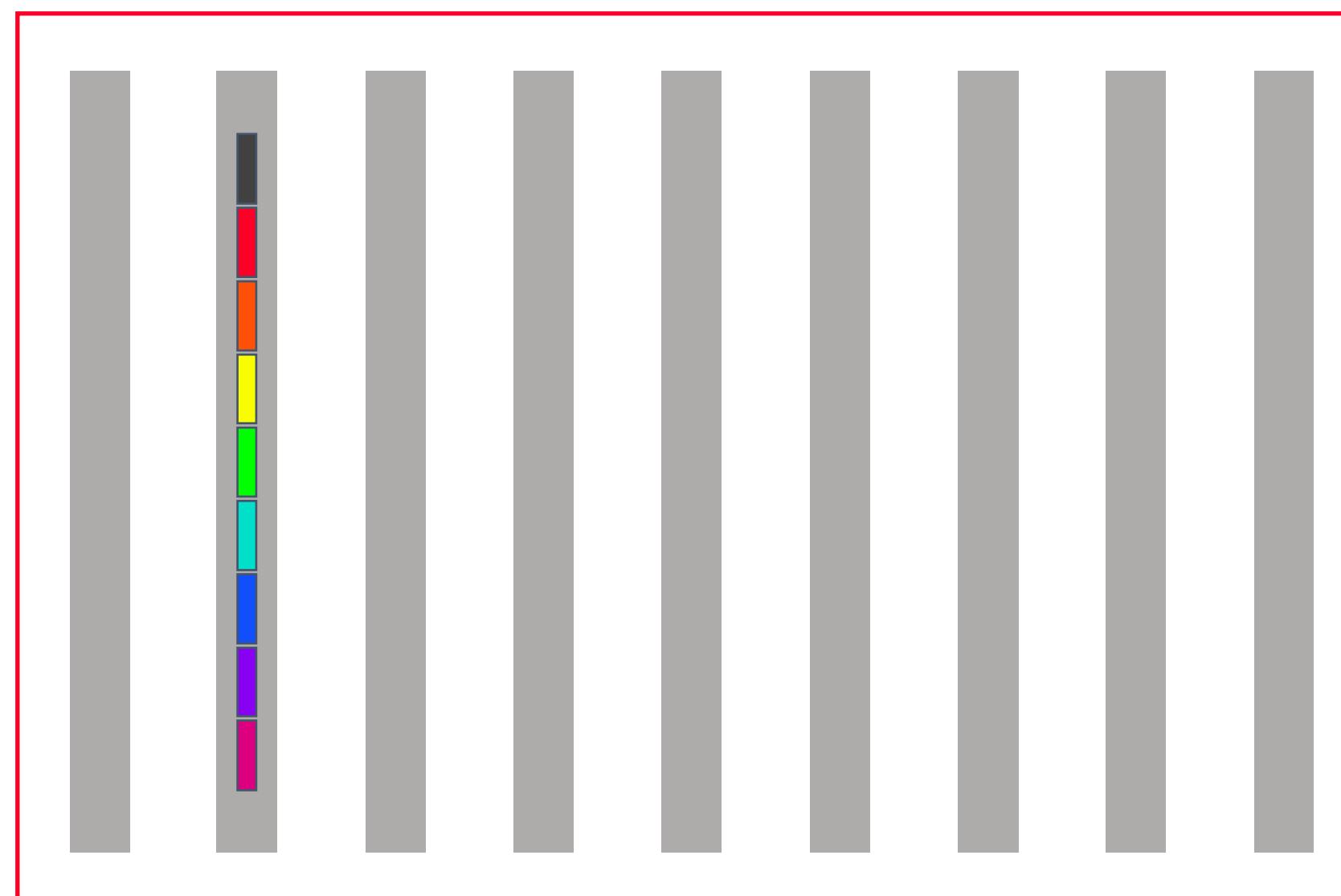


Collective Communications

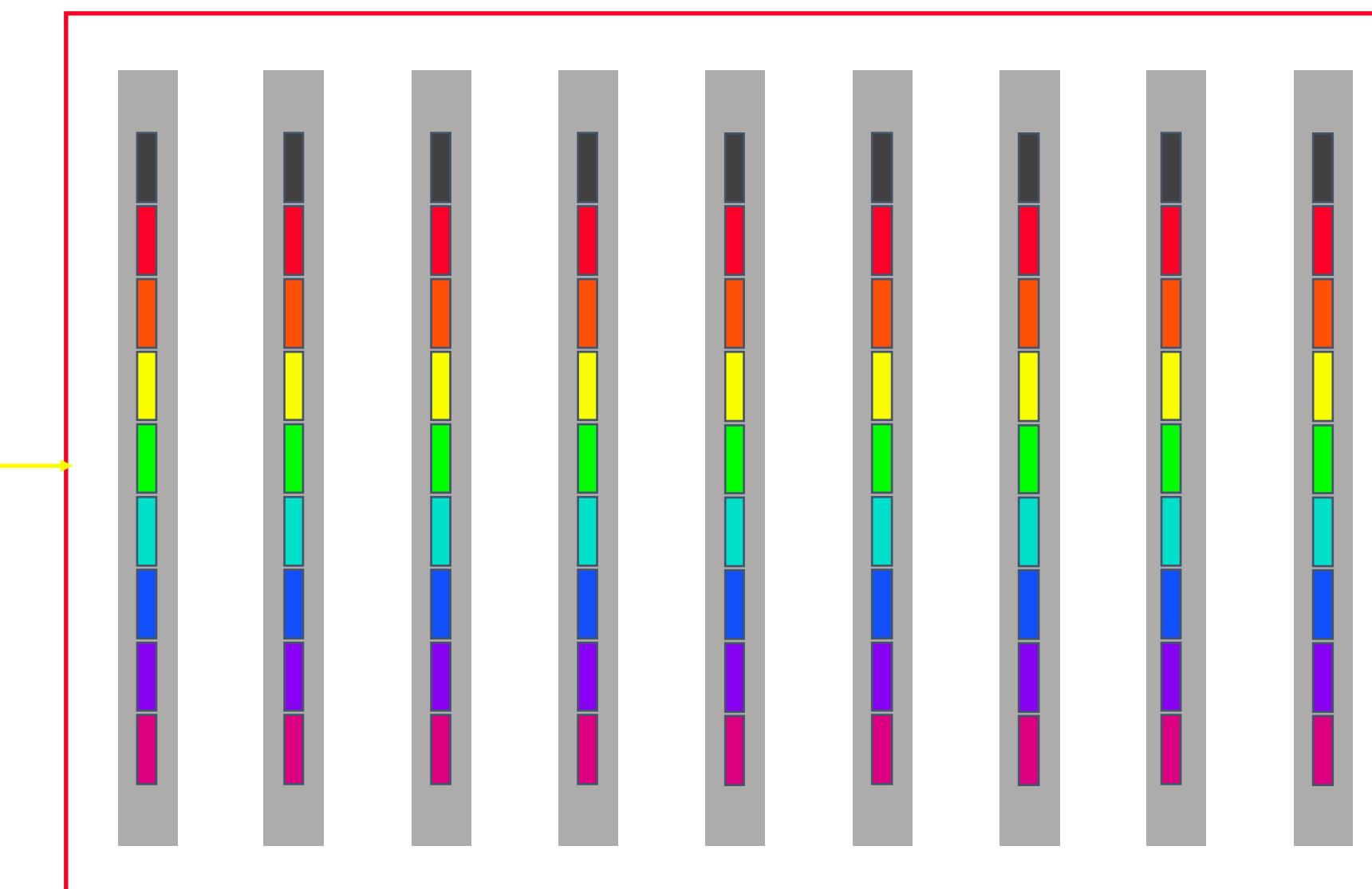
- Broadcast
- Reduce(-to-one)
- Scatter
- Gather
- Allgather
- Reduce-scatter
- Allreduce

Broadcast

Before

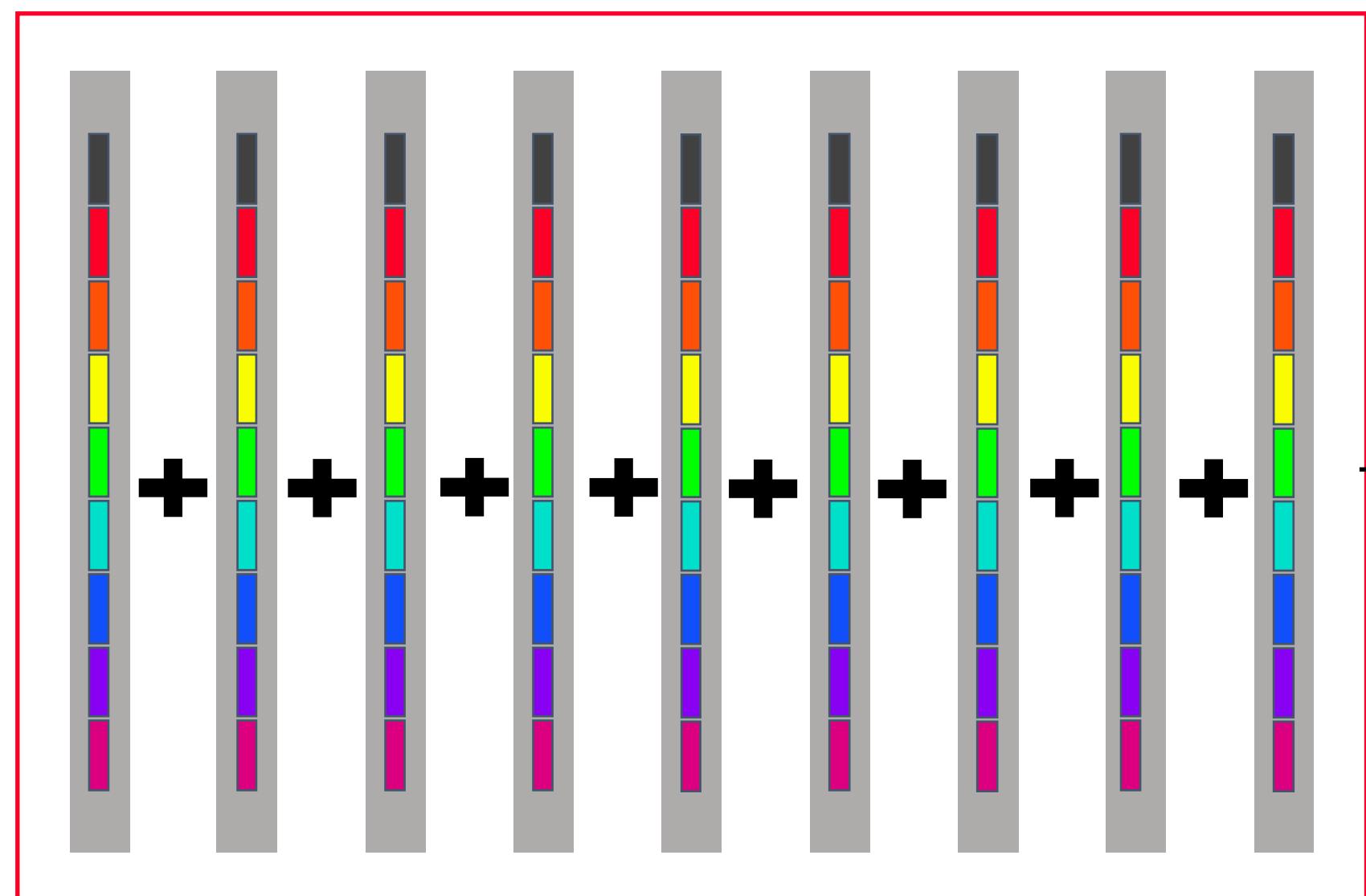


After

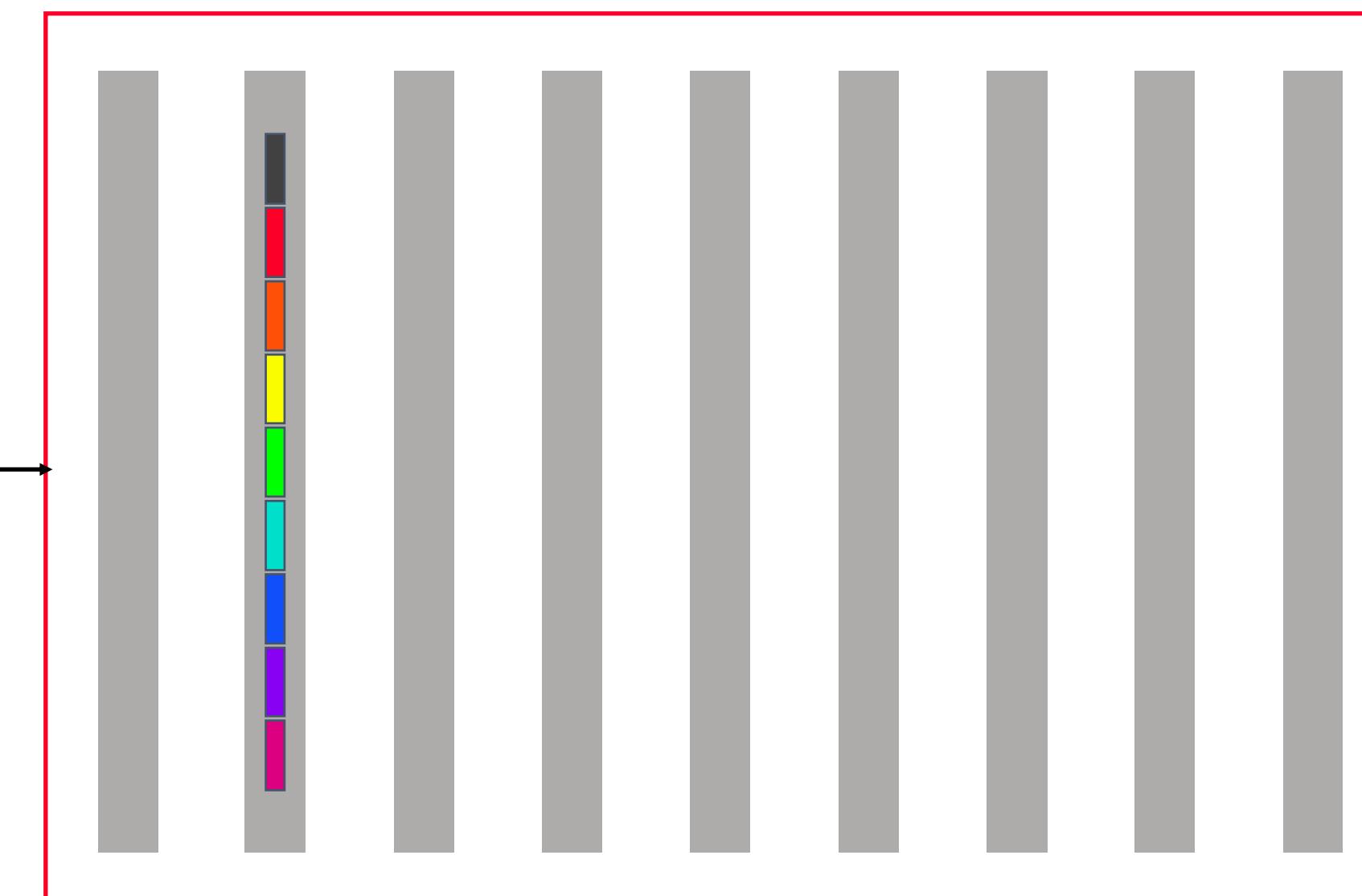


Reduce(-to-one)

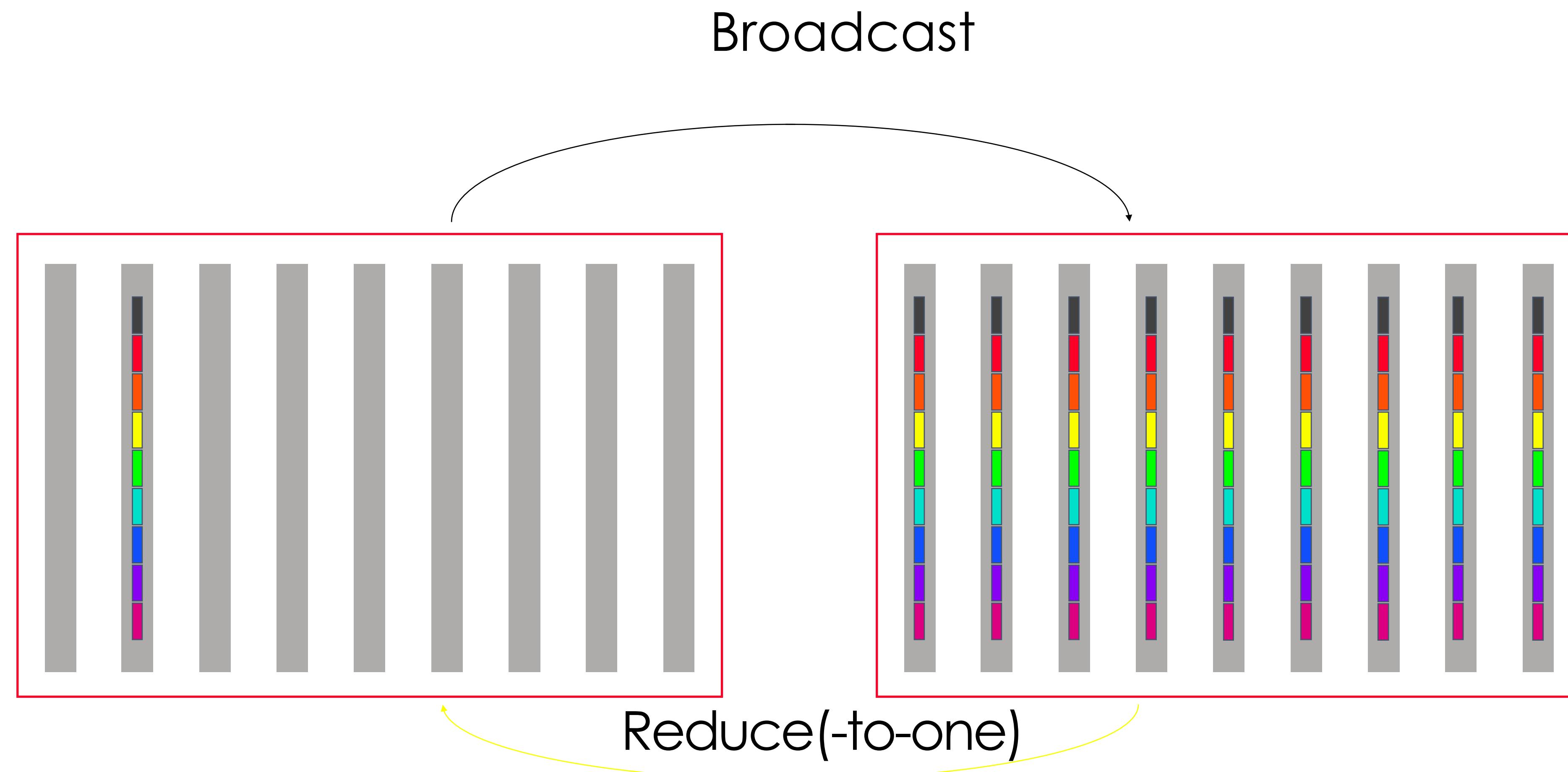
Before



After

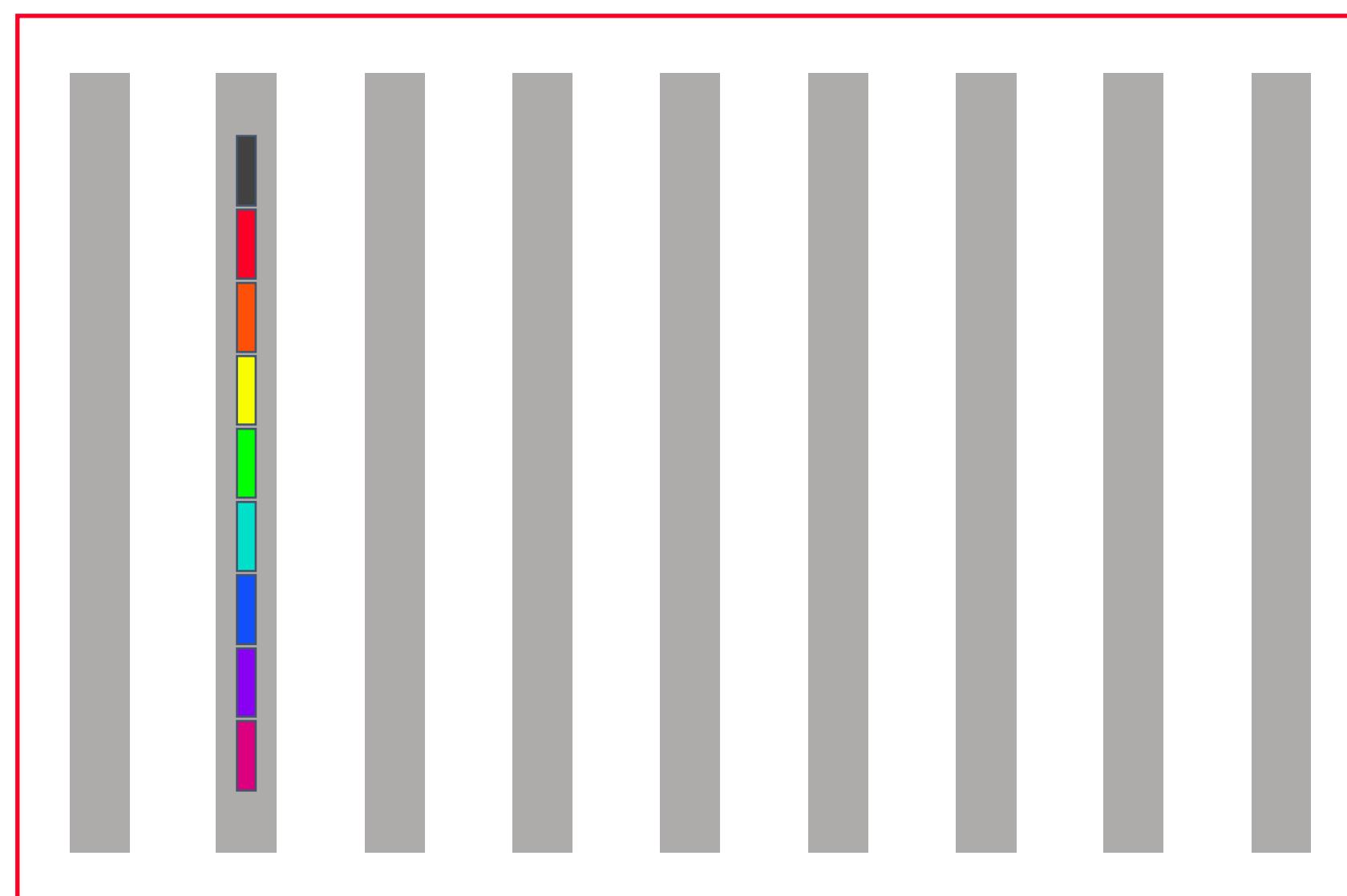


Broadcast/Reduce(-to-one)

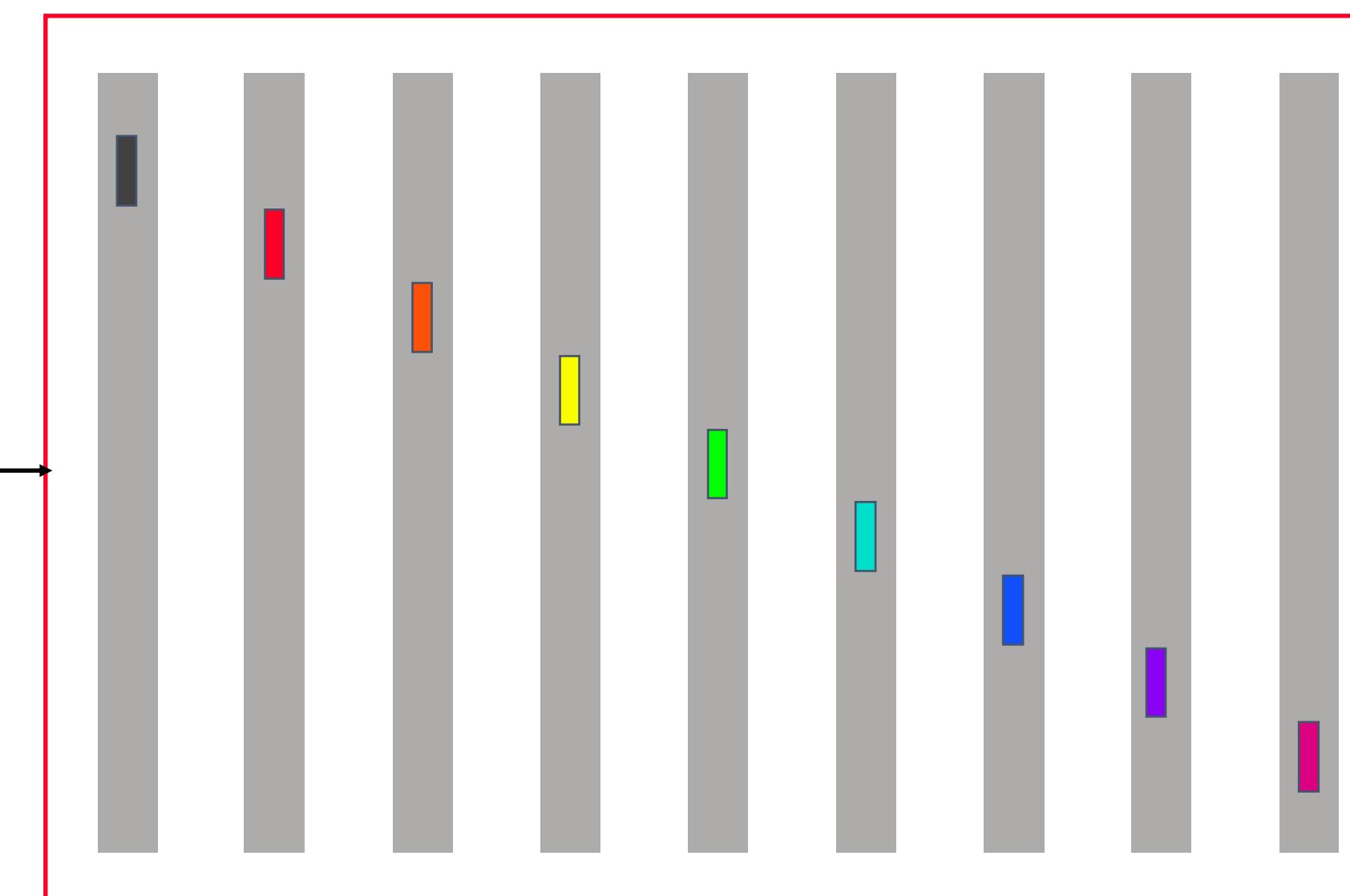


Scatter

Before

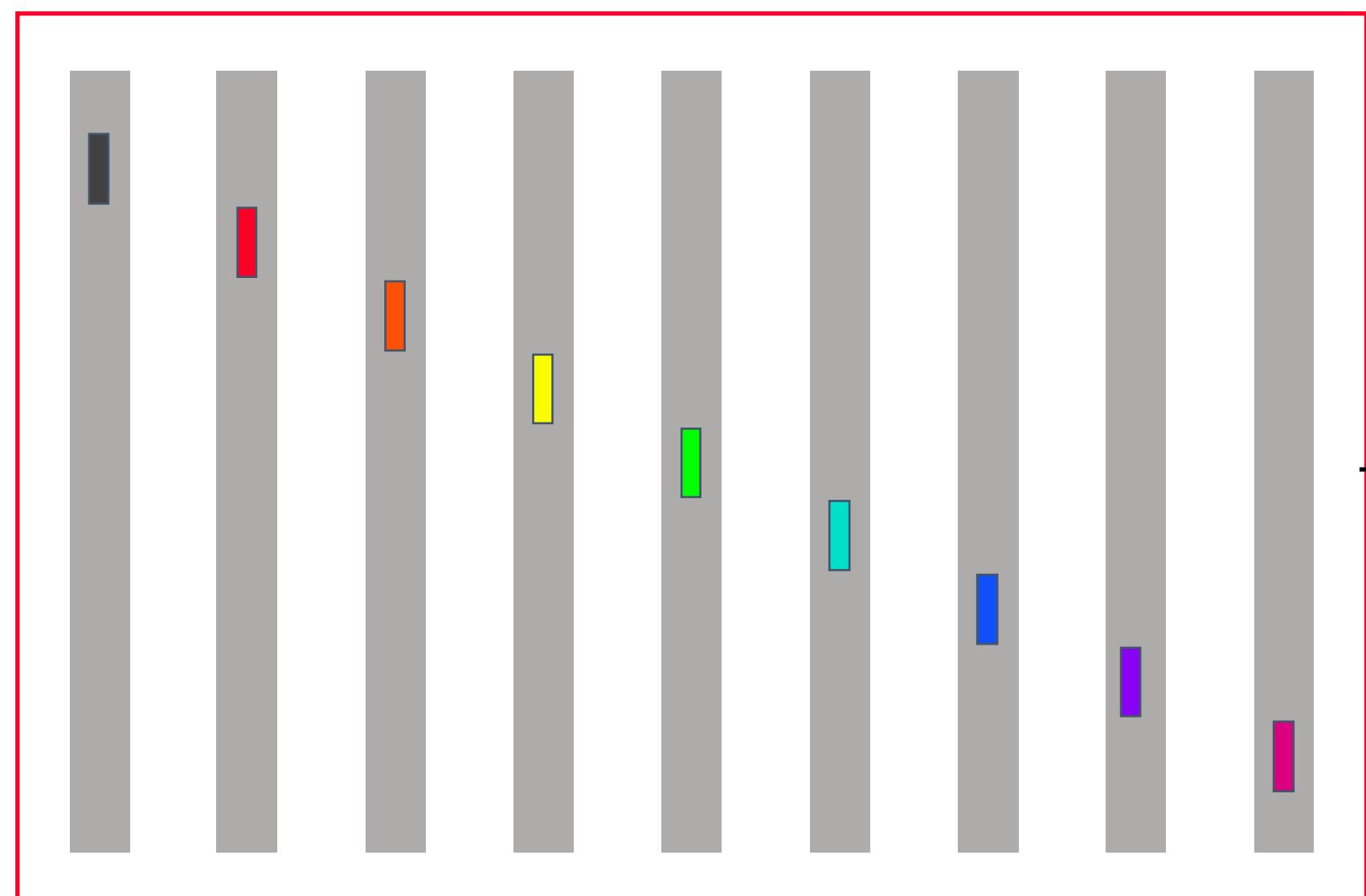


After

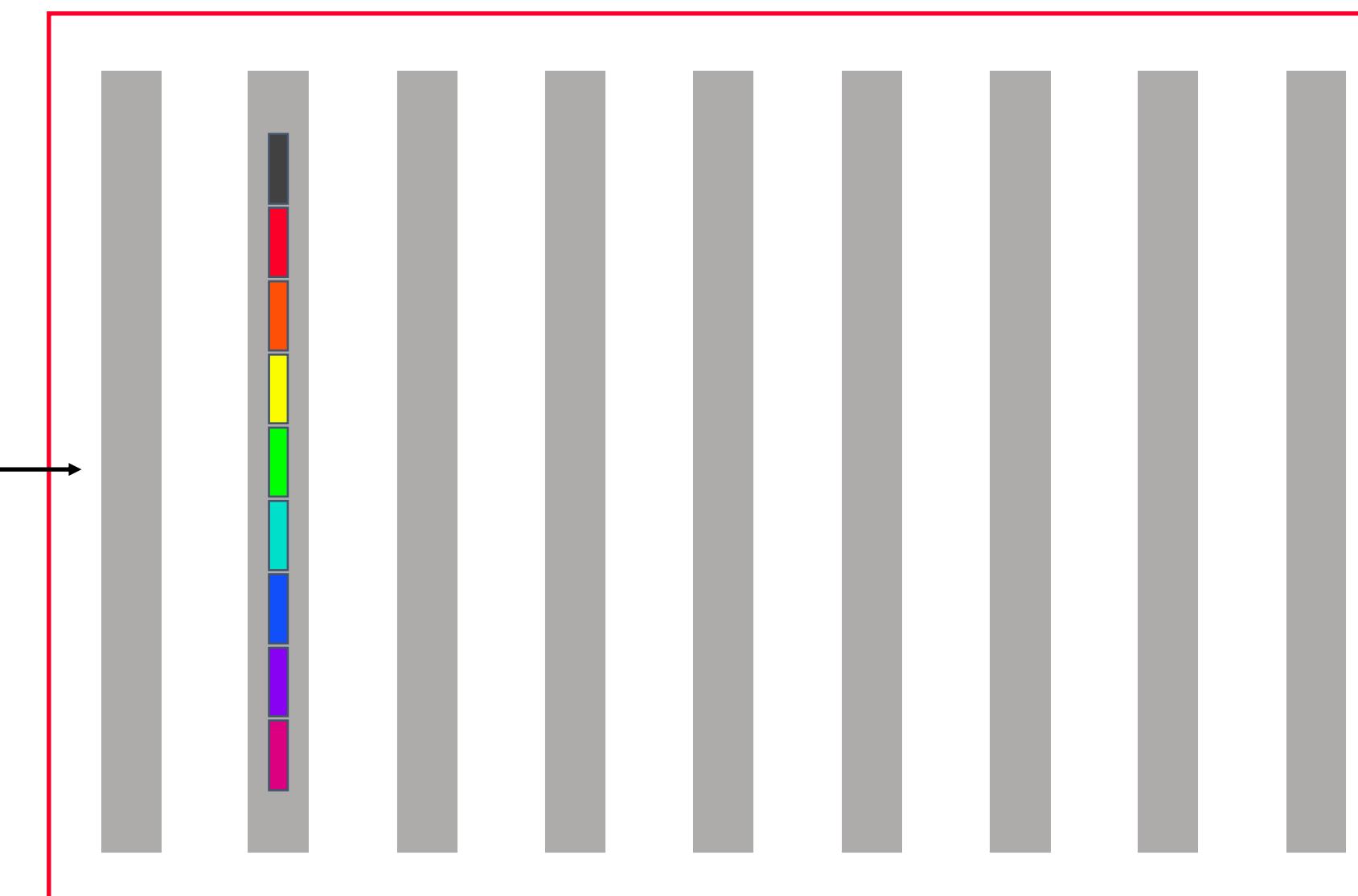


Gather

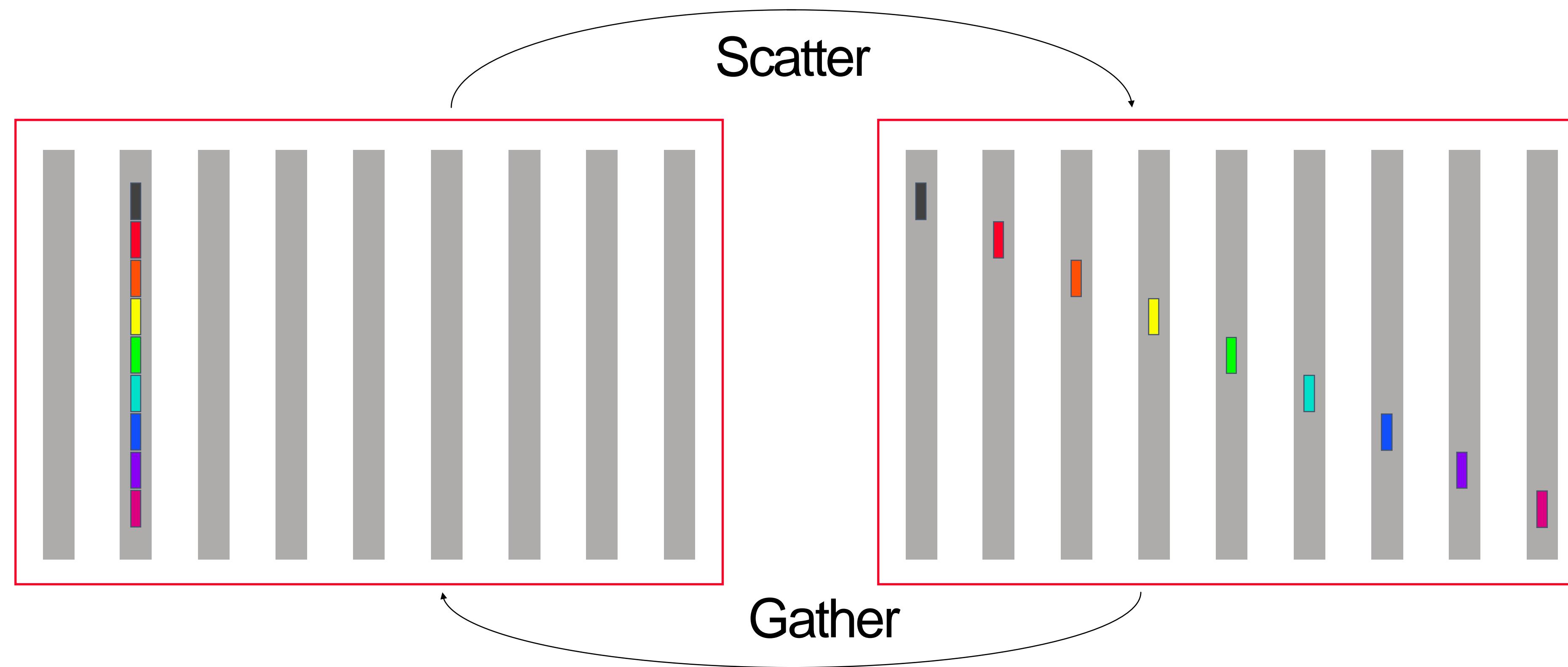
Before



After

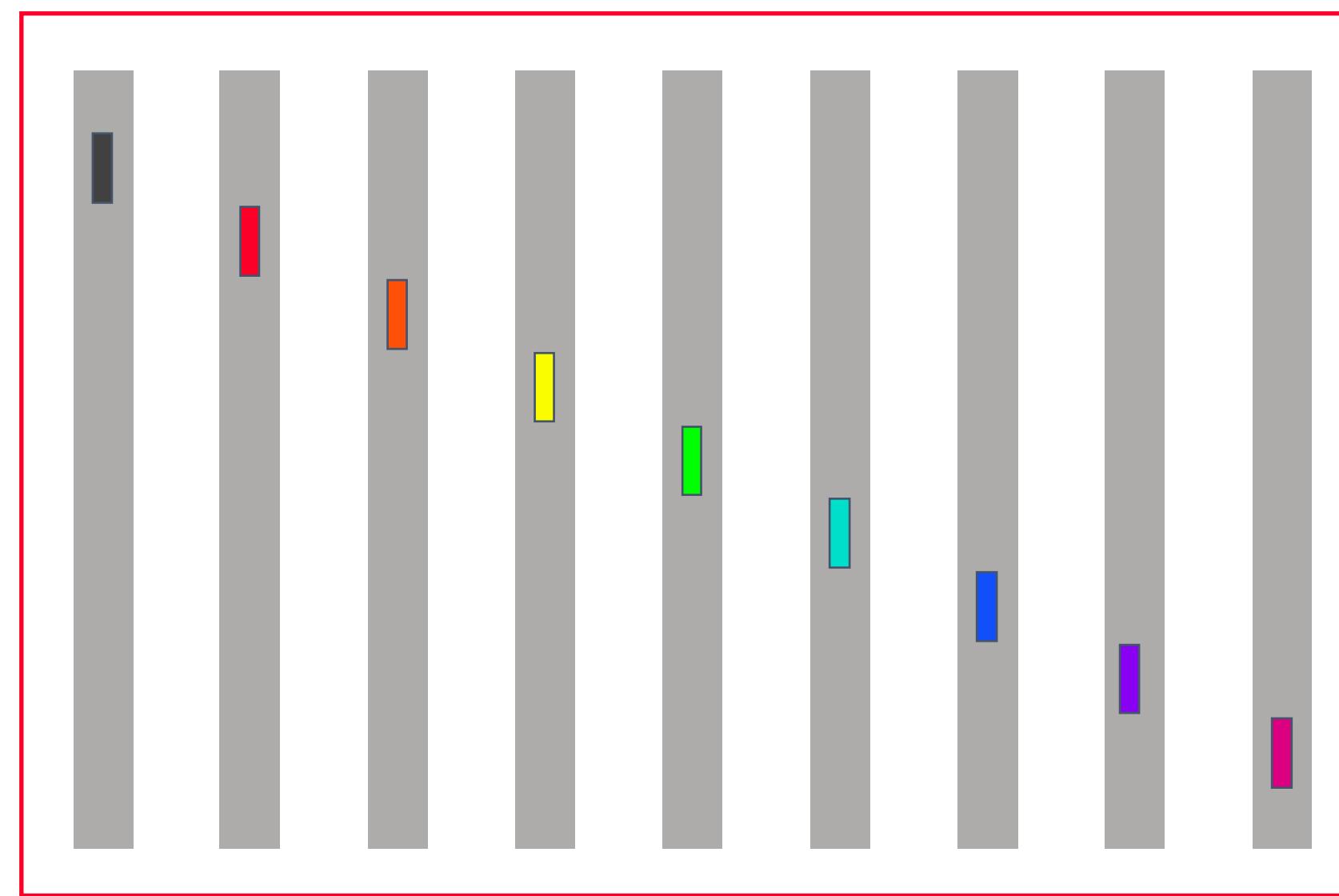


Scatter/Gather

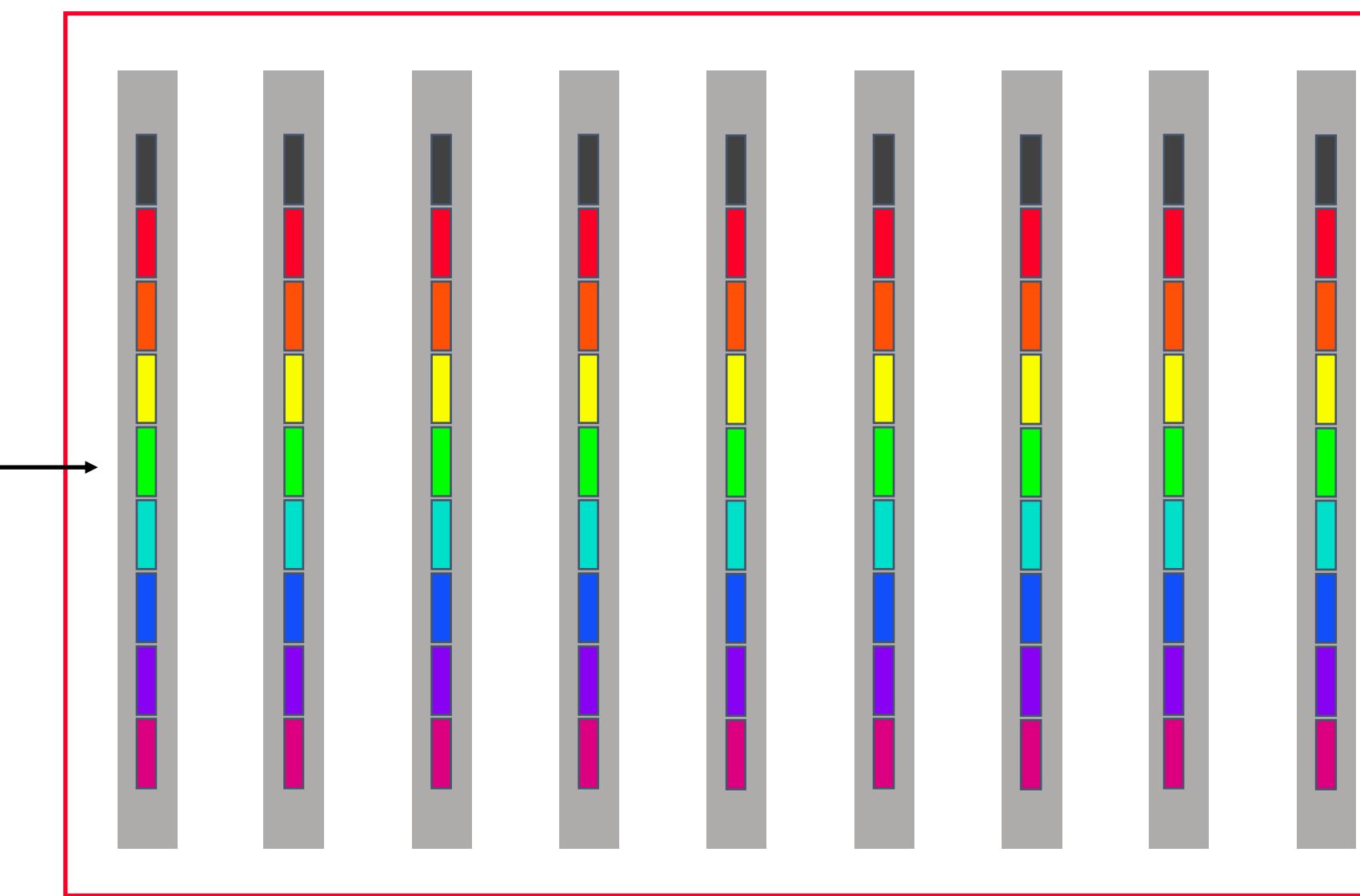


Allgather

Before

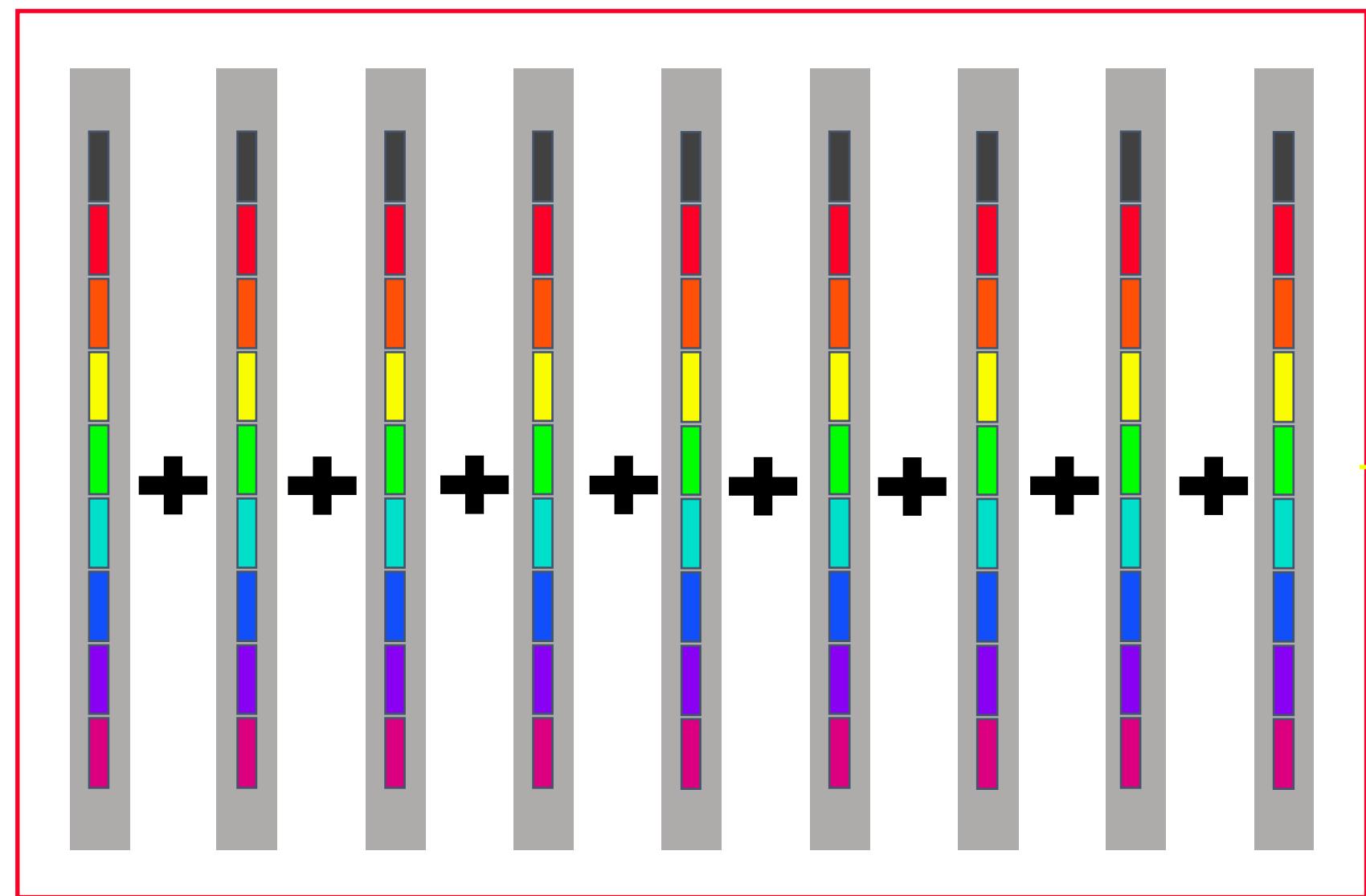


After

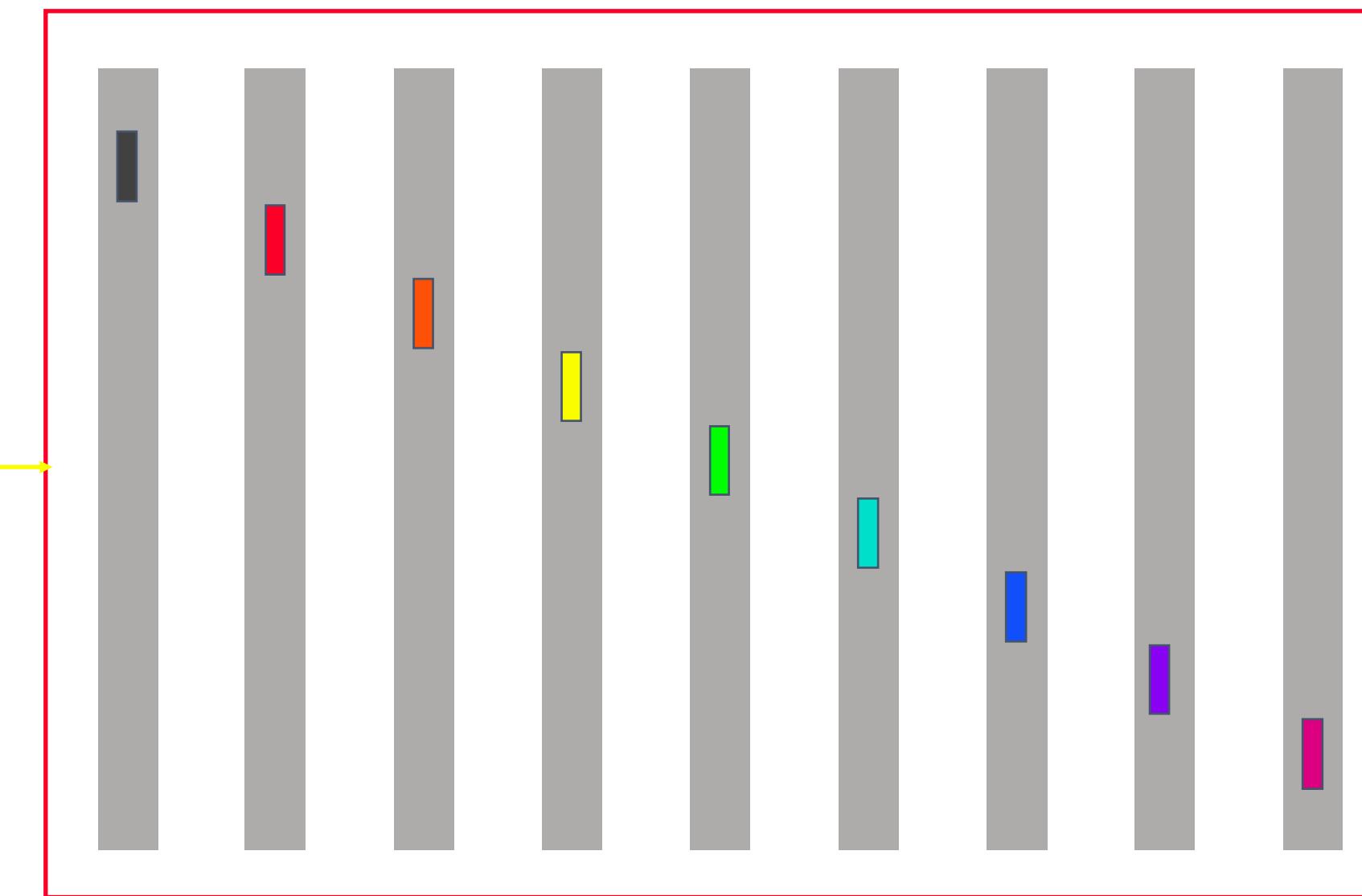


Reduce-scatter

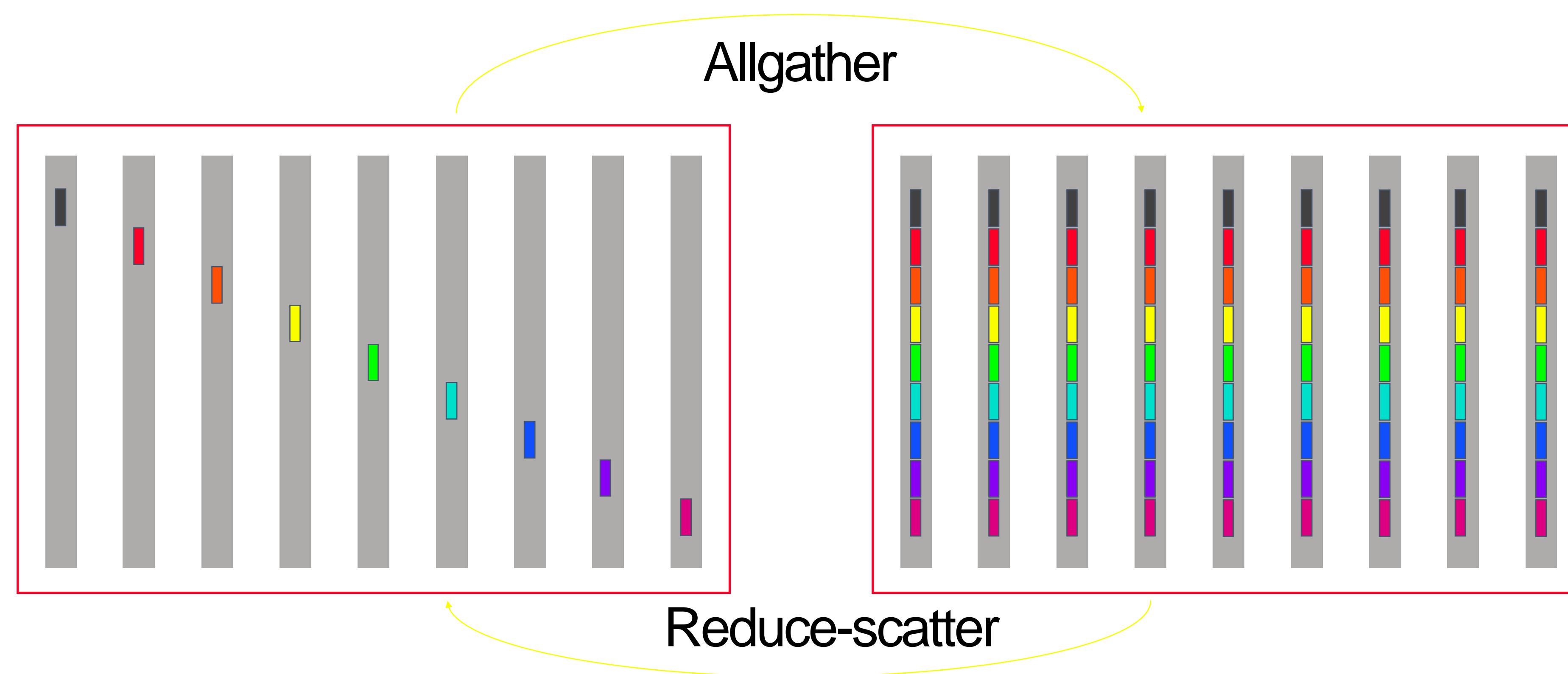
Before



After

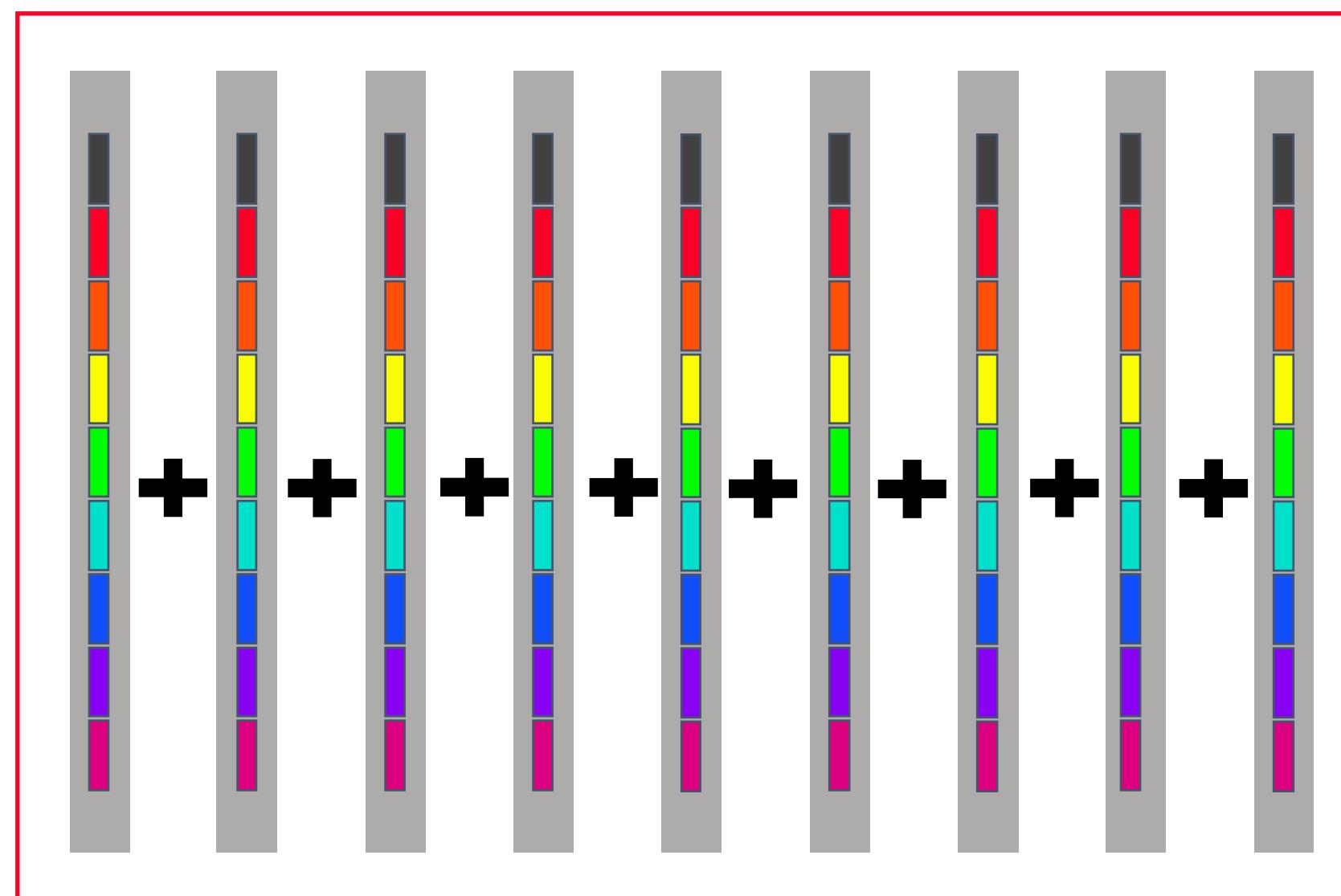


Allgather/Reduce-scatter

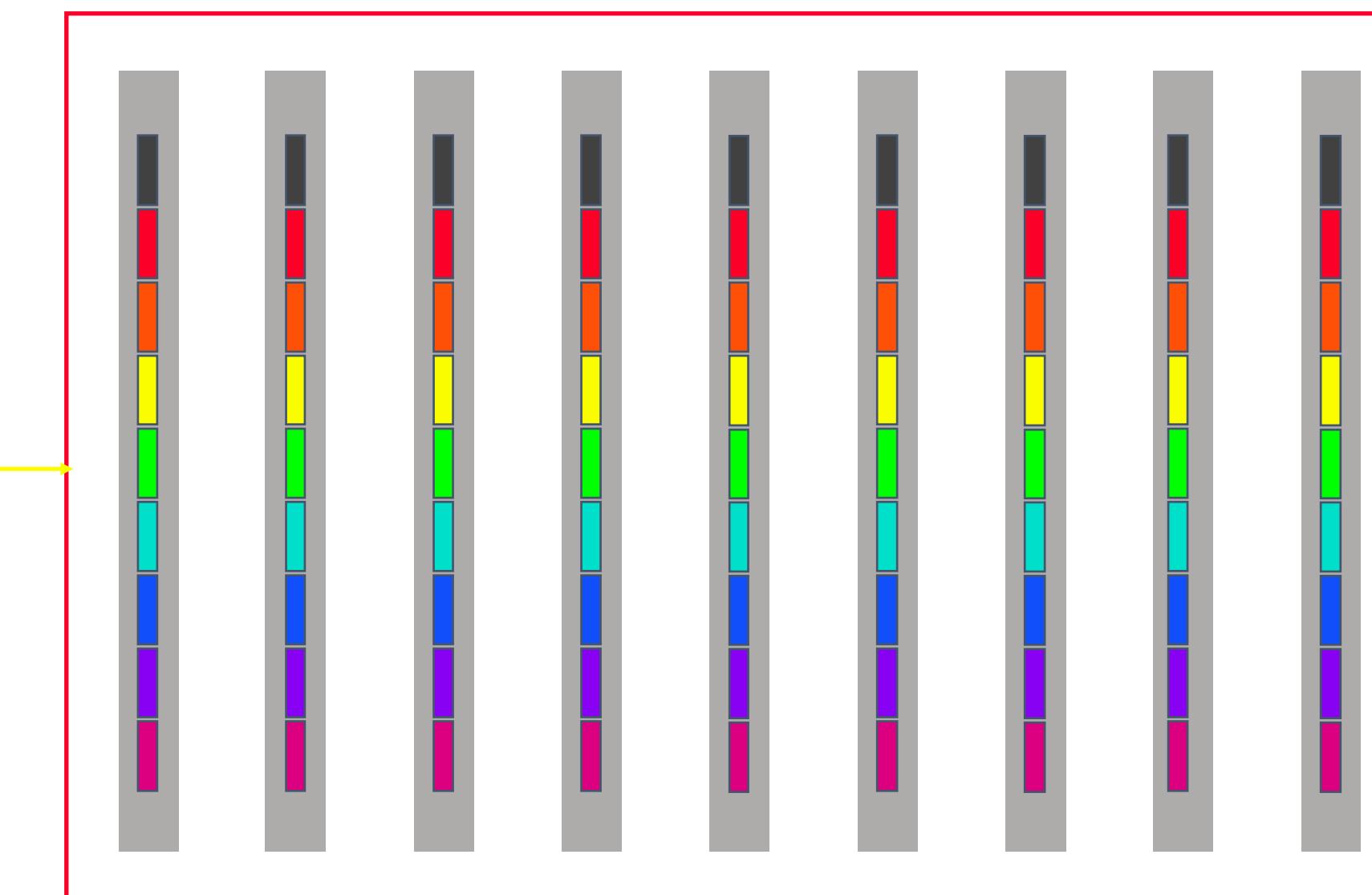


Allreduce

Before



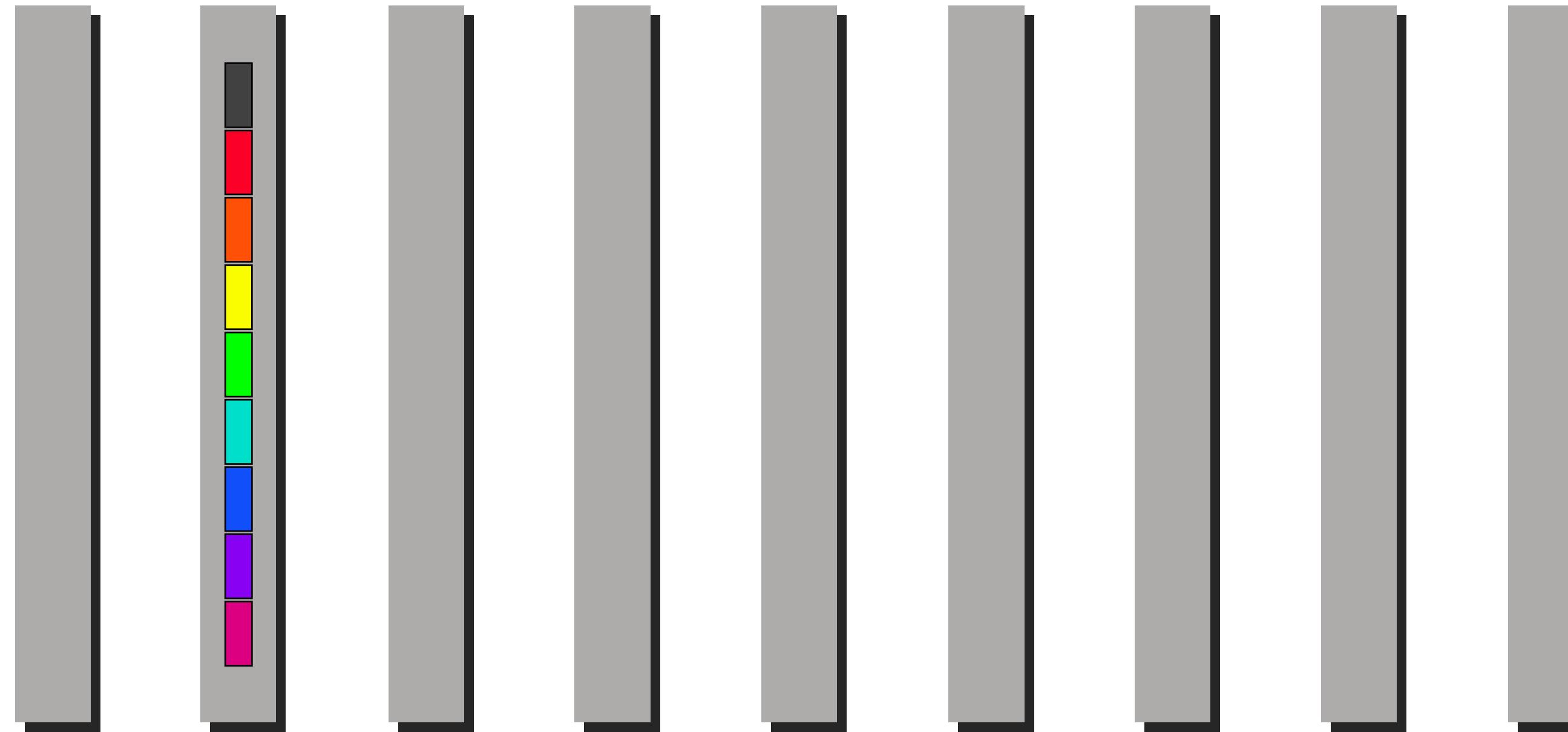
After



Two Family of Mainstream Algorithms/Implementations

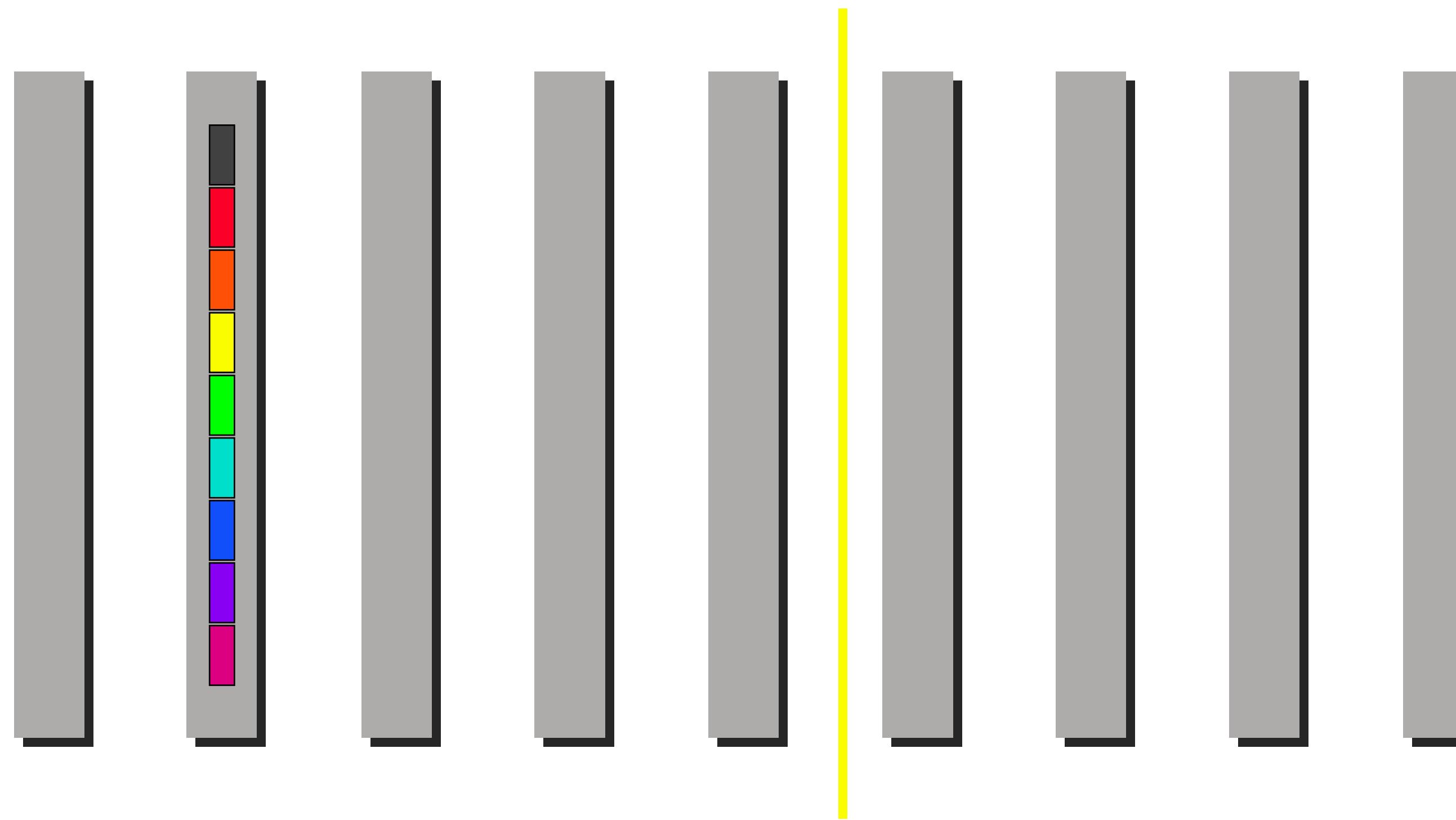
- Small message: Minimum Spanning Tree algorithm
 - Emphasize **low latency**
- Large Message: Ring algorithm
 - Emphasize **bandwidth utilization**
- There are 50+ different algorithms developed in the past 50 years by a community called “High-performance computing”
 - Last year Turing award

General principles



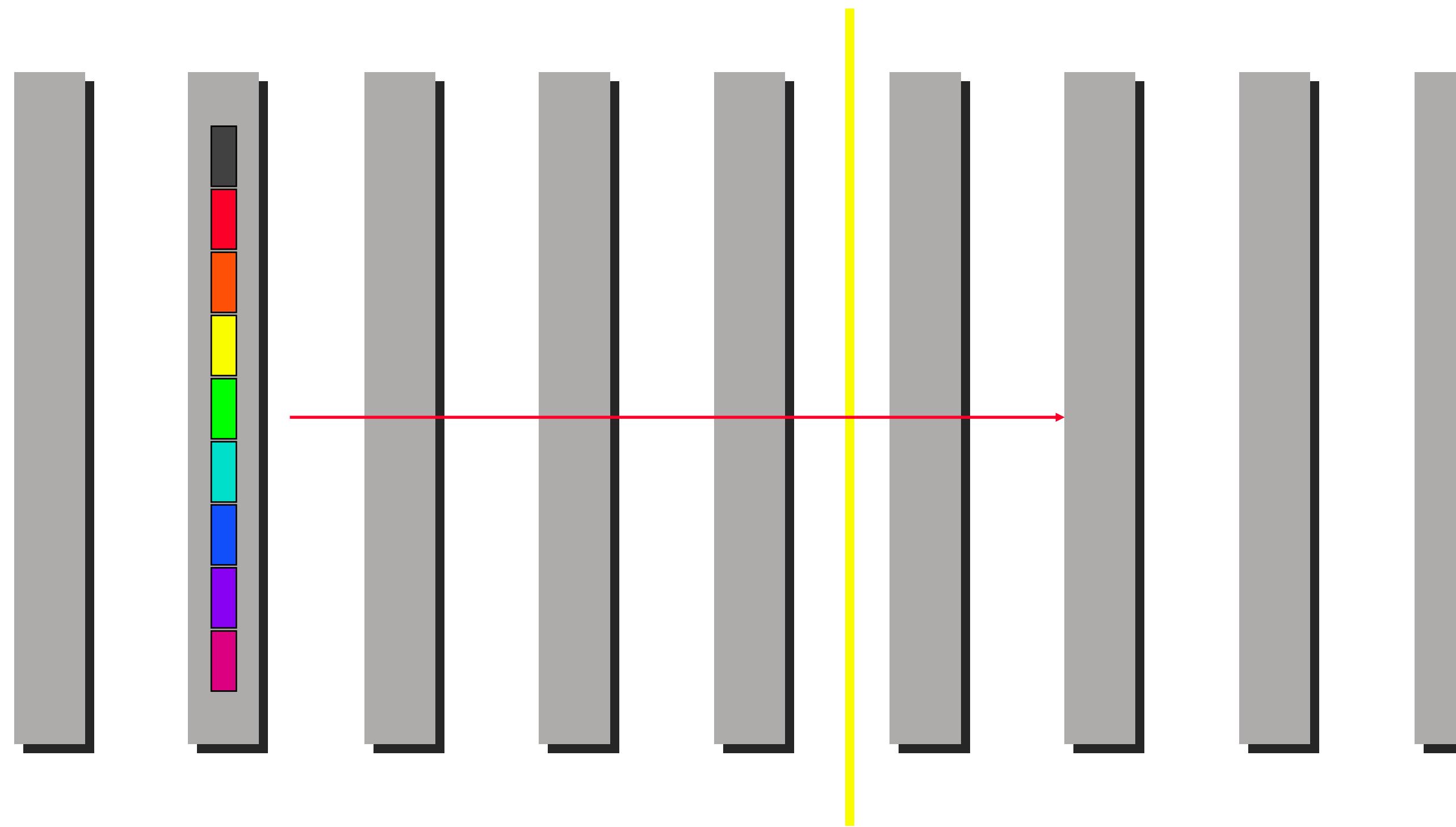
- message starts on one processor

General principles



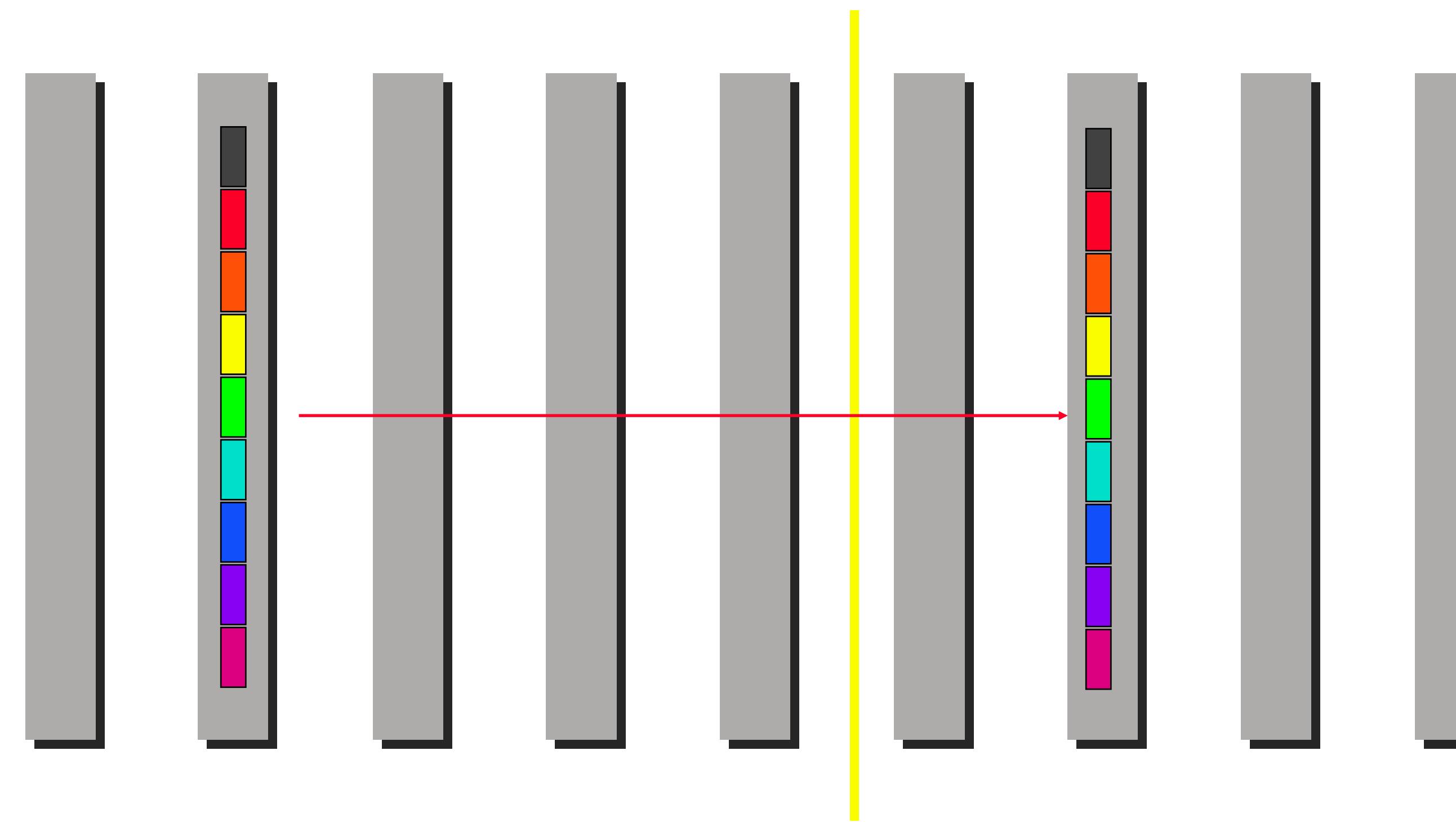
- divide logical linear array in half

General principles



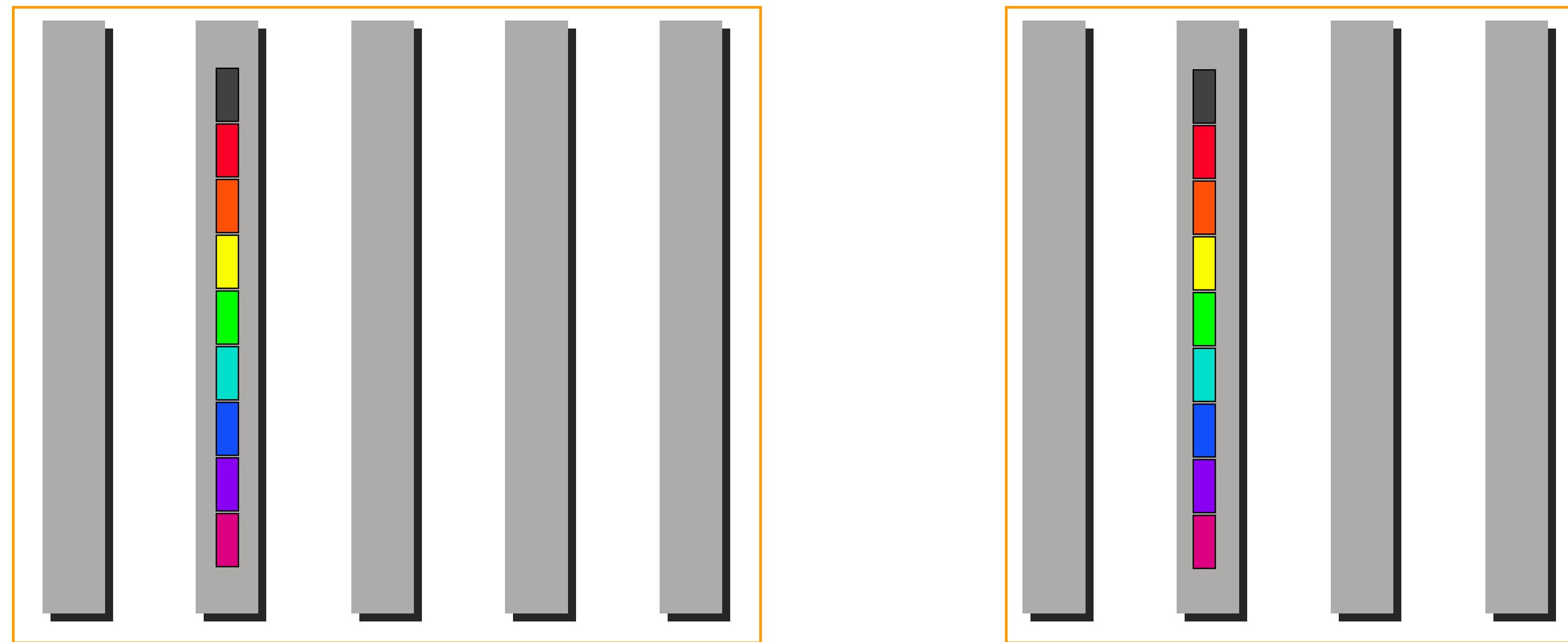
- send message to the half of the network that does not contain the current node (root) that holds the message

General principles



- send message to the half of the network that does not contain the current node (root) that holds the message

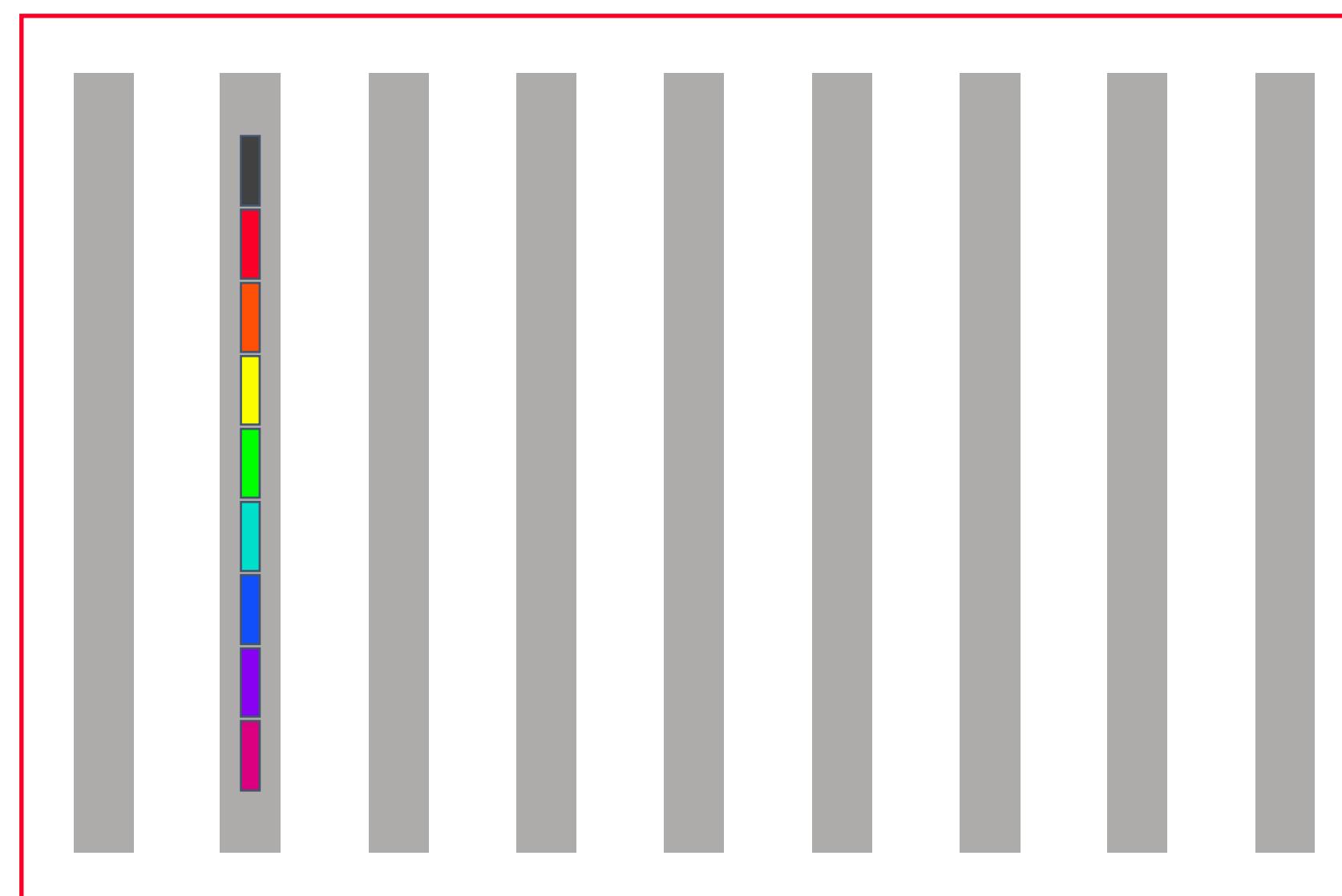
General principles



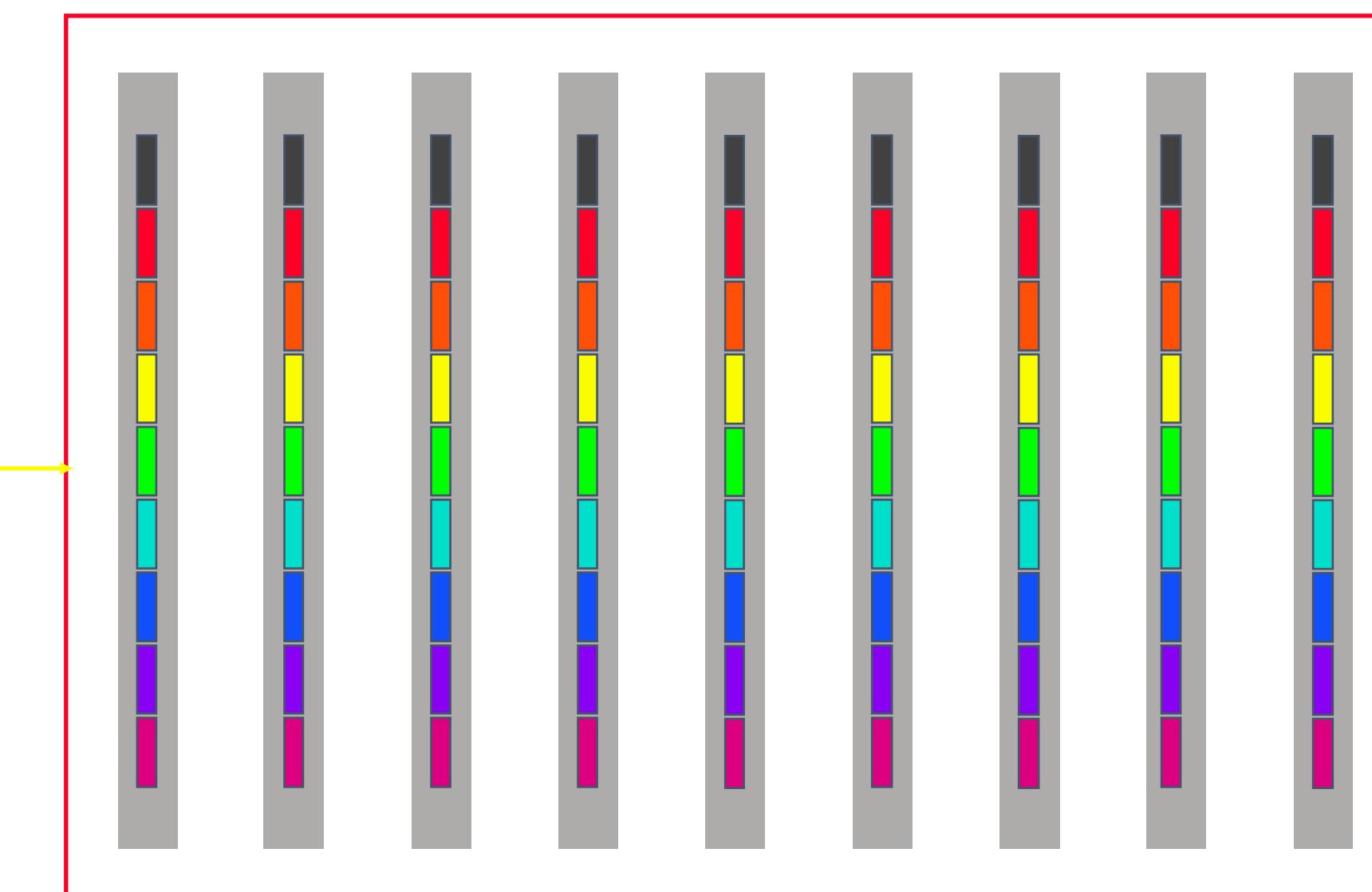
- continue recursively in each of the two halves

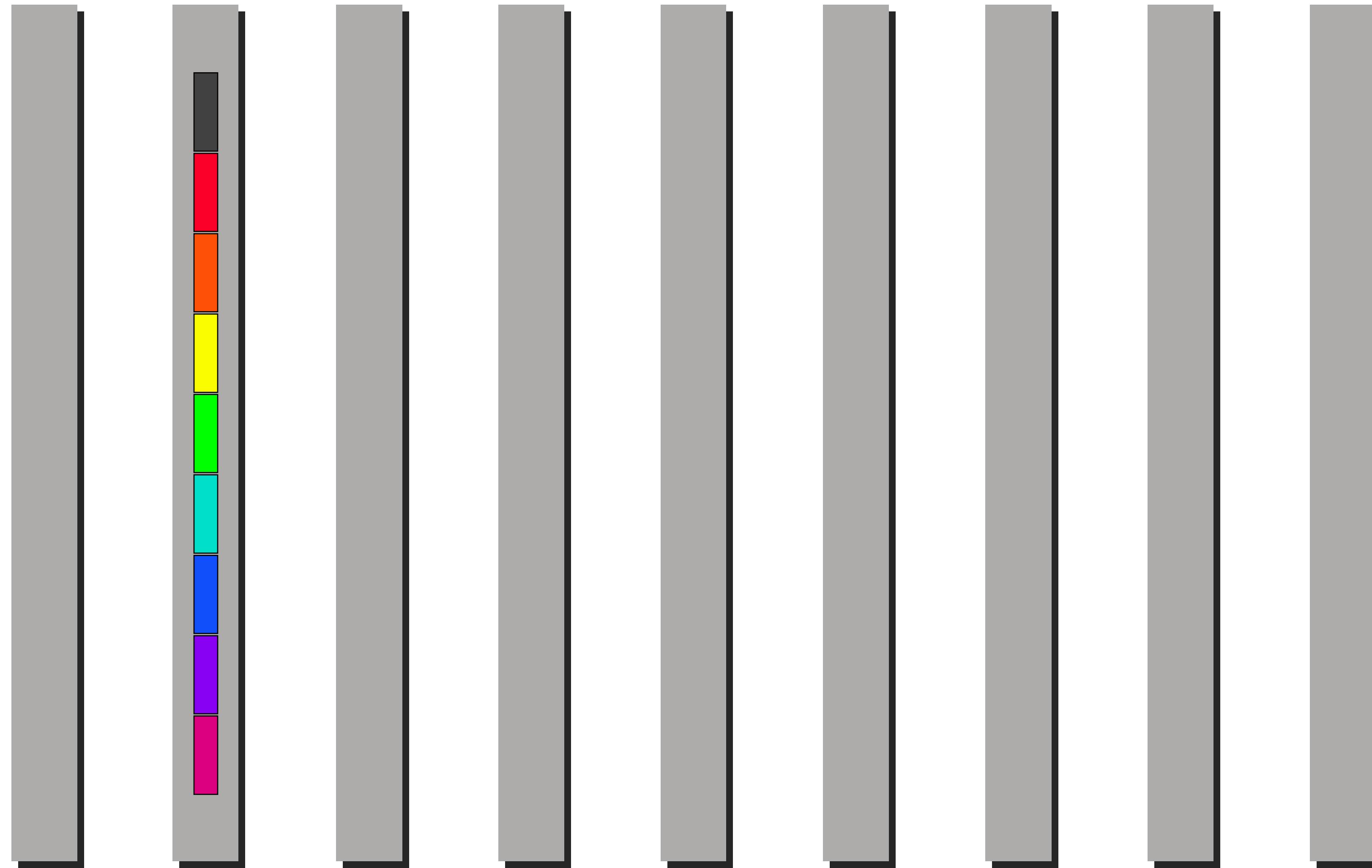
Broadcast

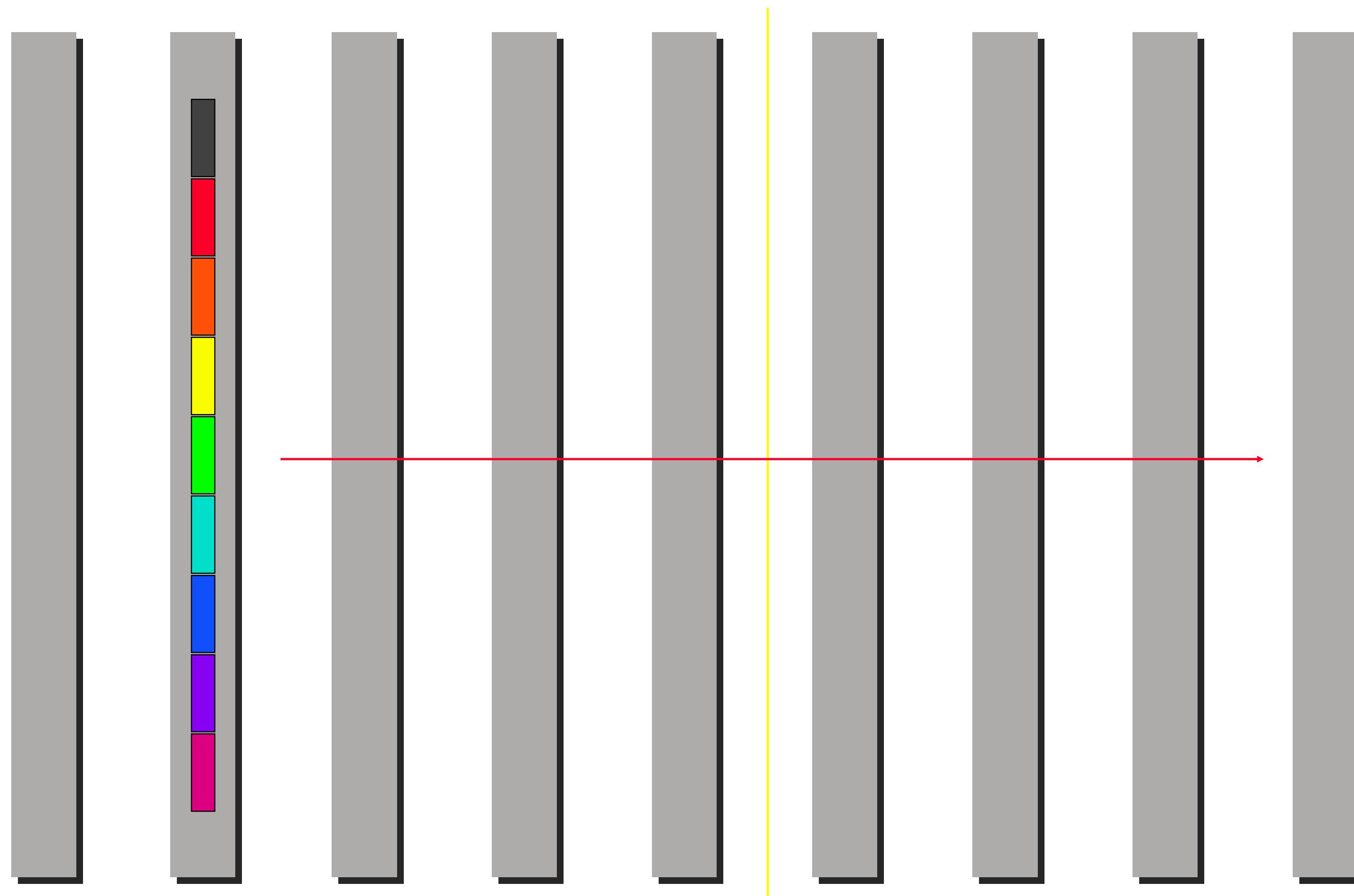
Before

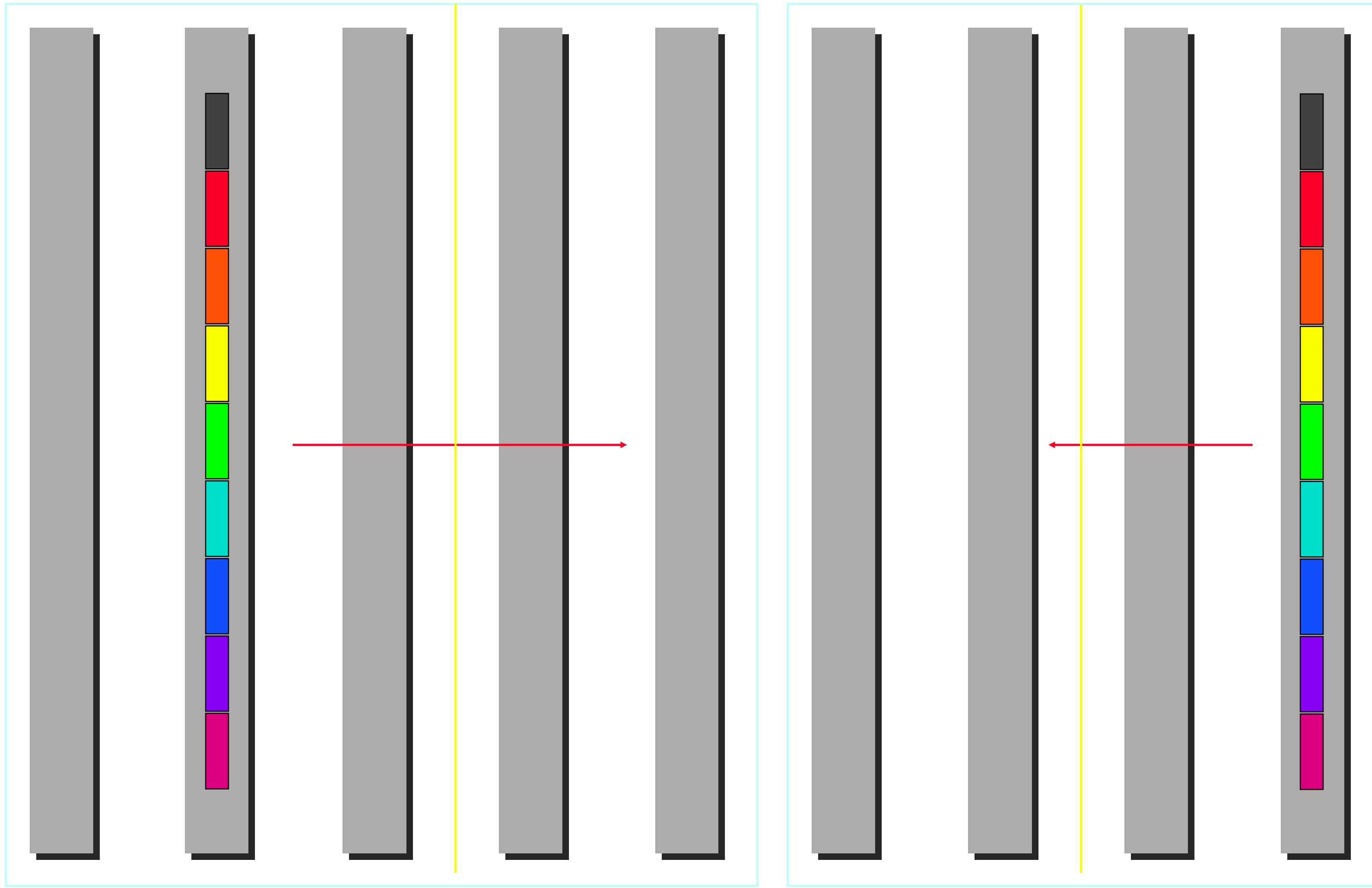


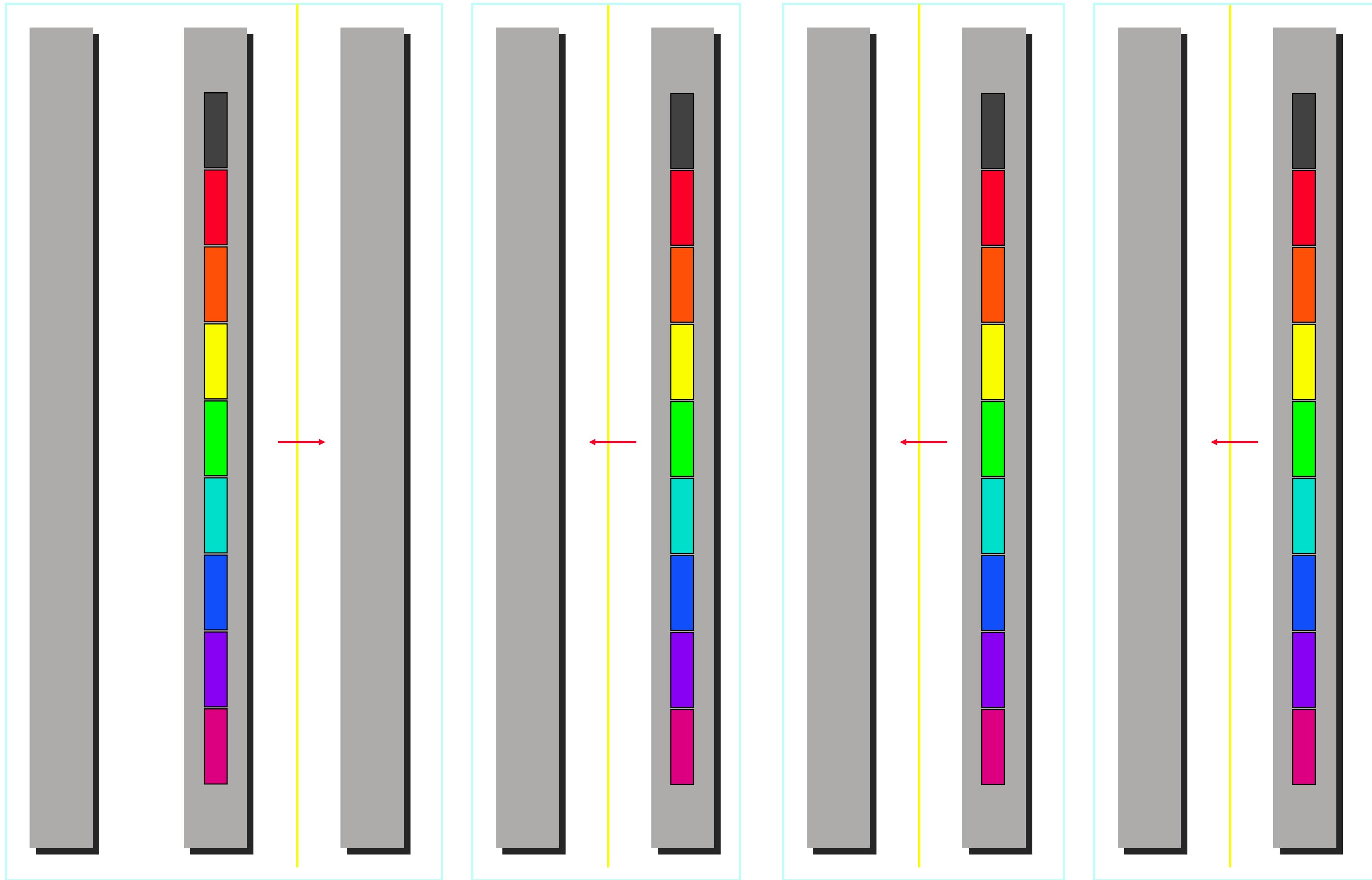
After

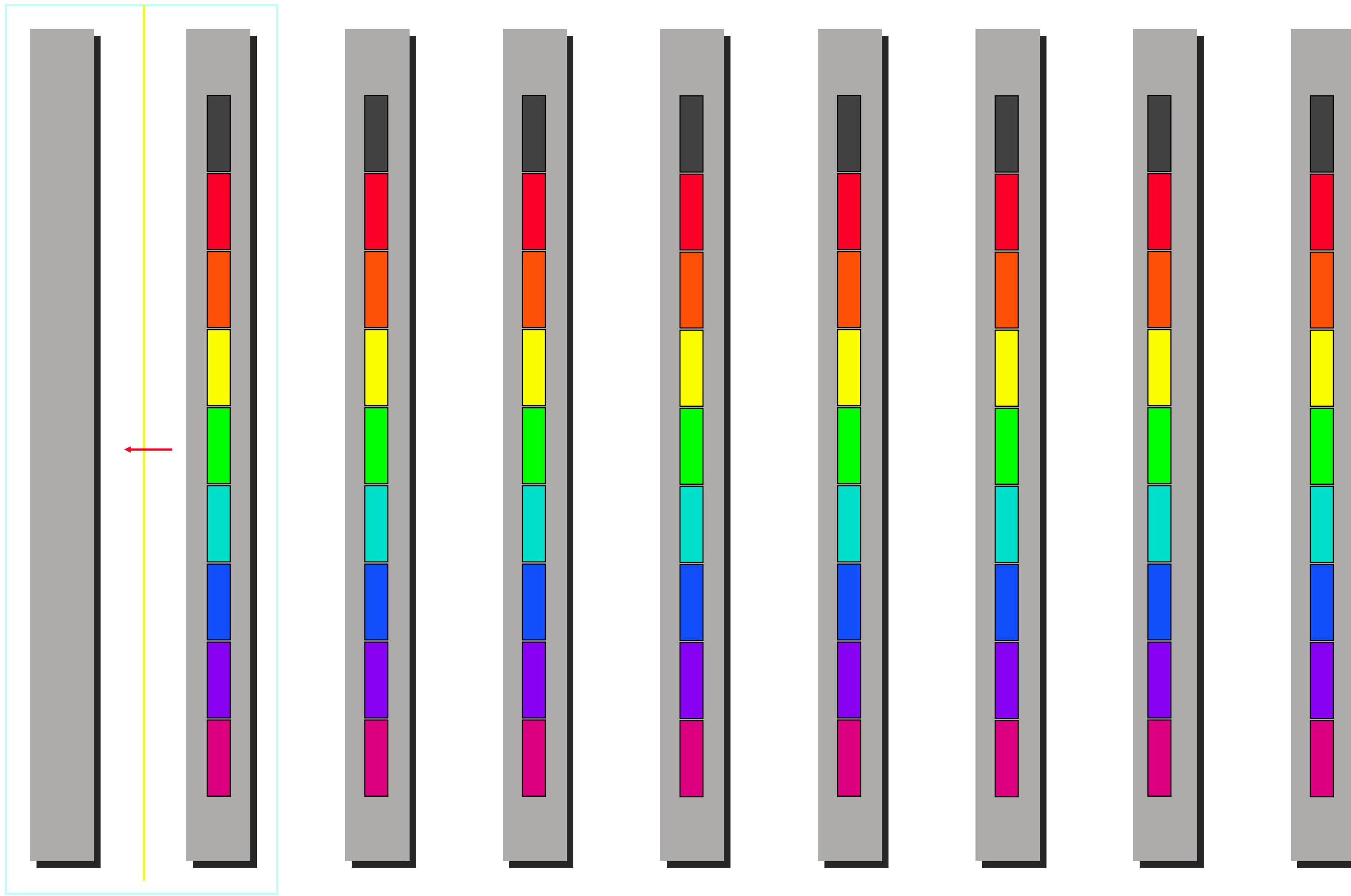


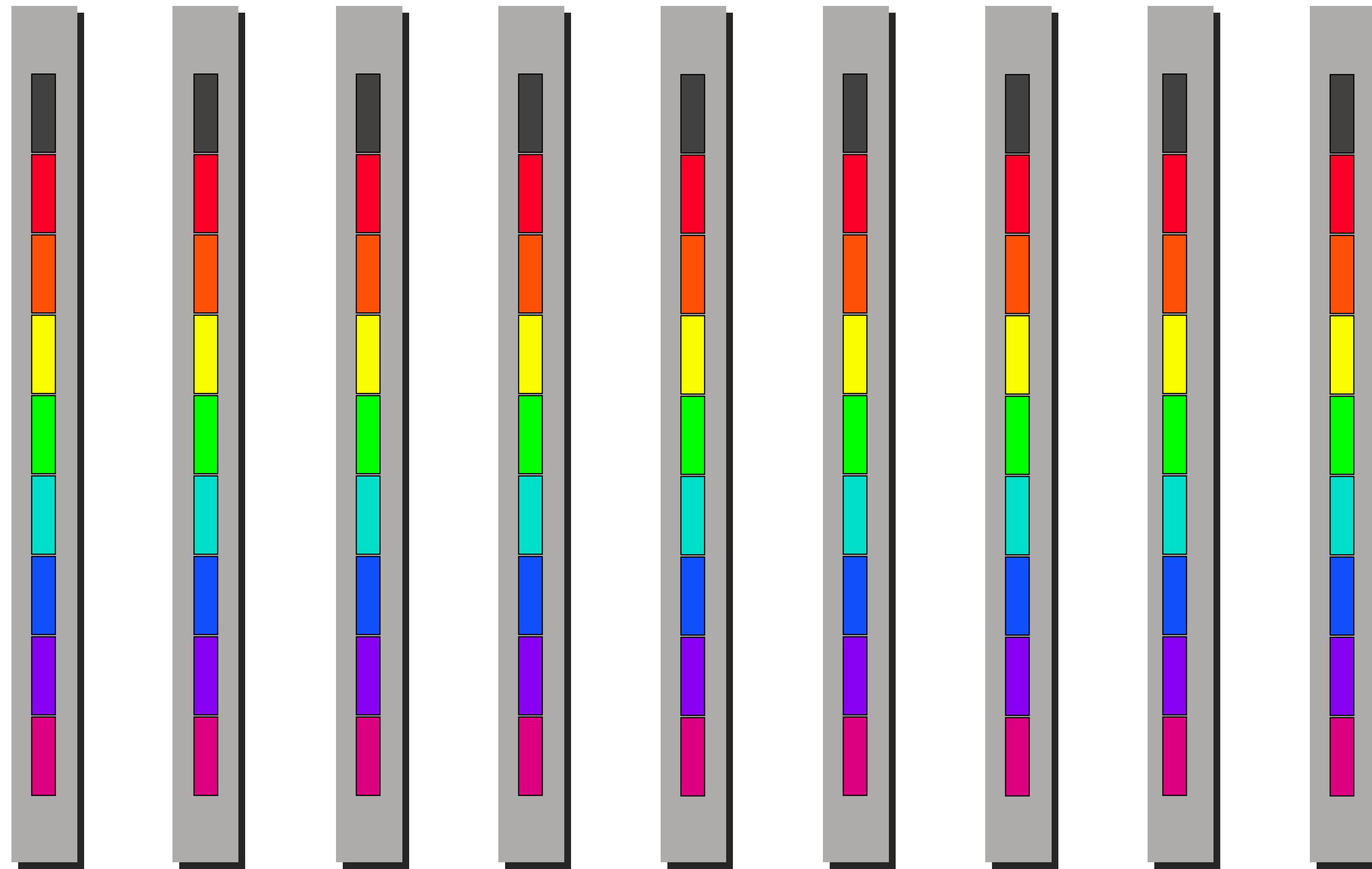






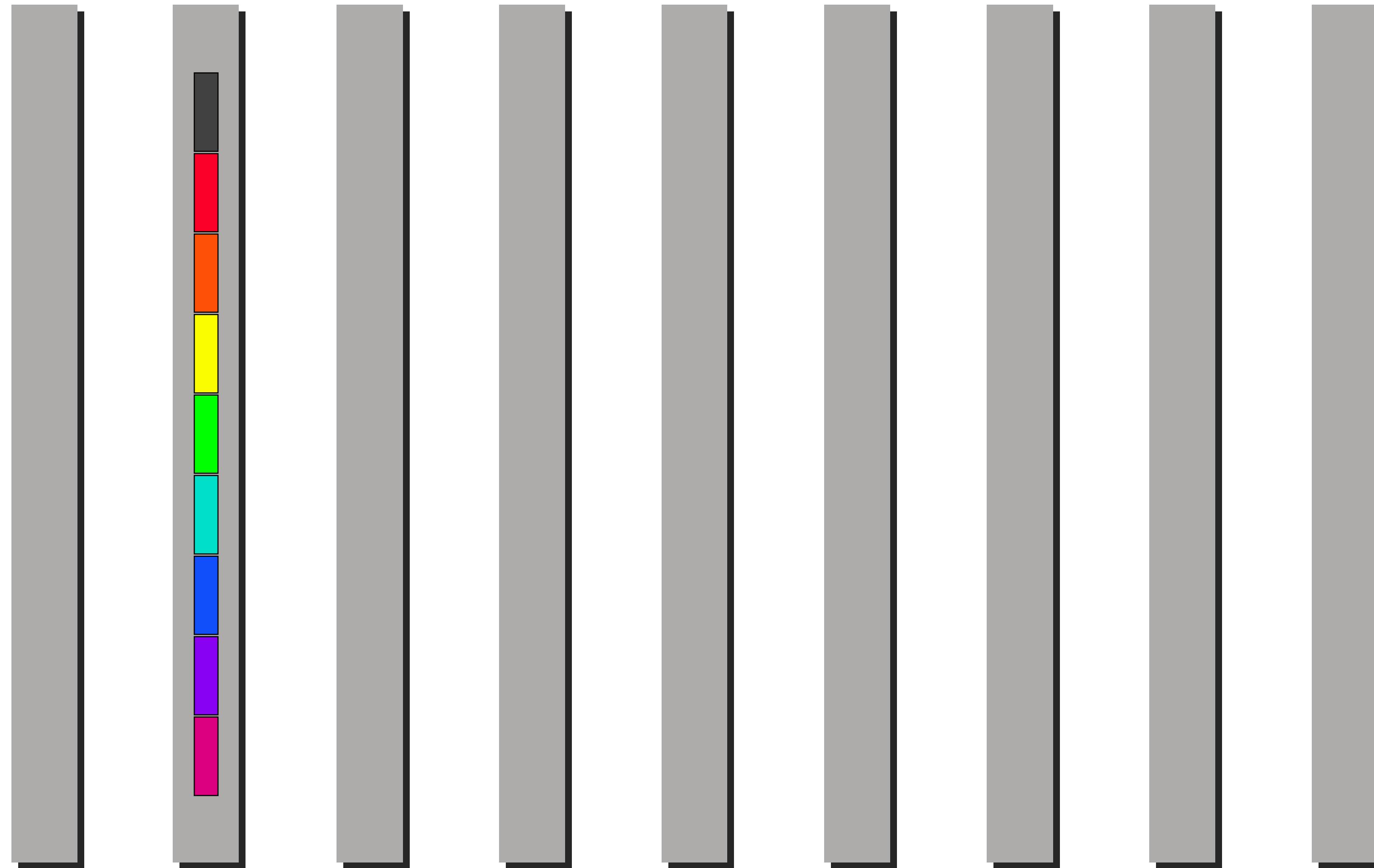


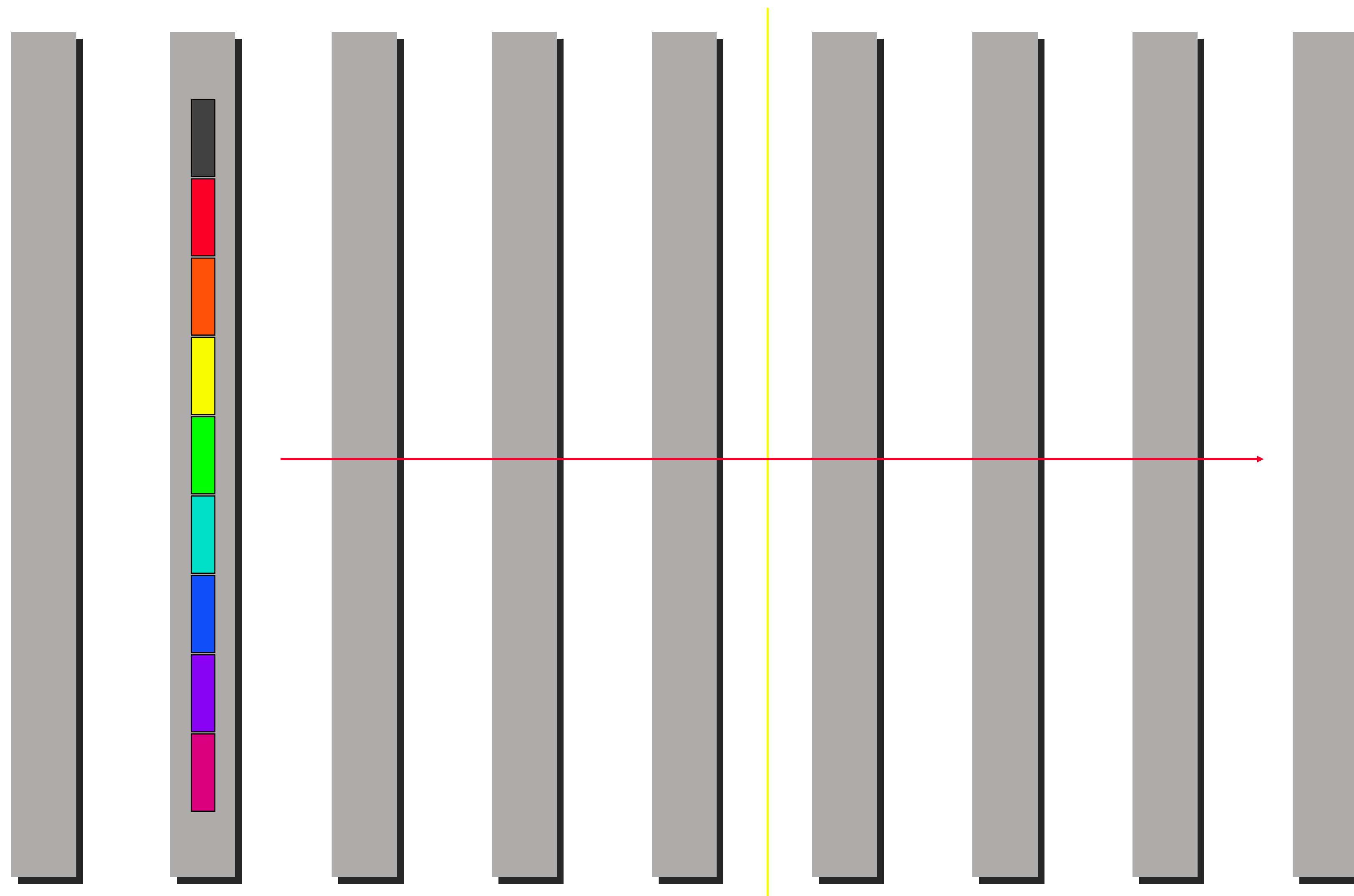


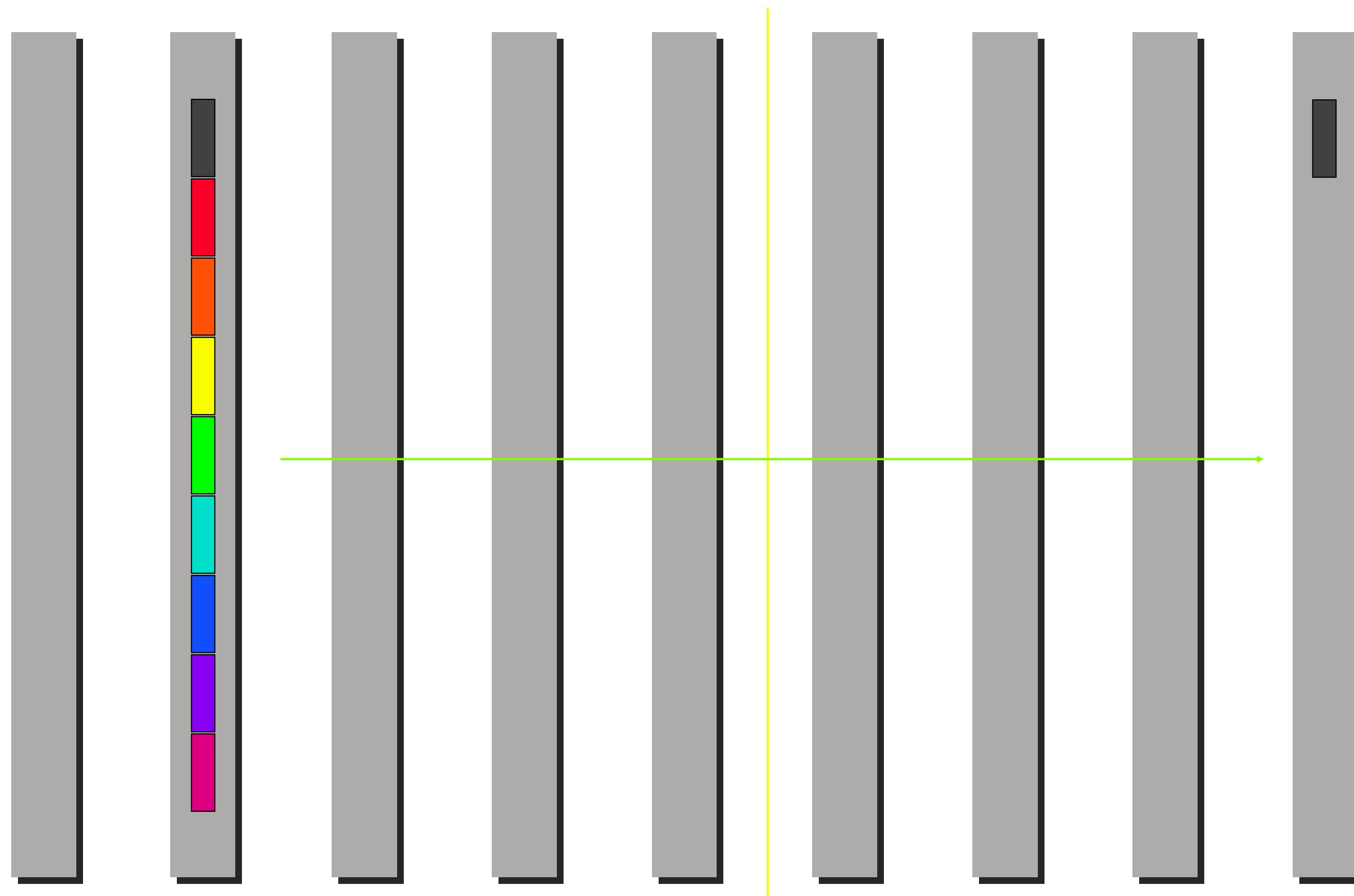


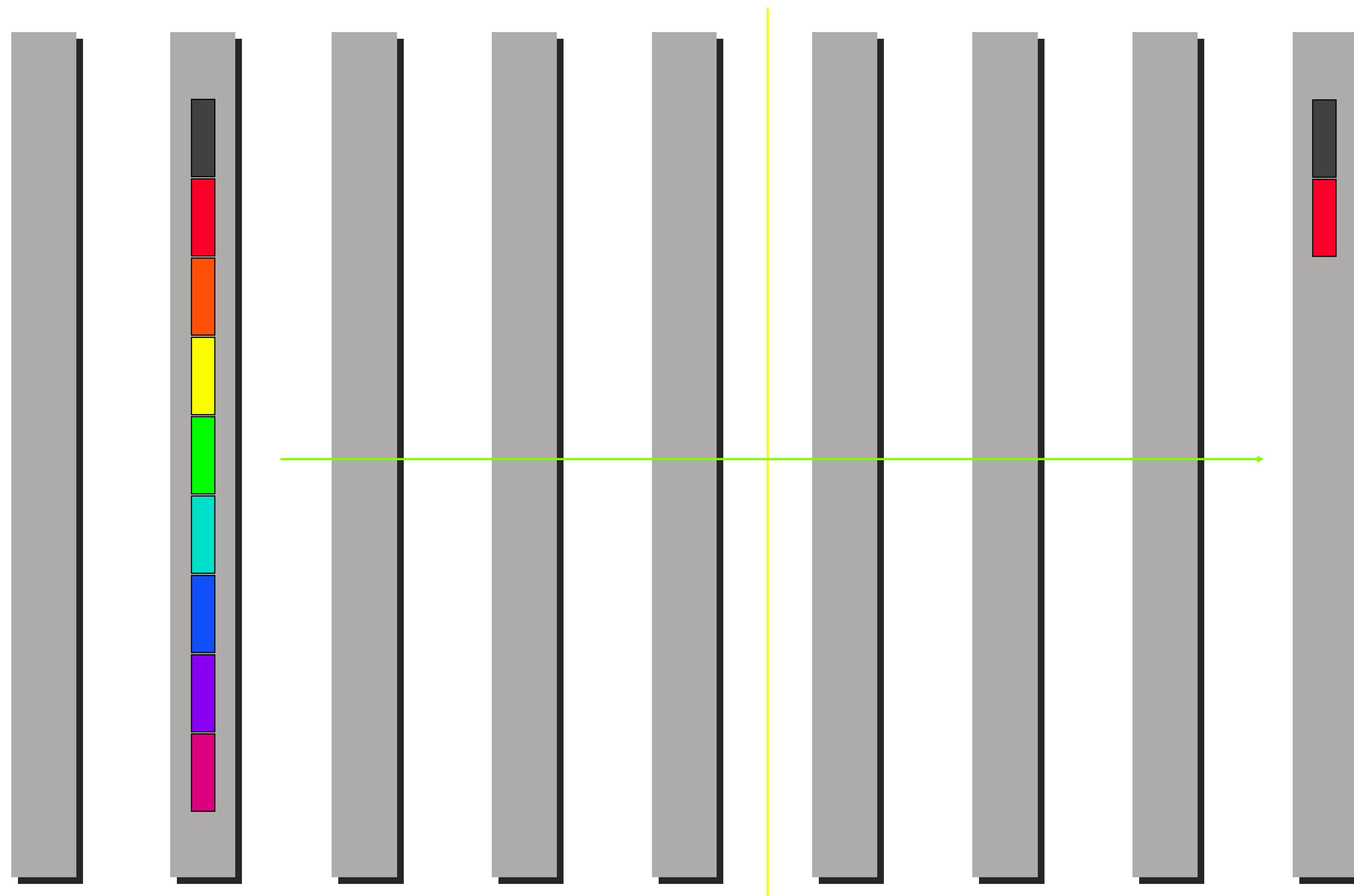
Let us view this more closely

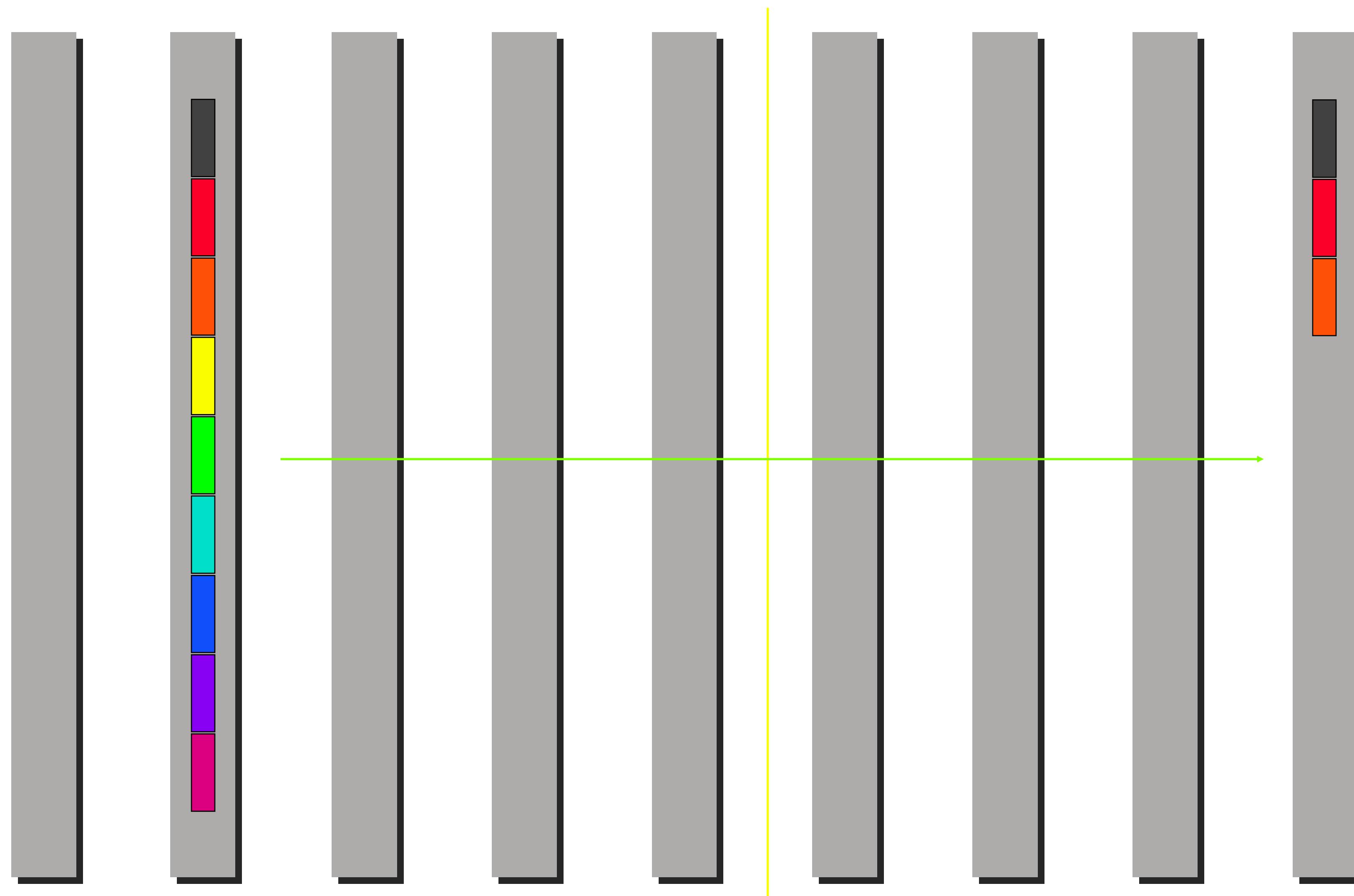
- Red arrows indicate startup of communication (leading to latency, α)
- Green arrows indicate packets in transit (leading to a bandwidth related cost proportional to β and the length of the packet)

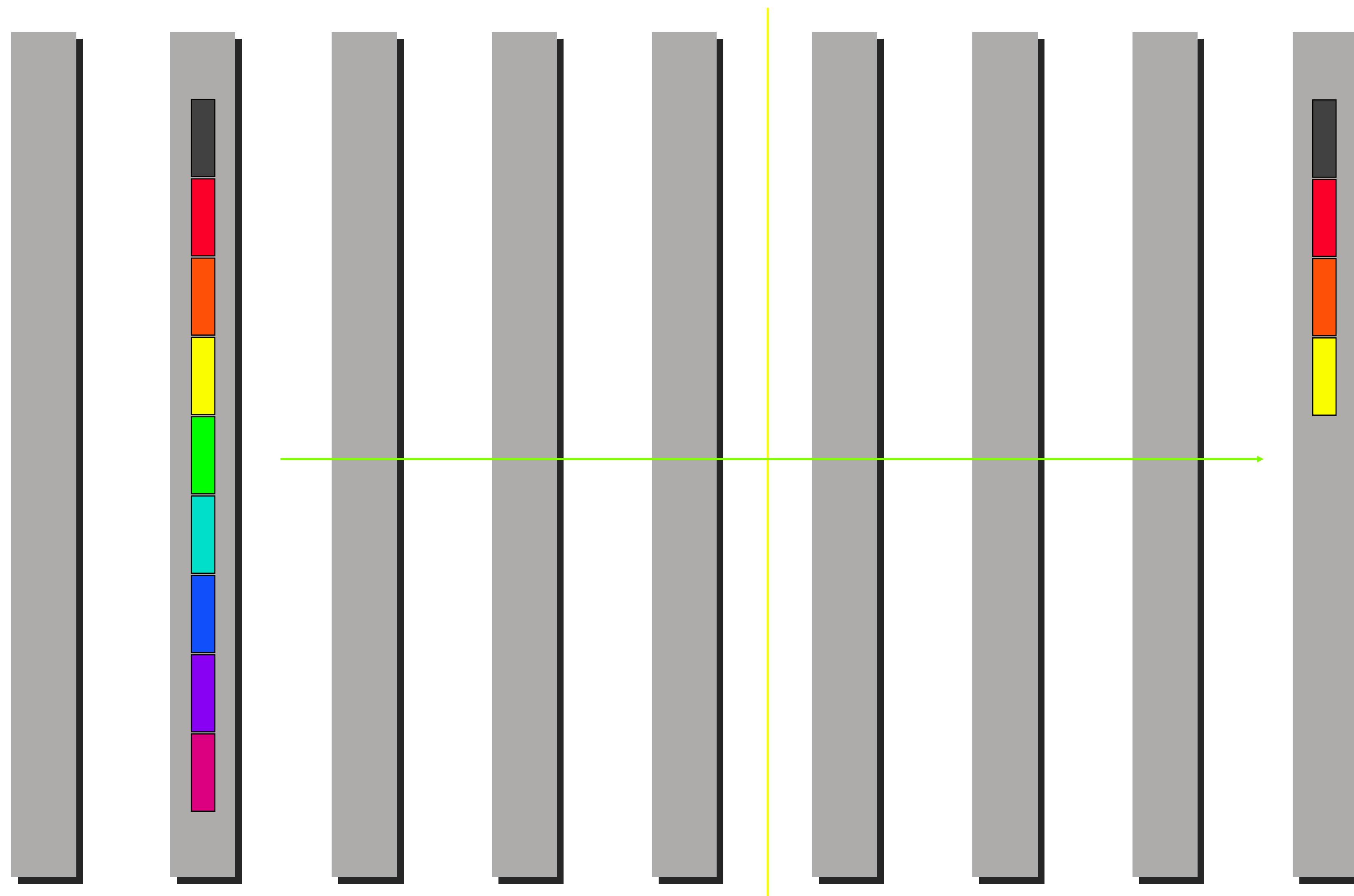


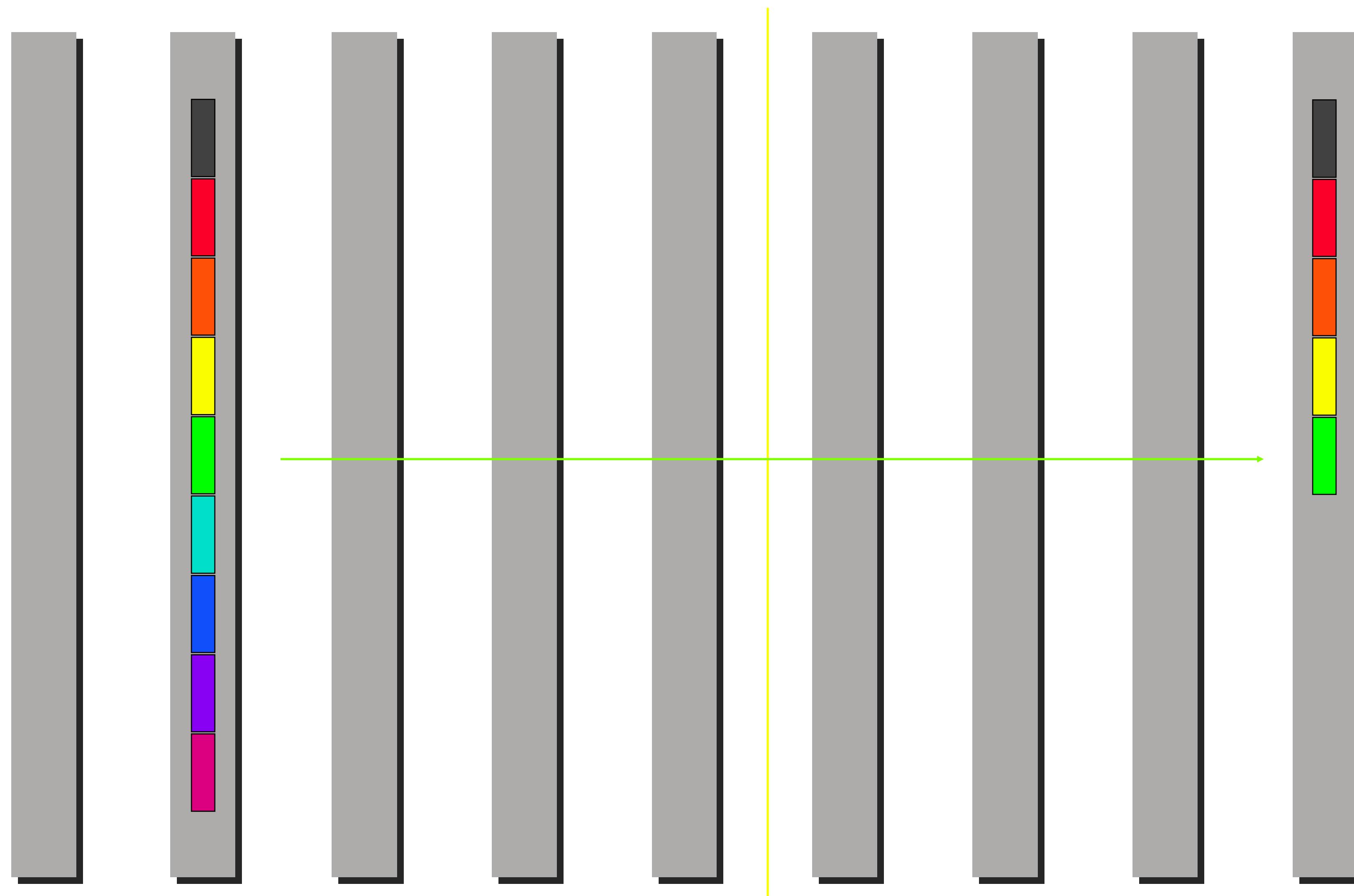


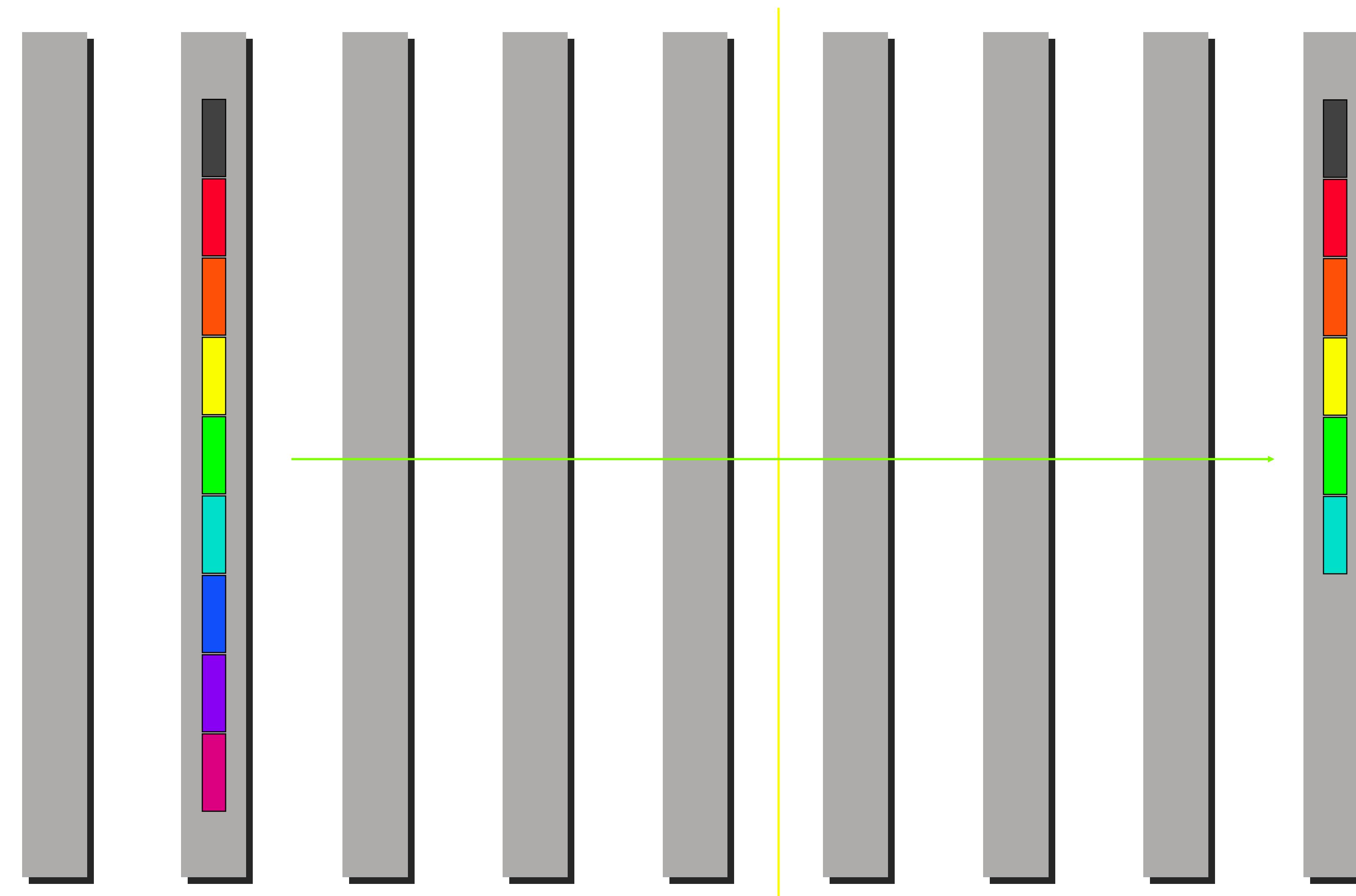


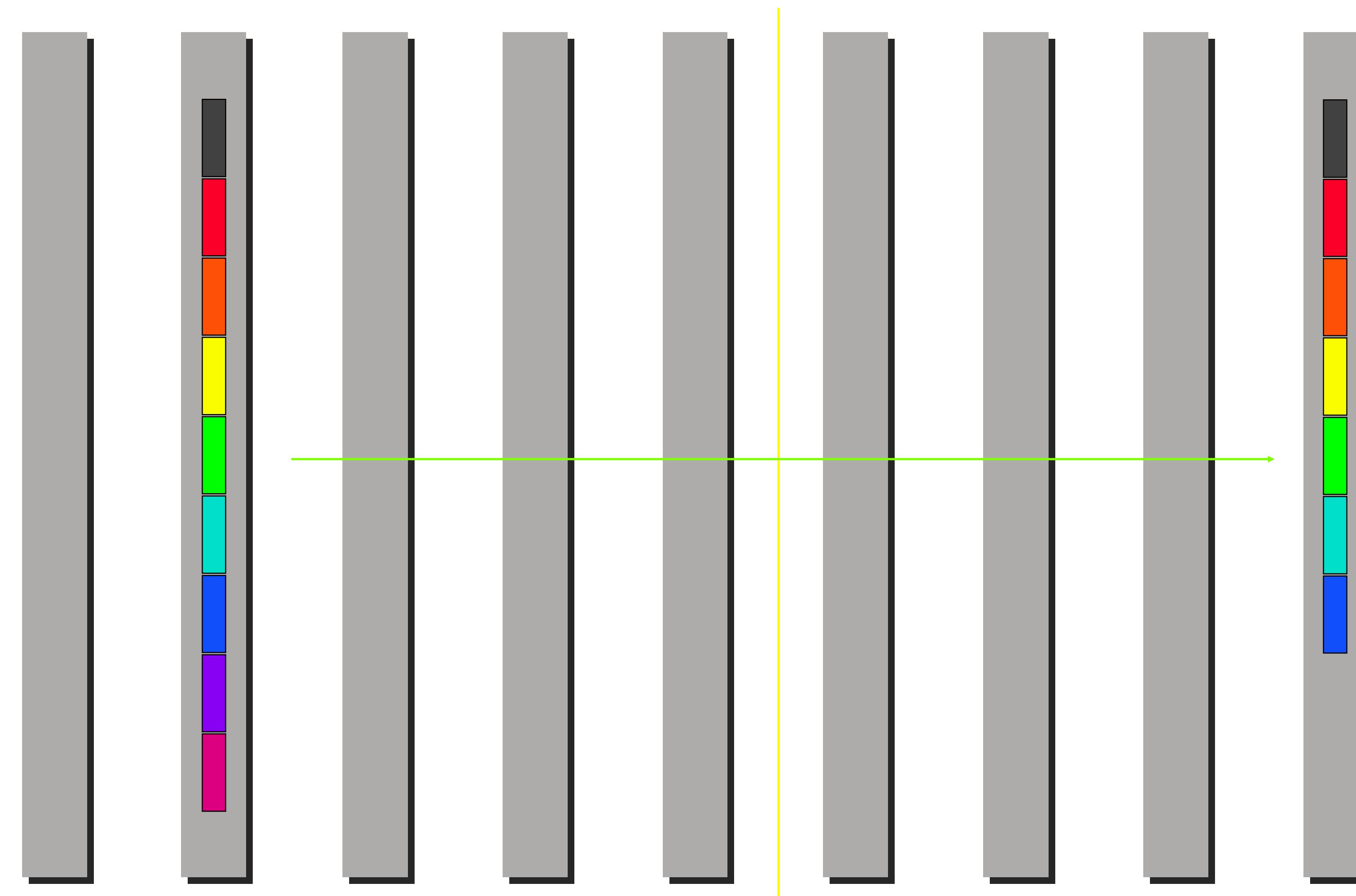


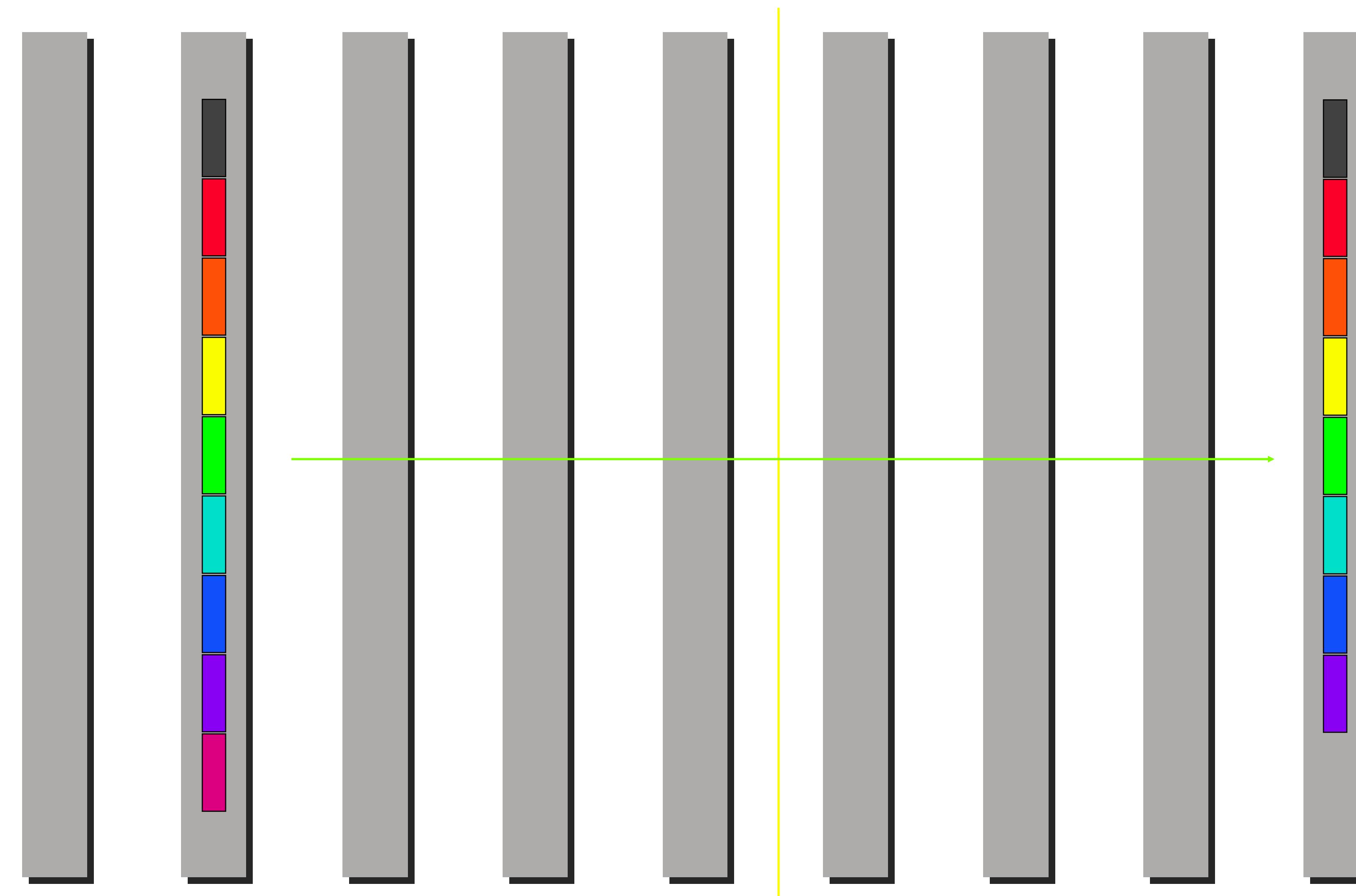


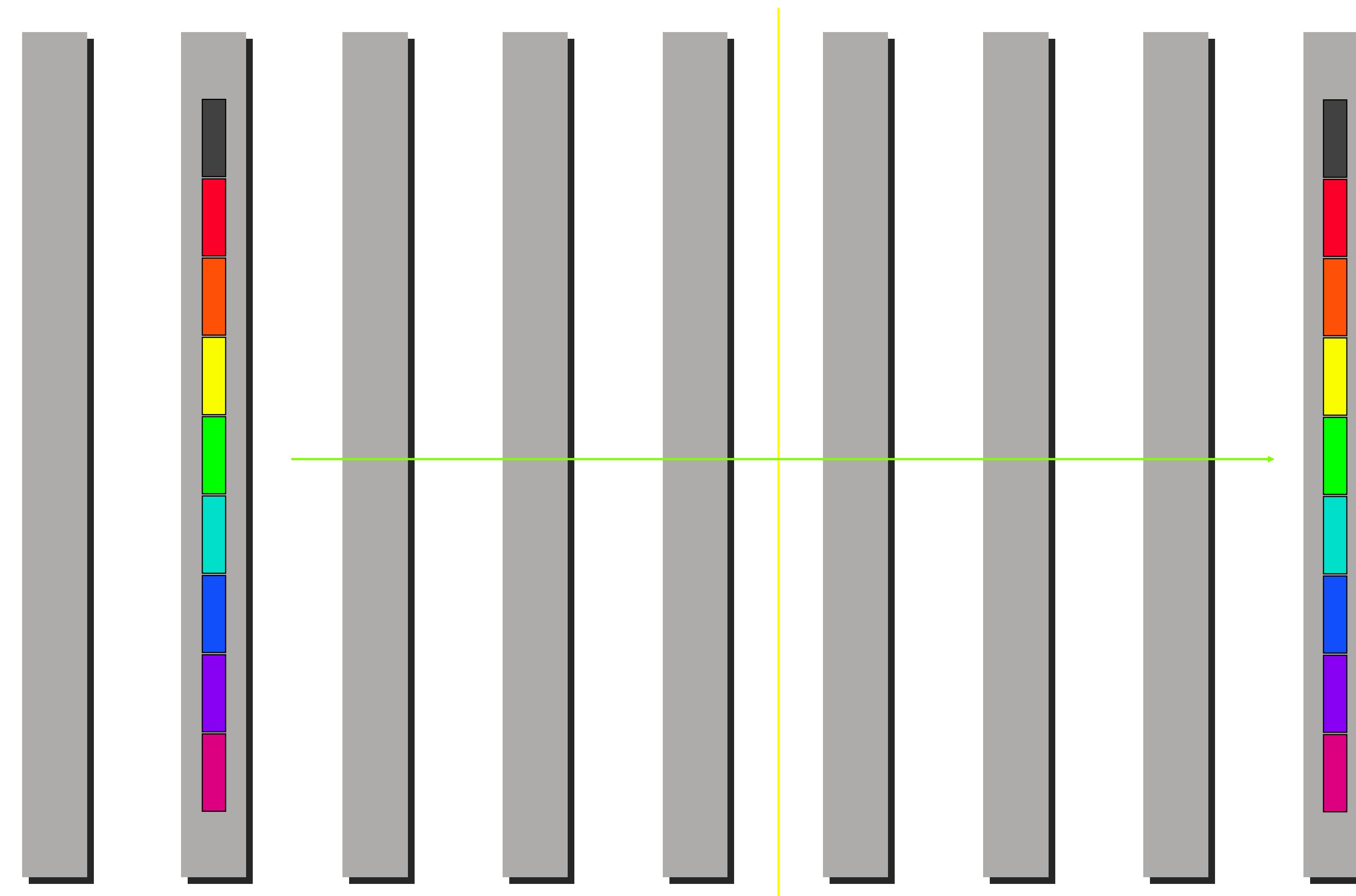


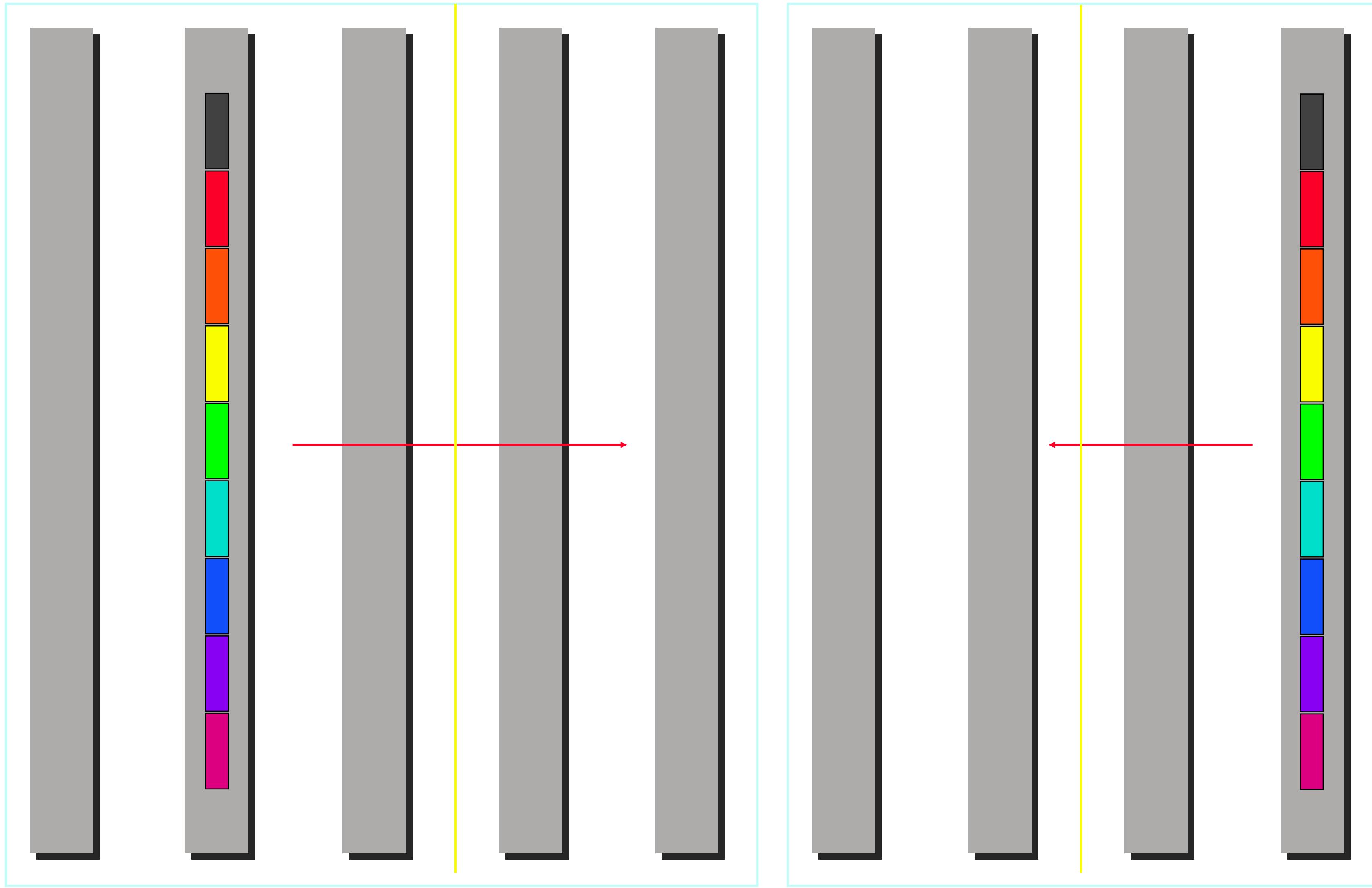


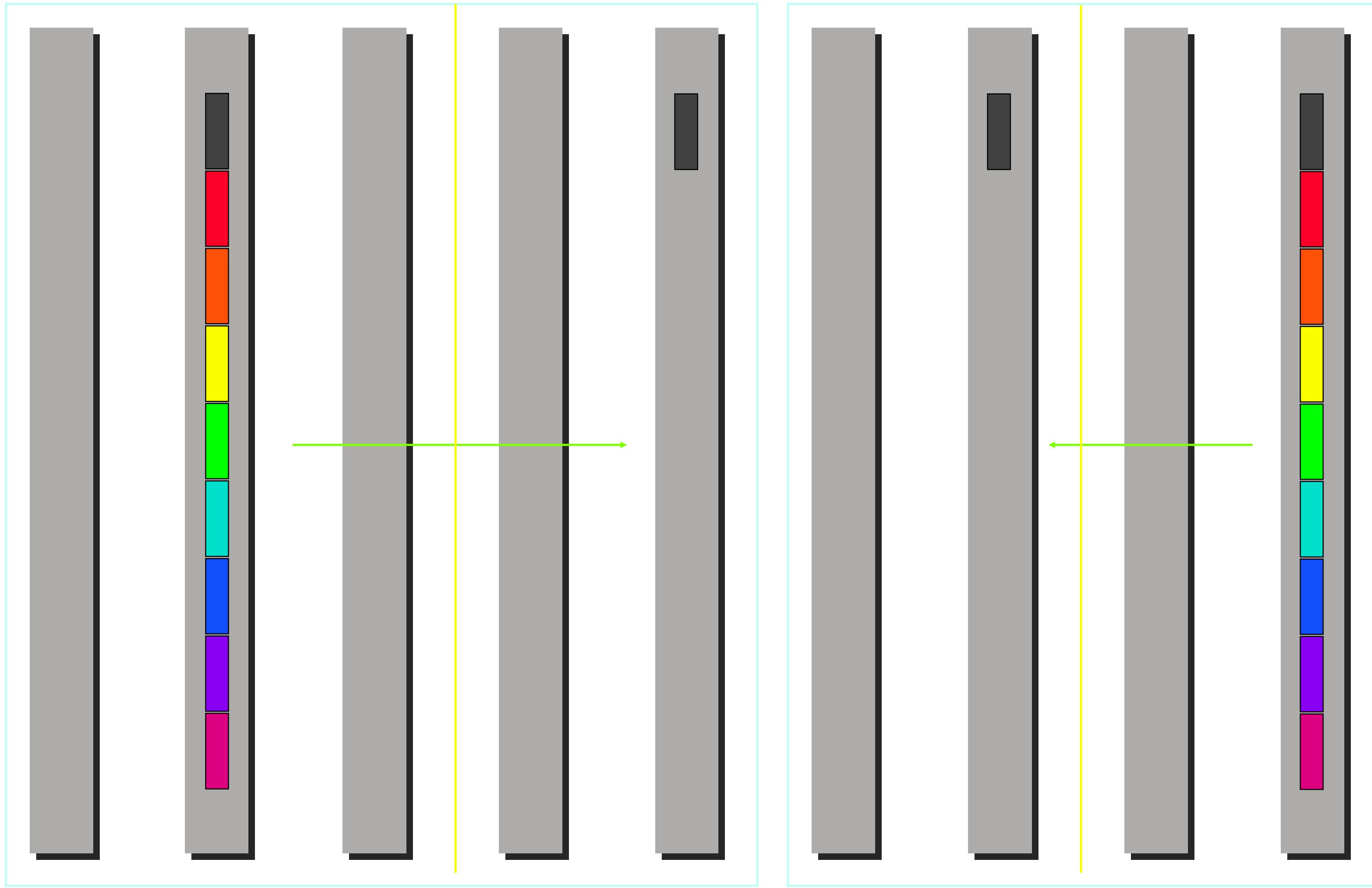


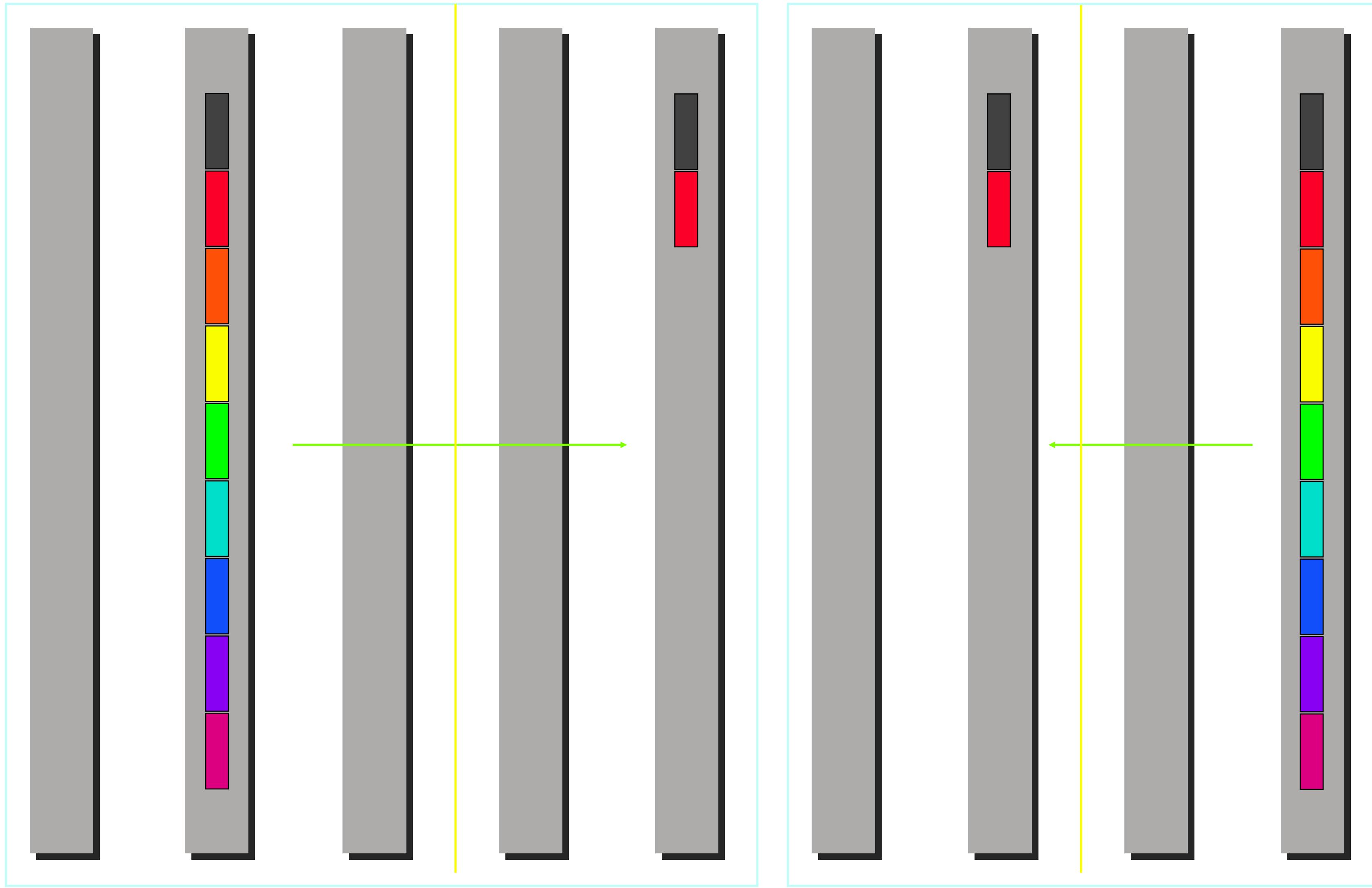


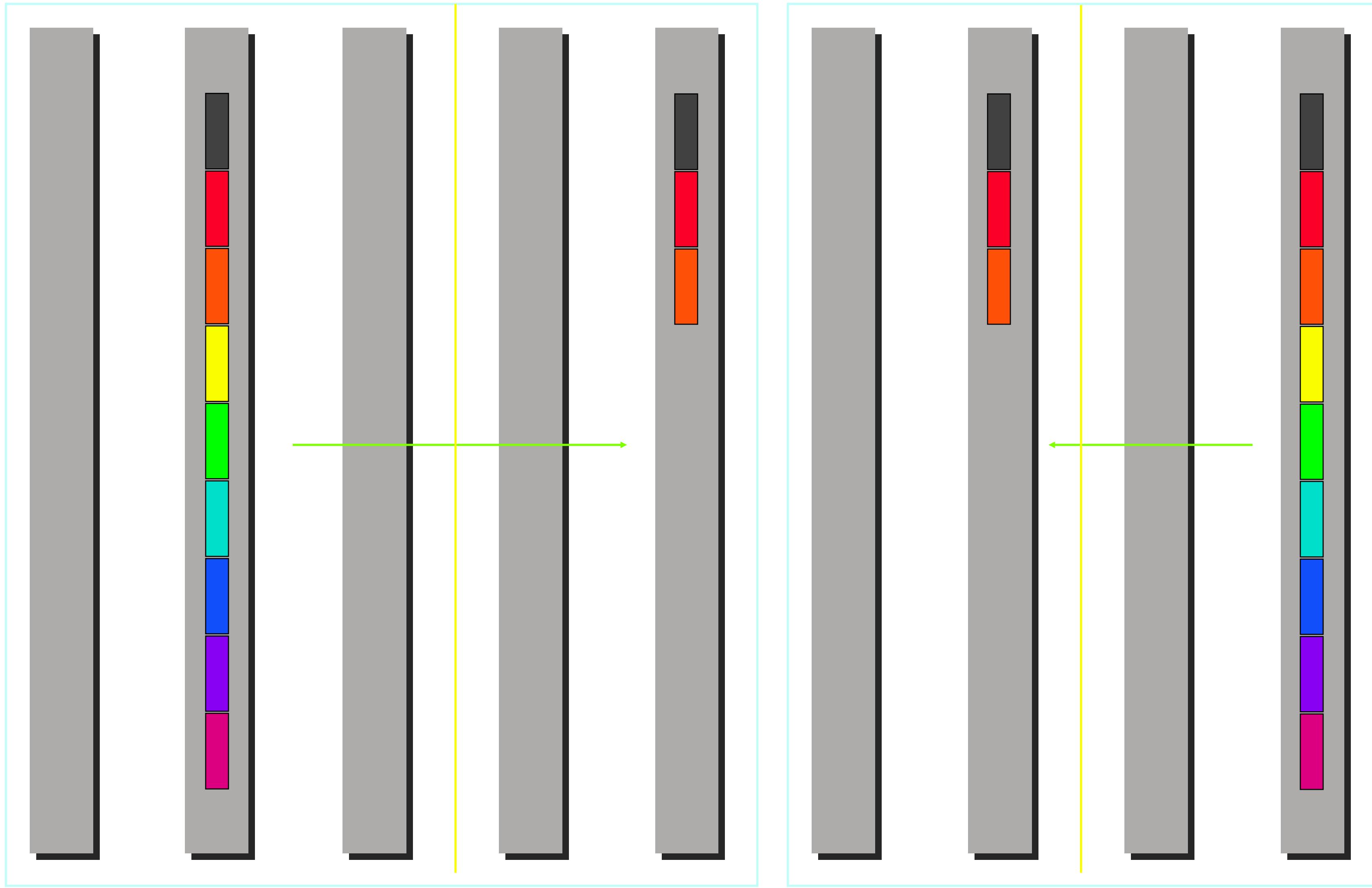


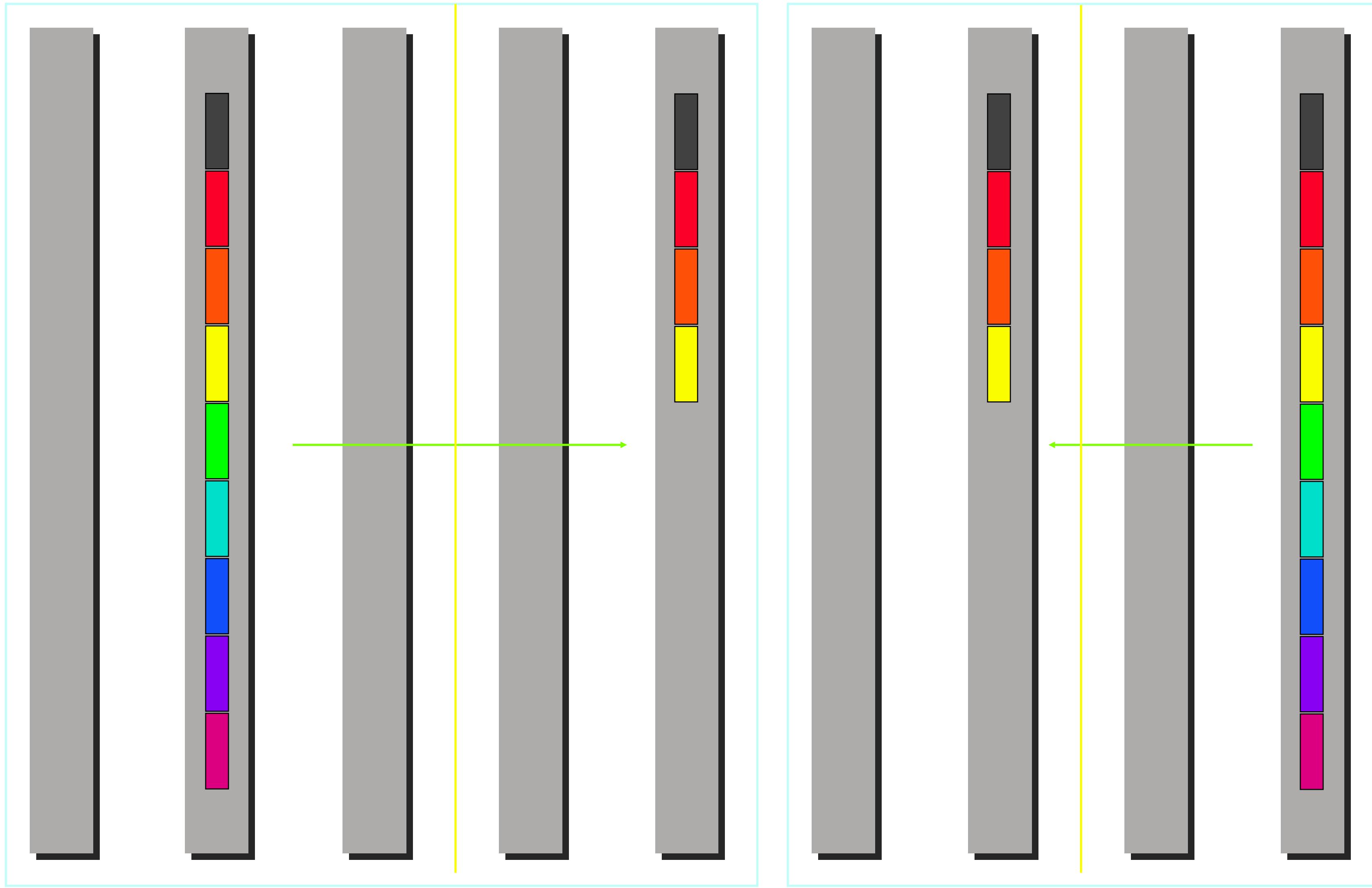


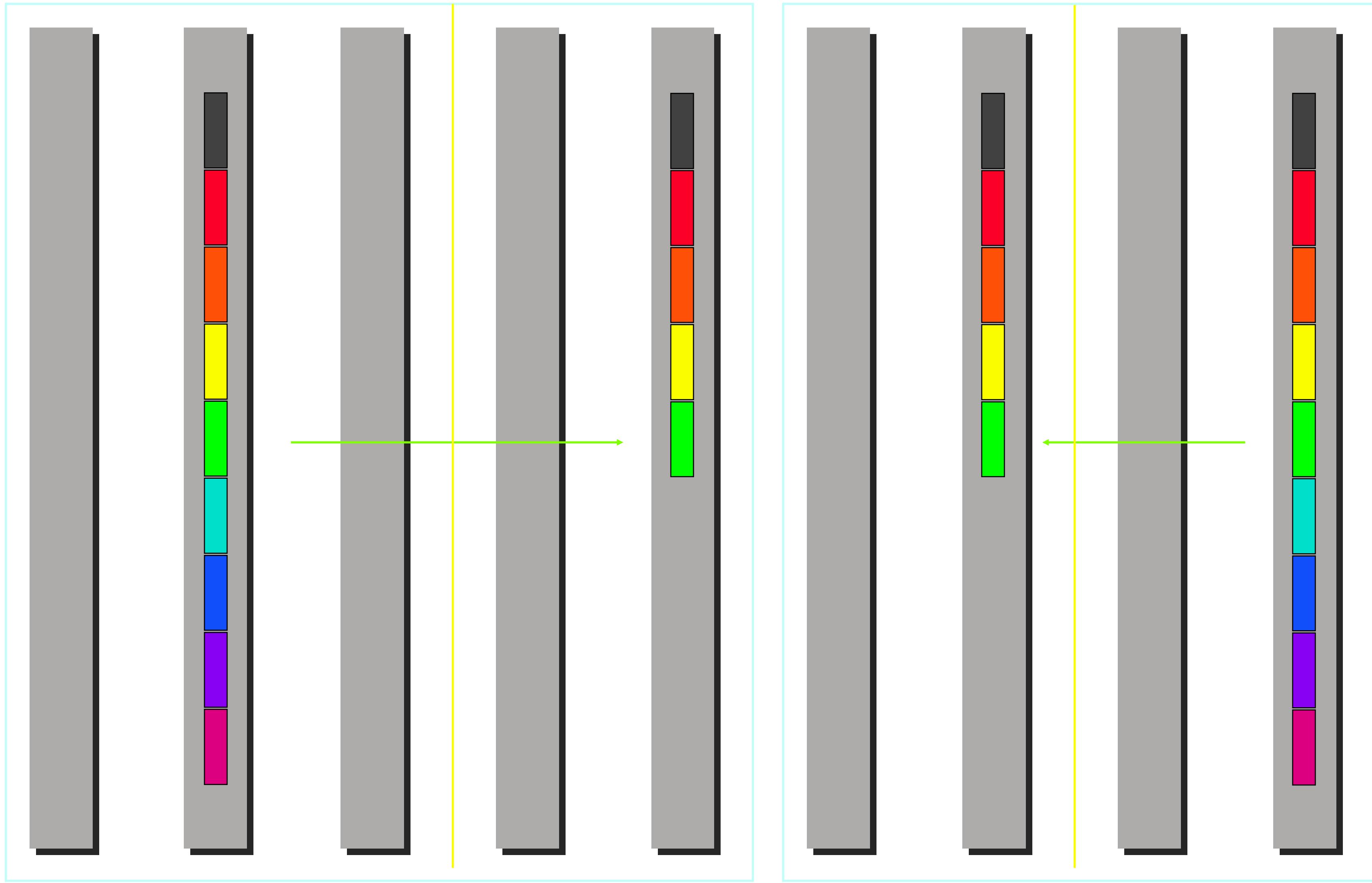


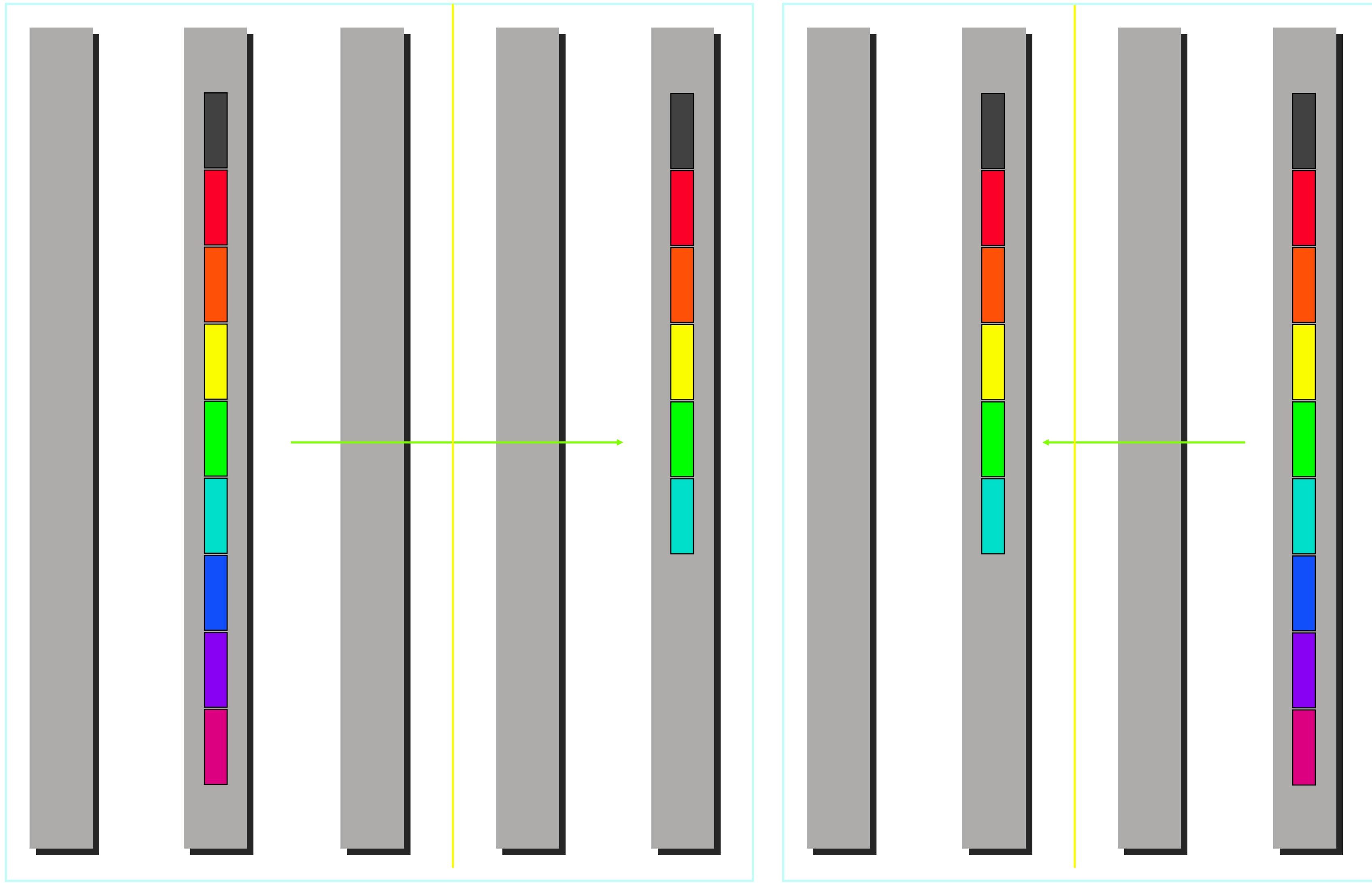


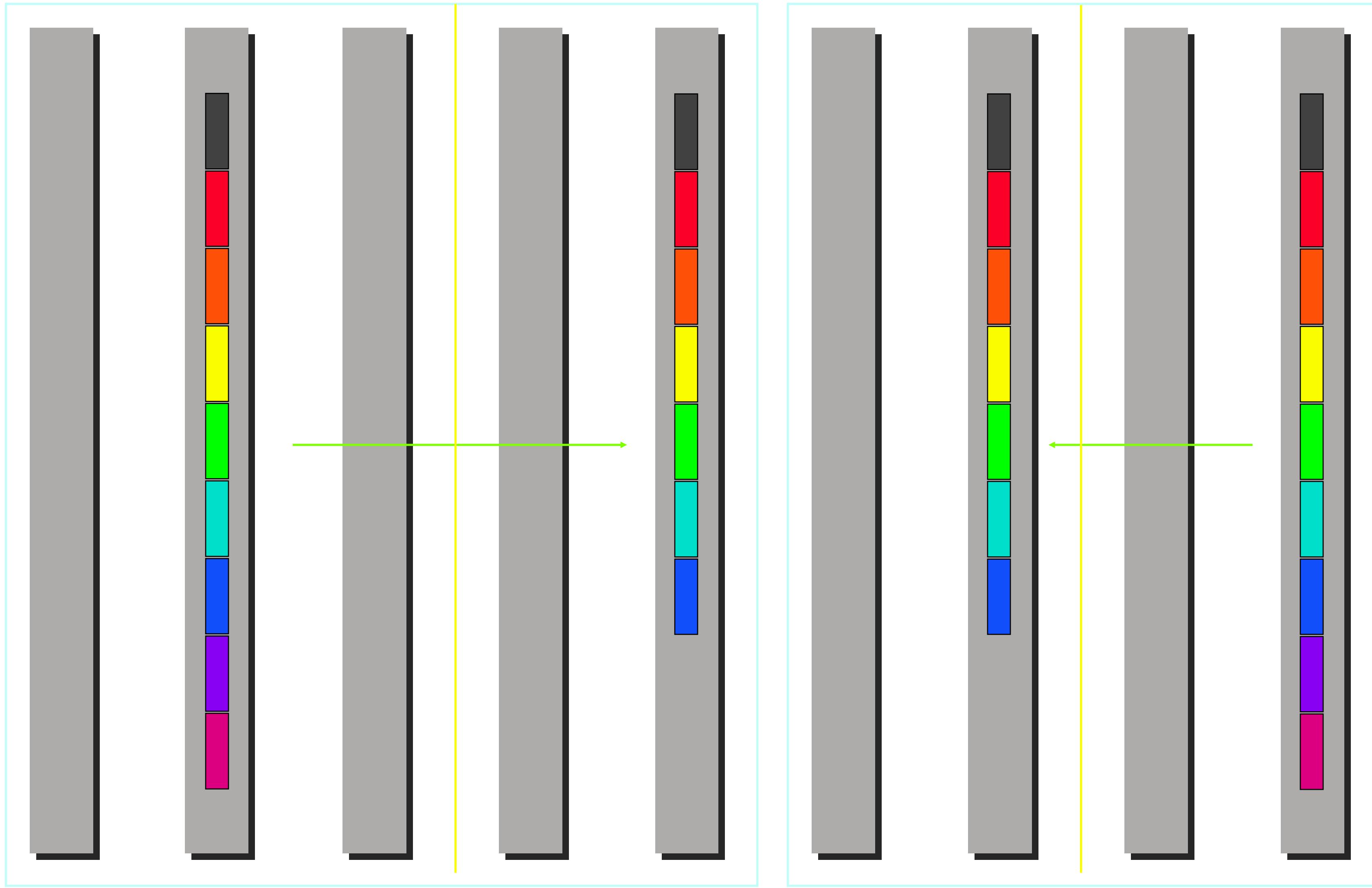


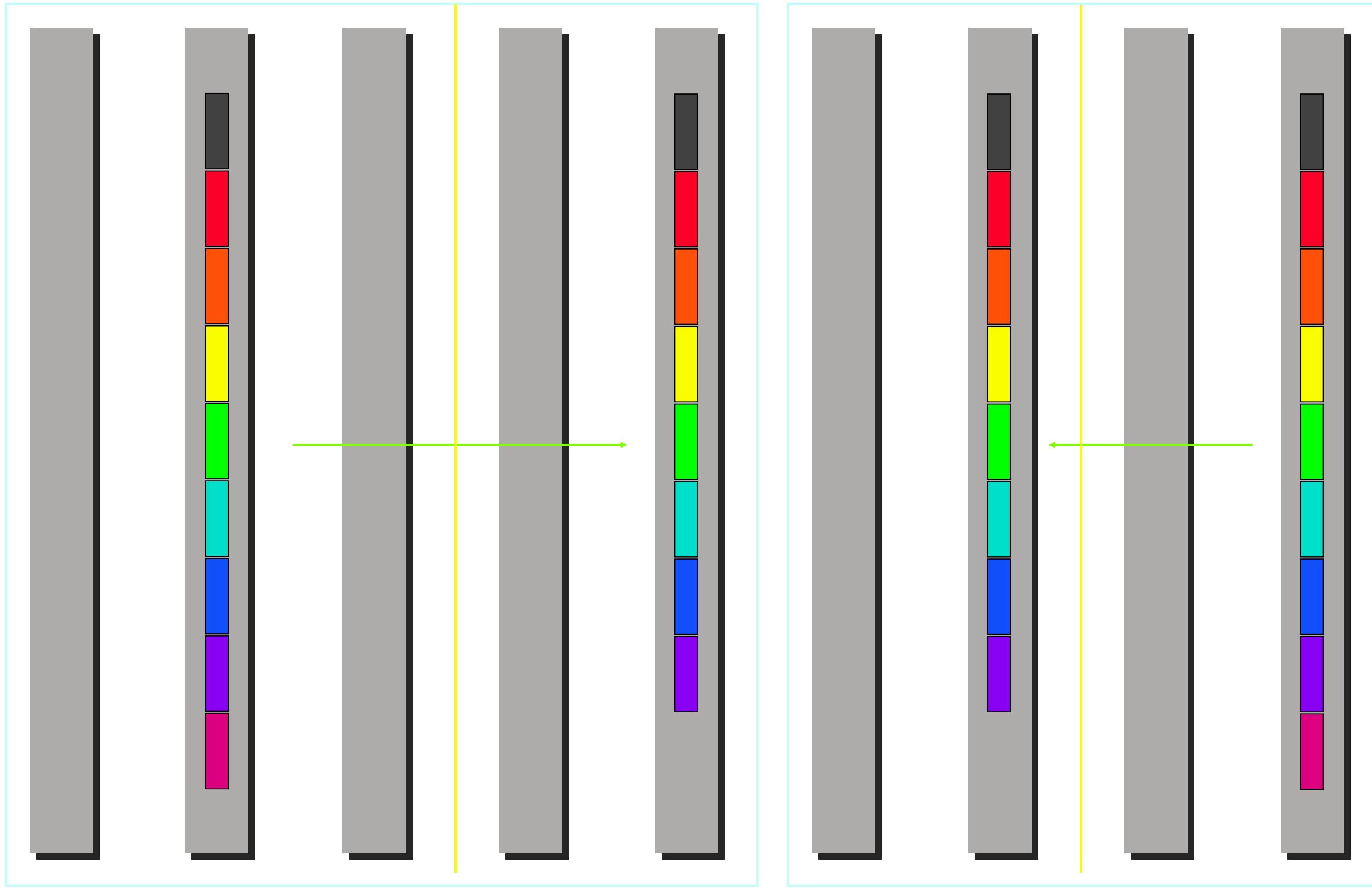


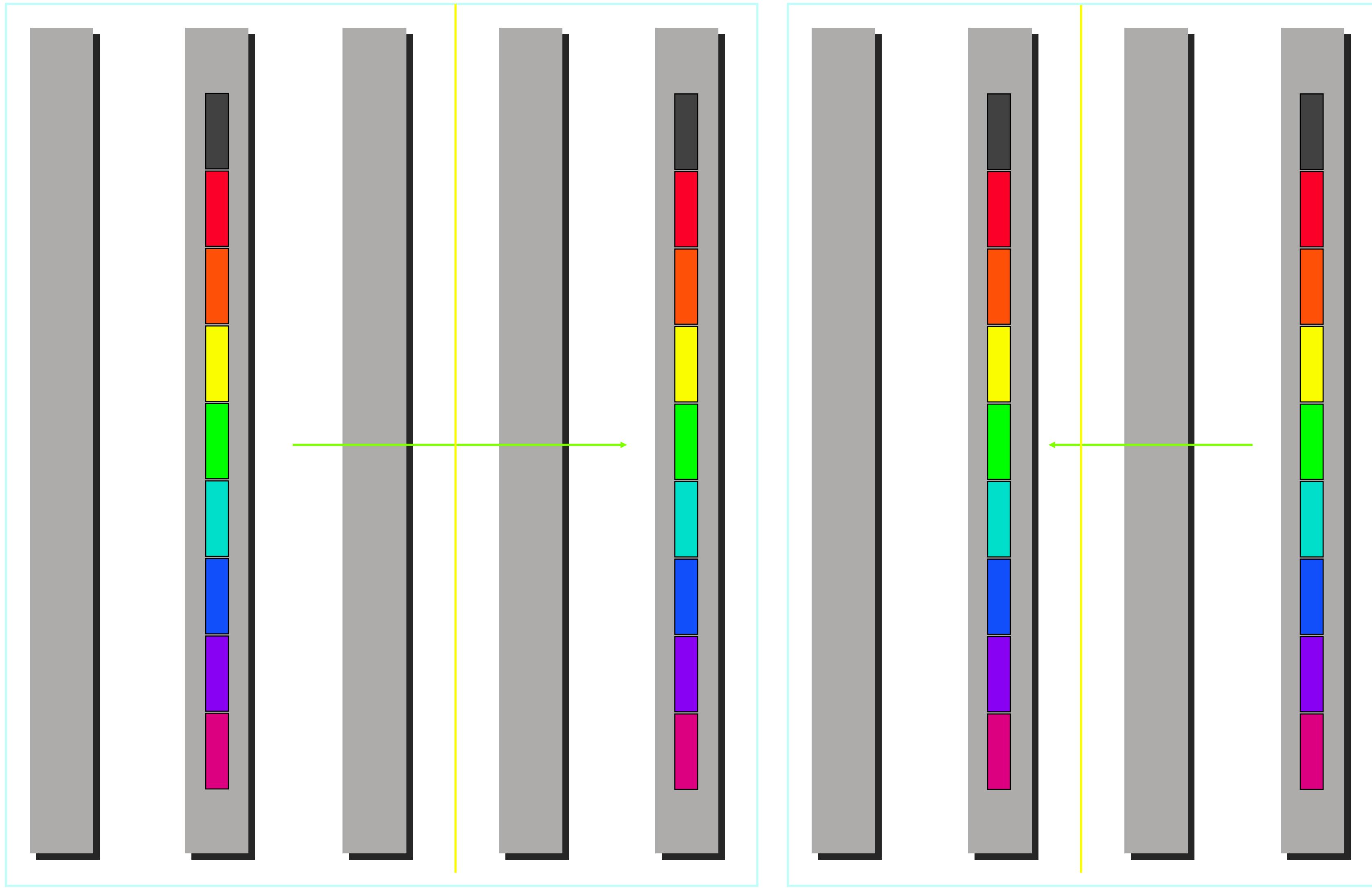


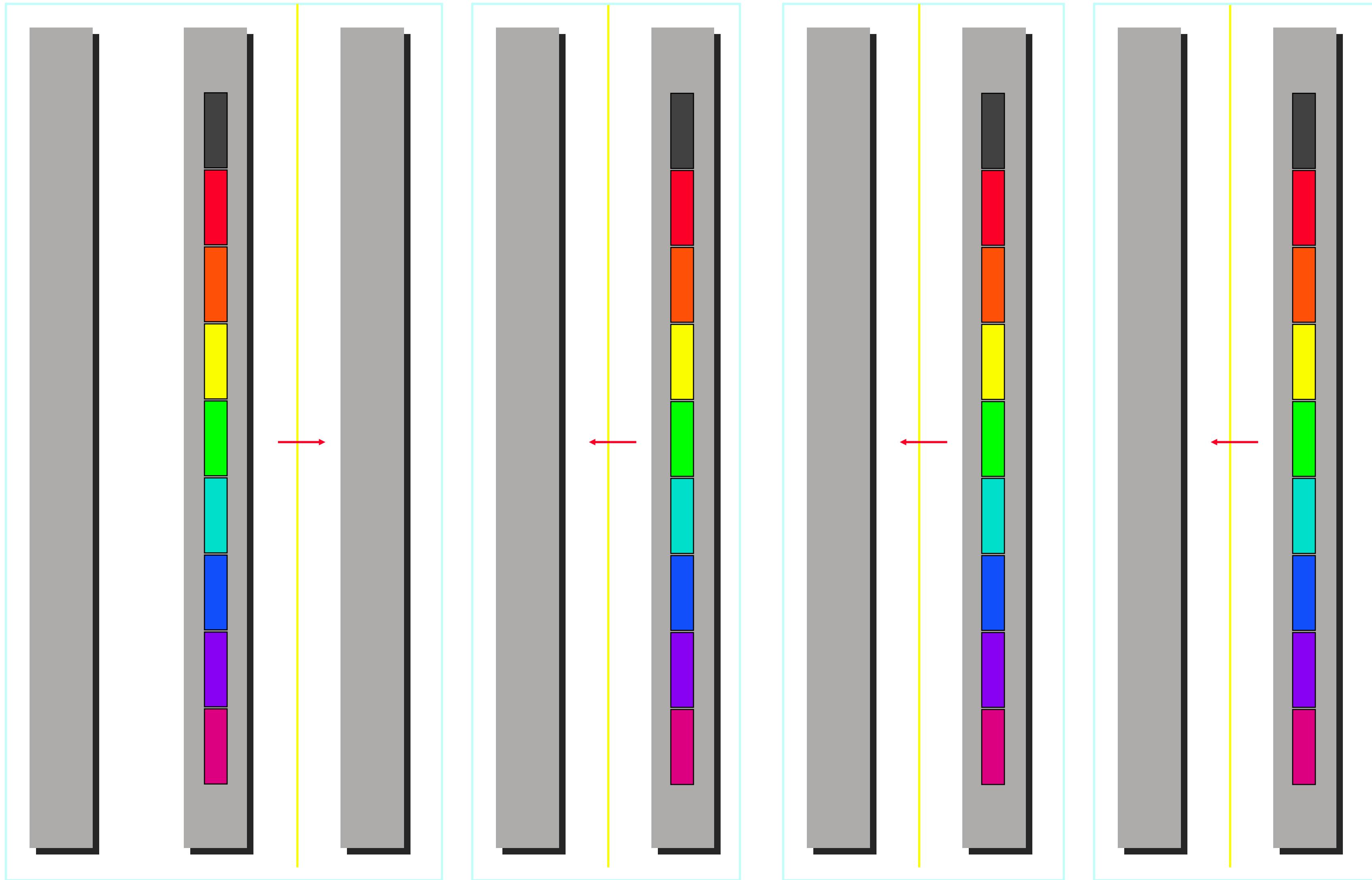


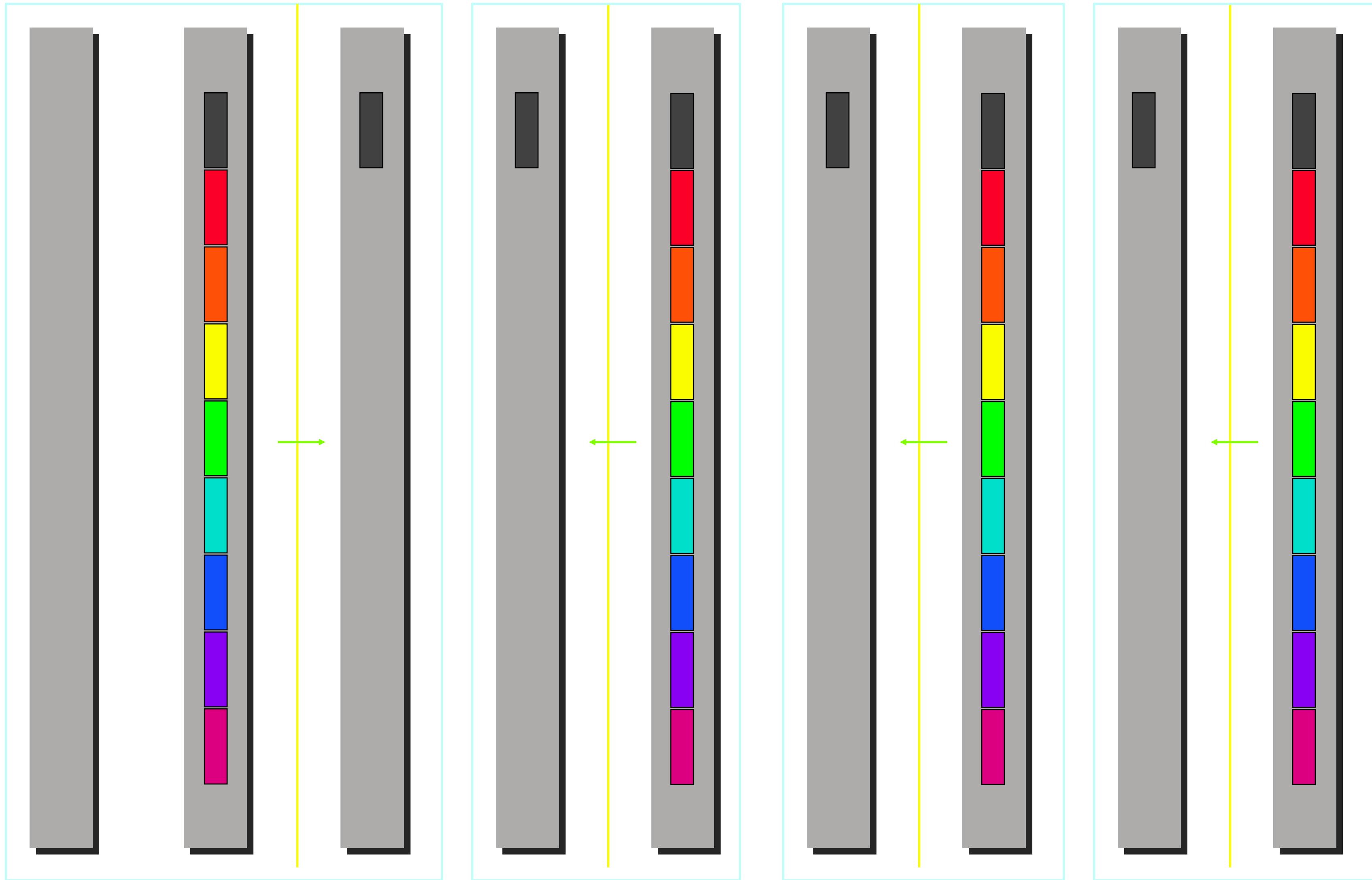


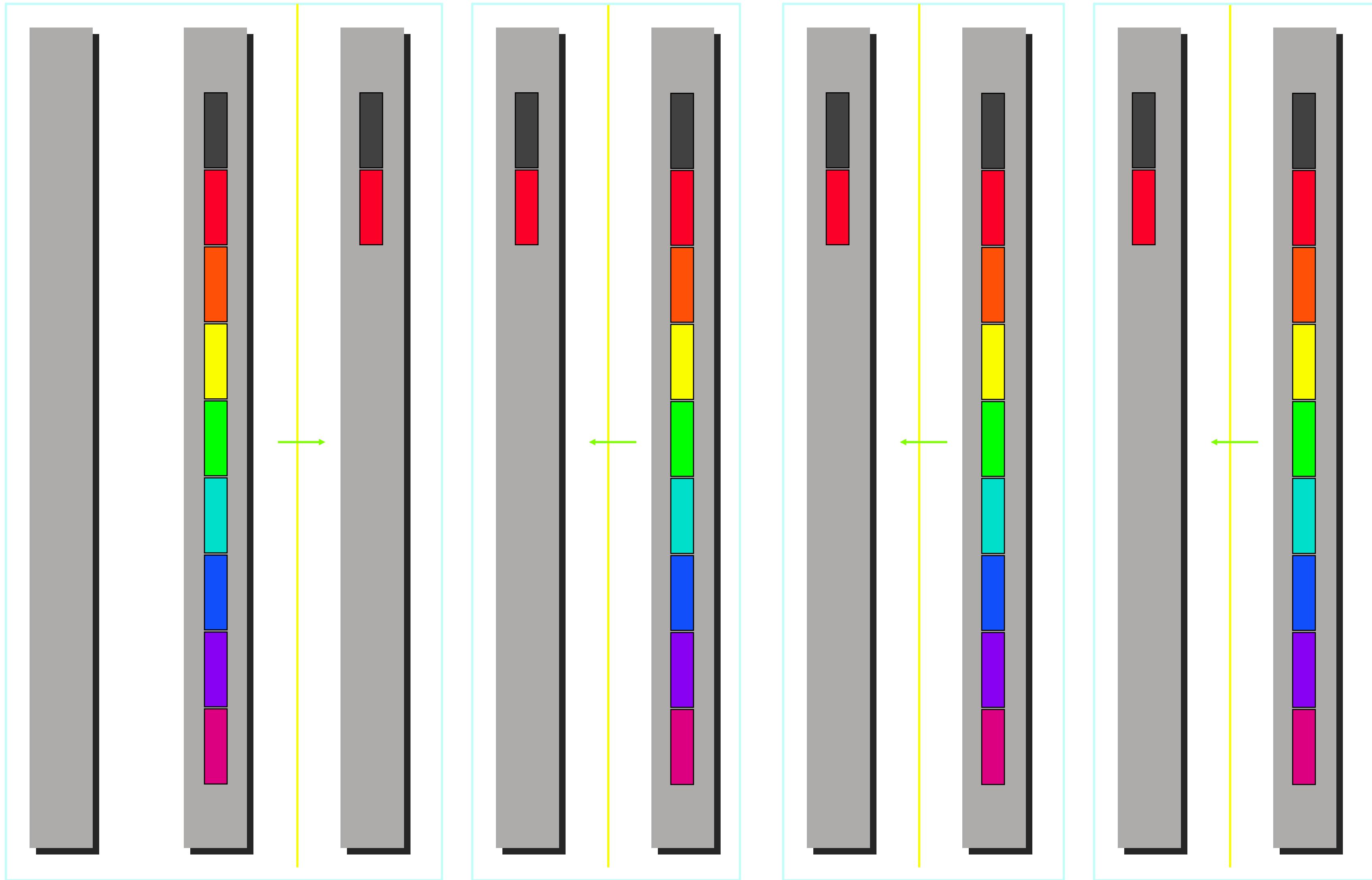


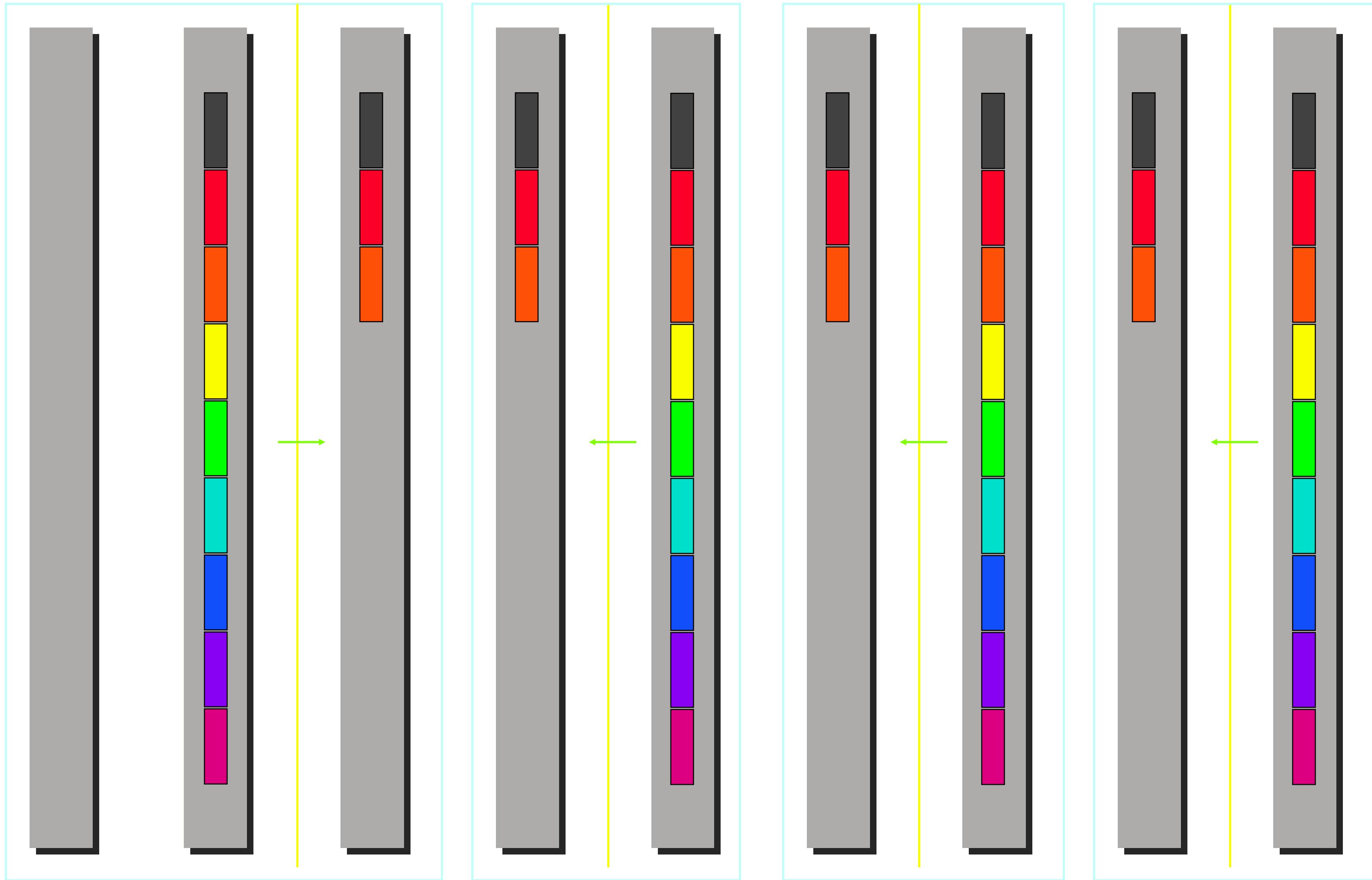


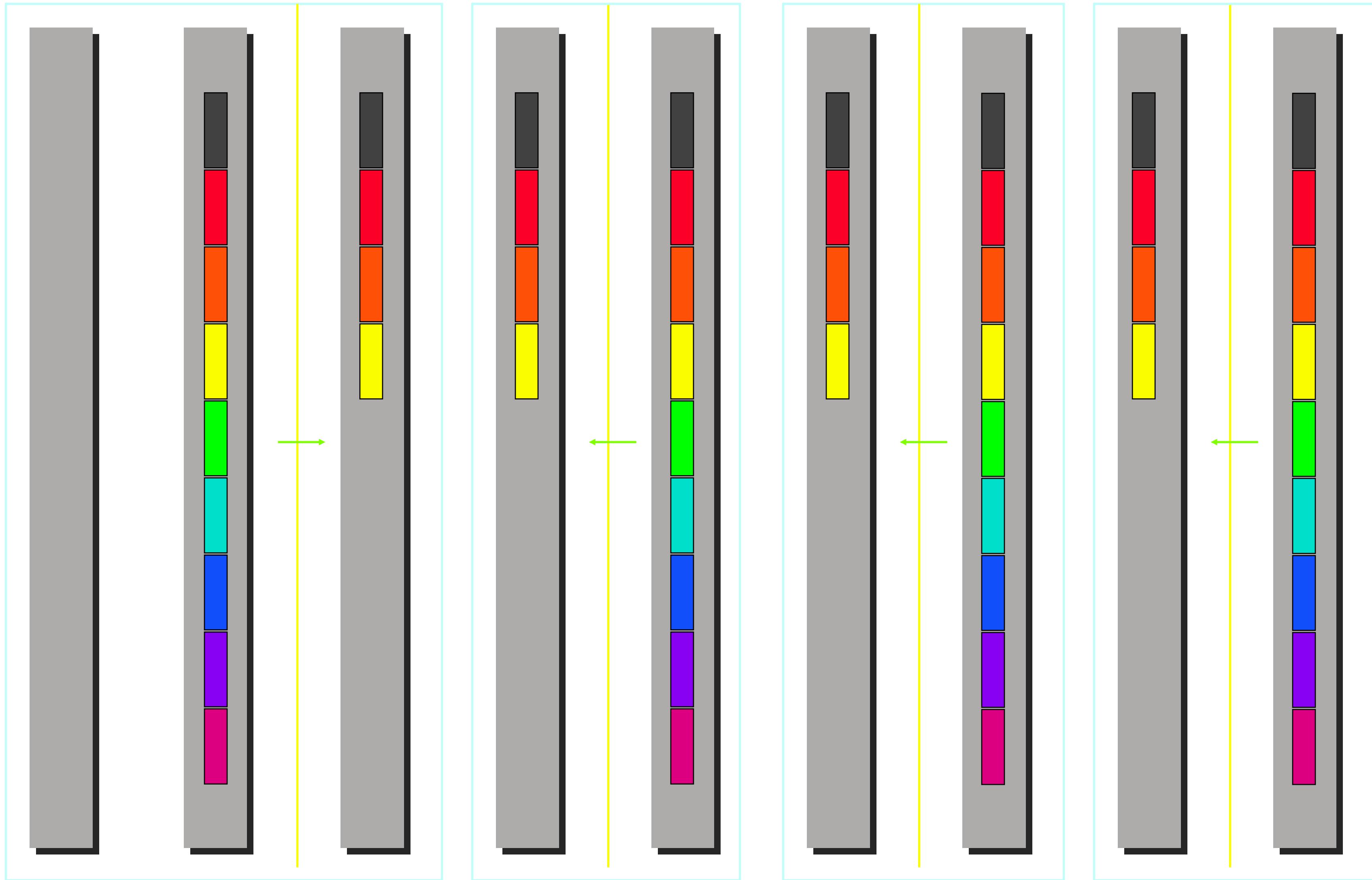


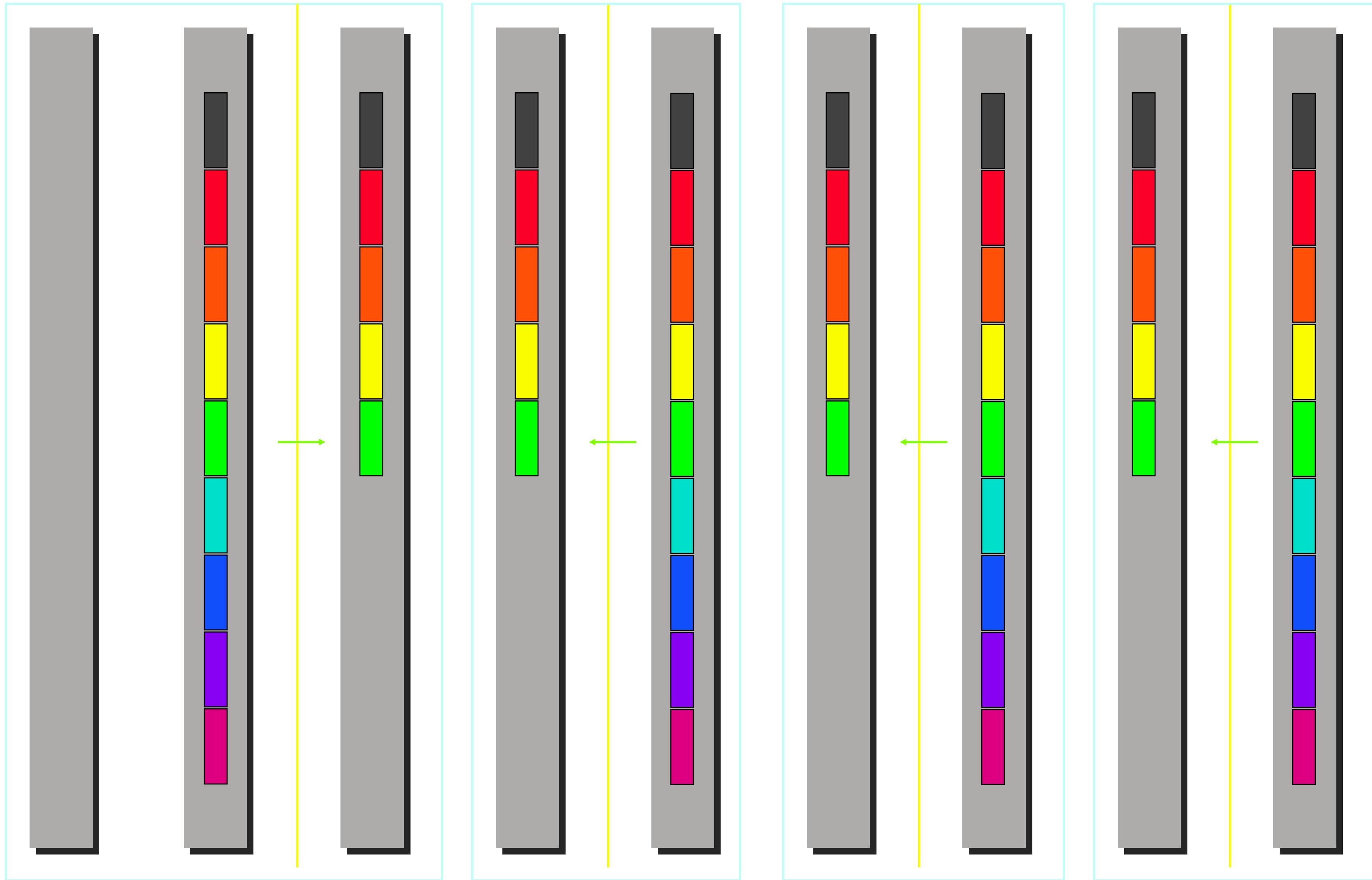


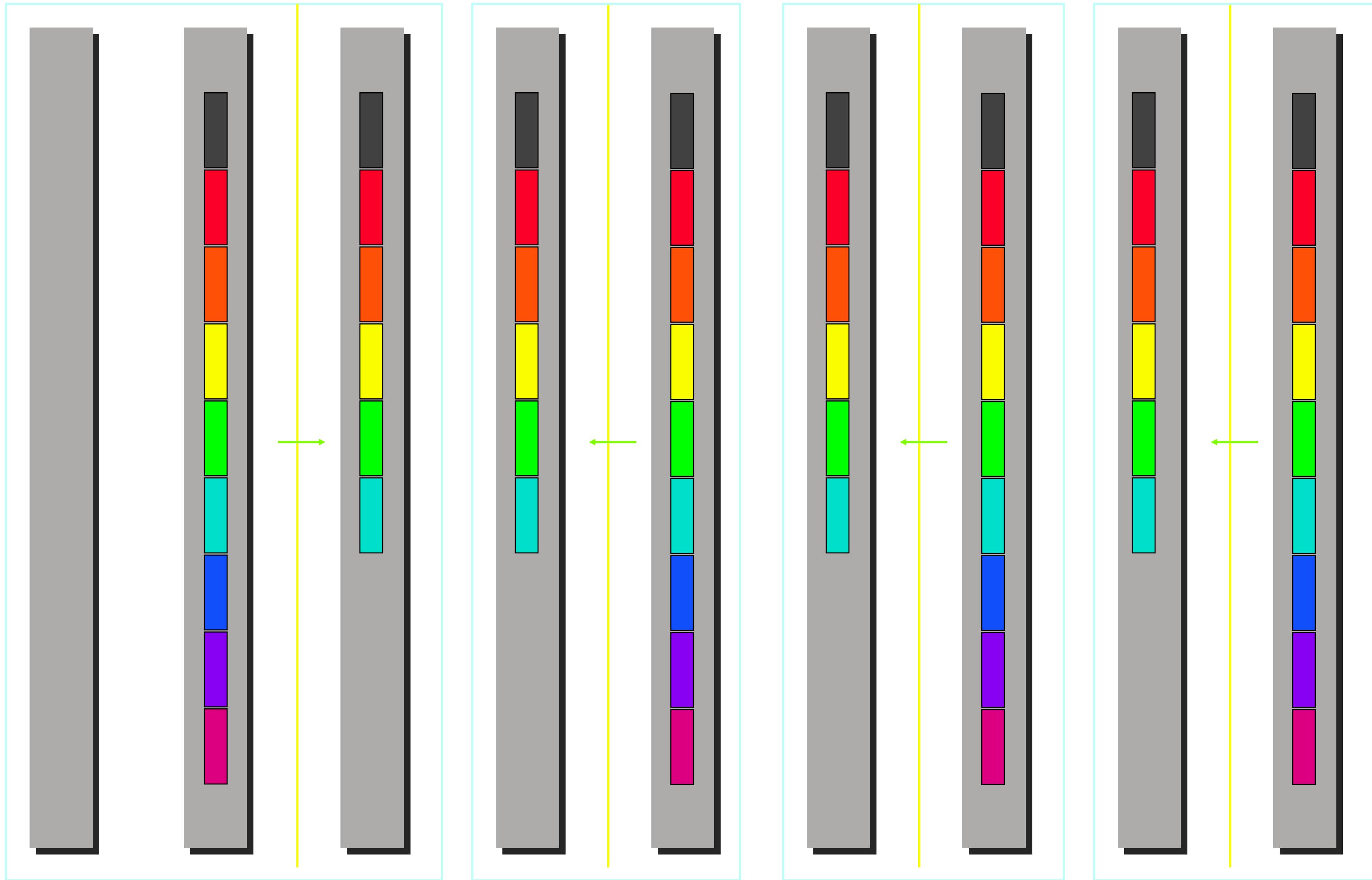


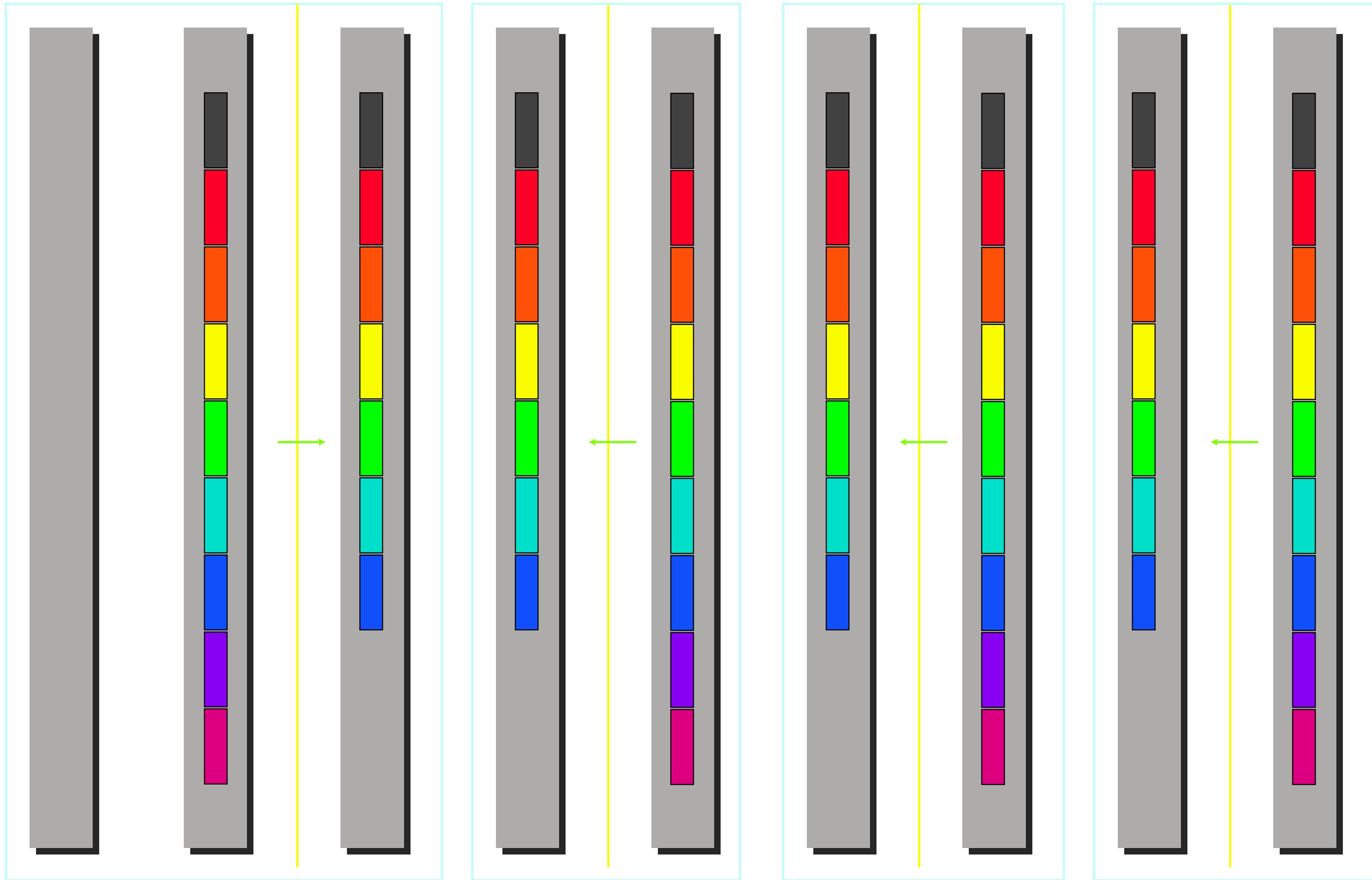


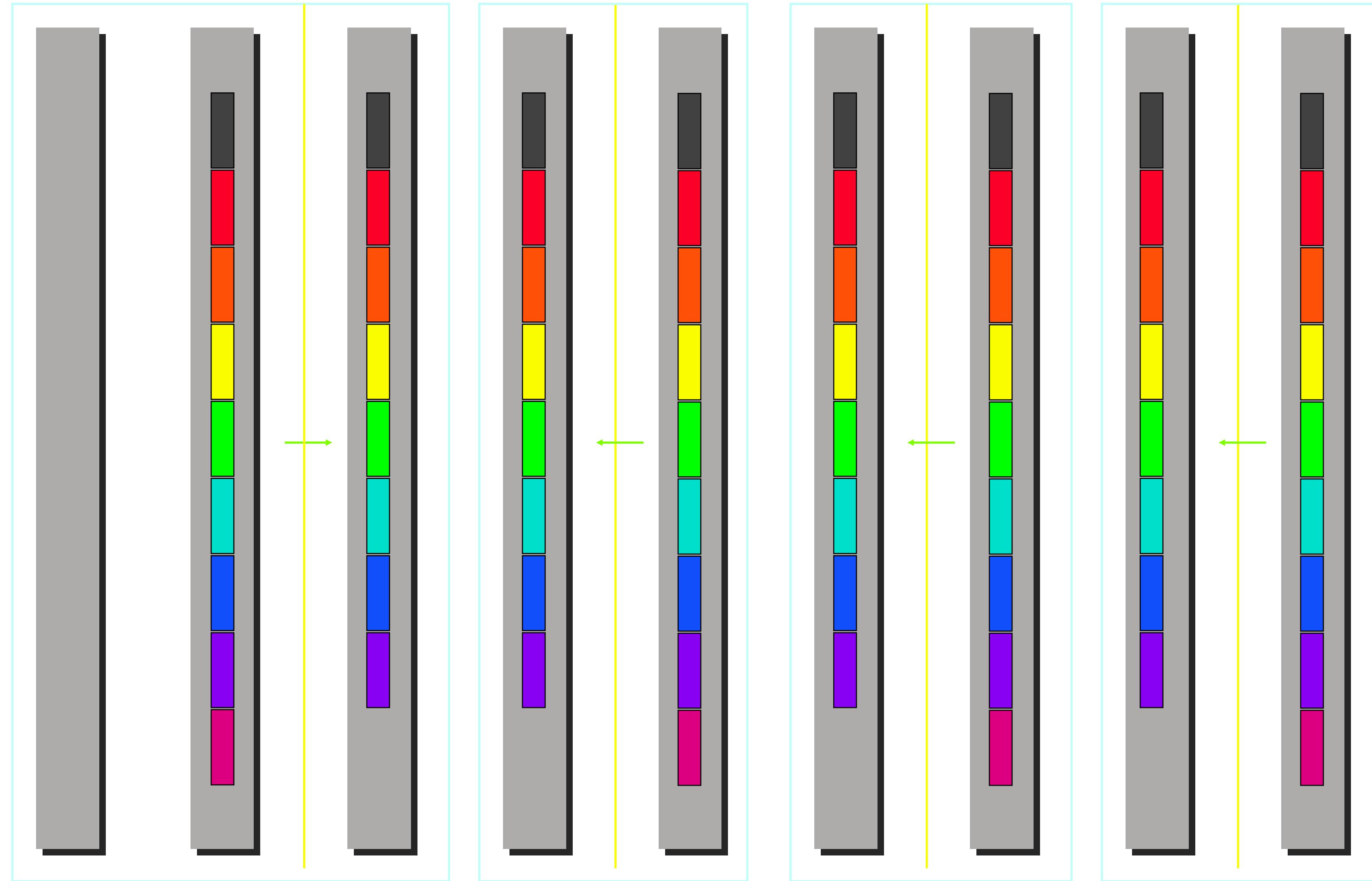


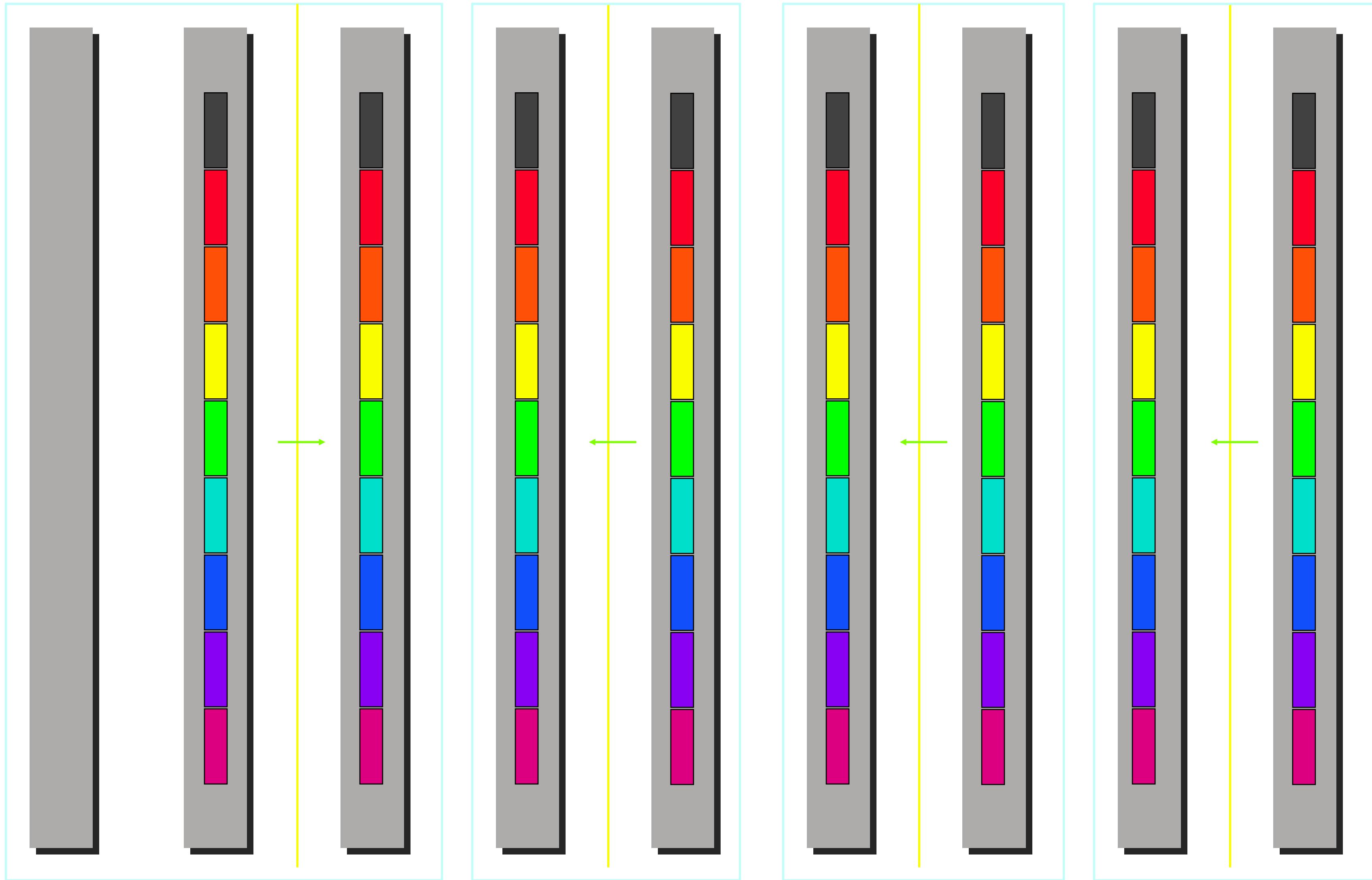


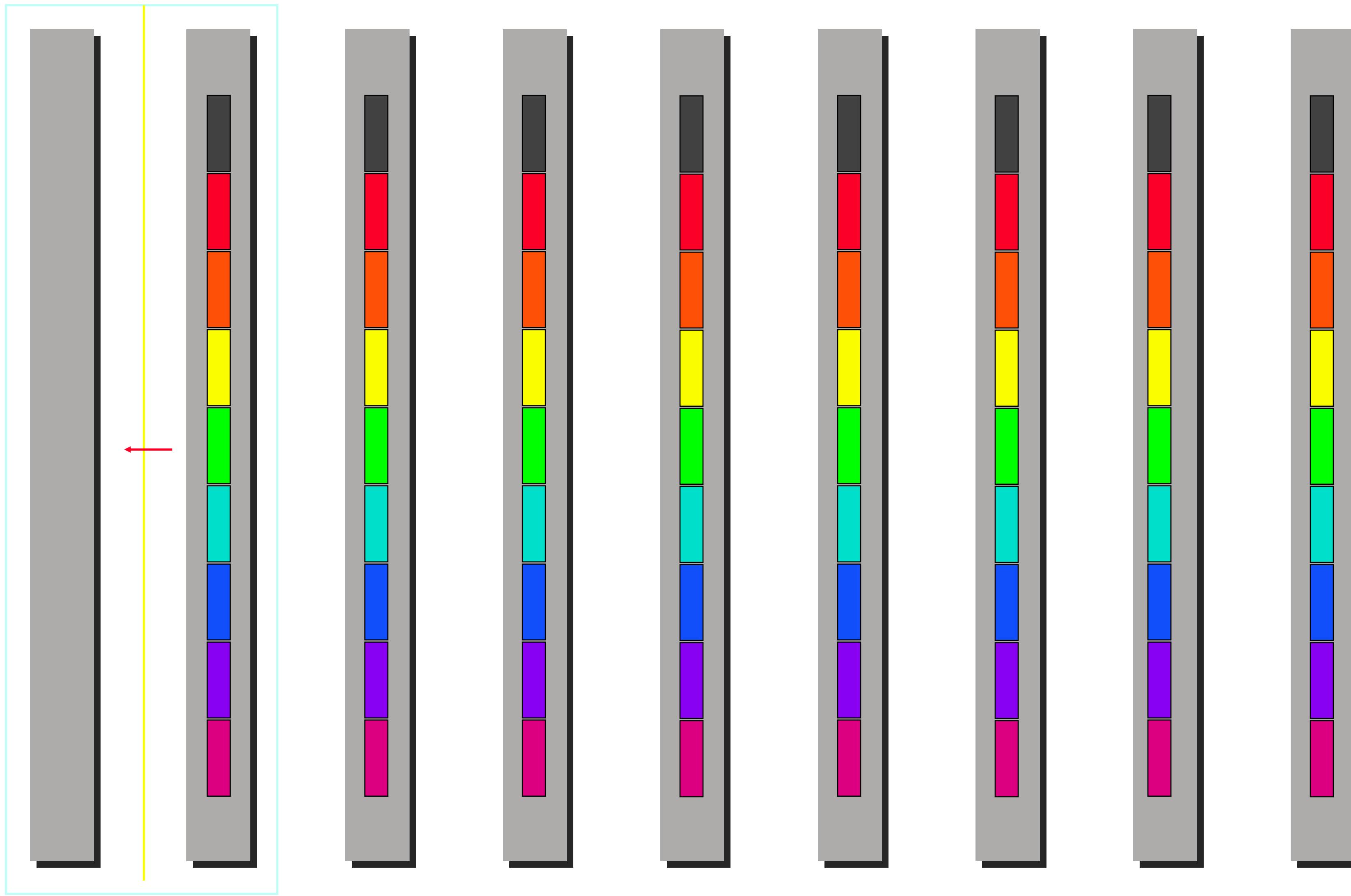


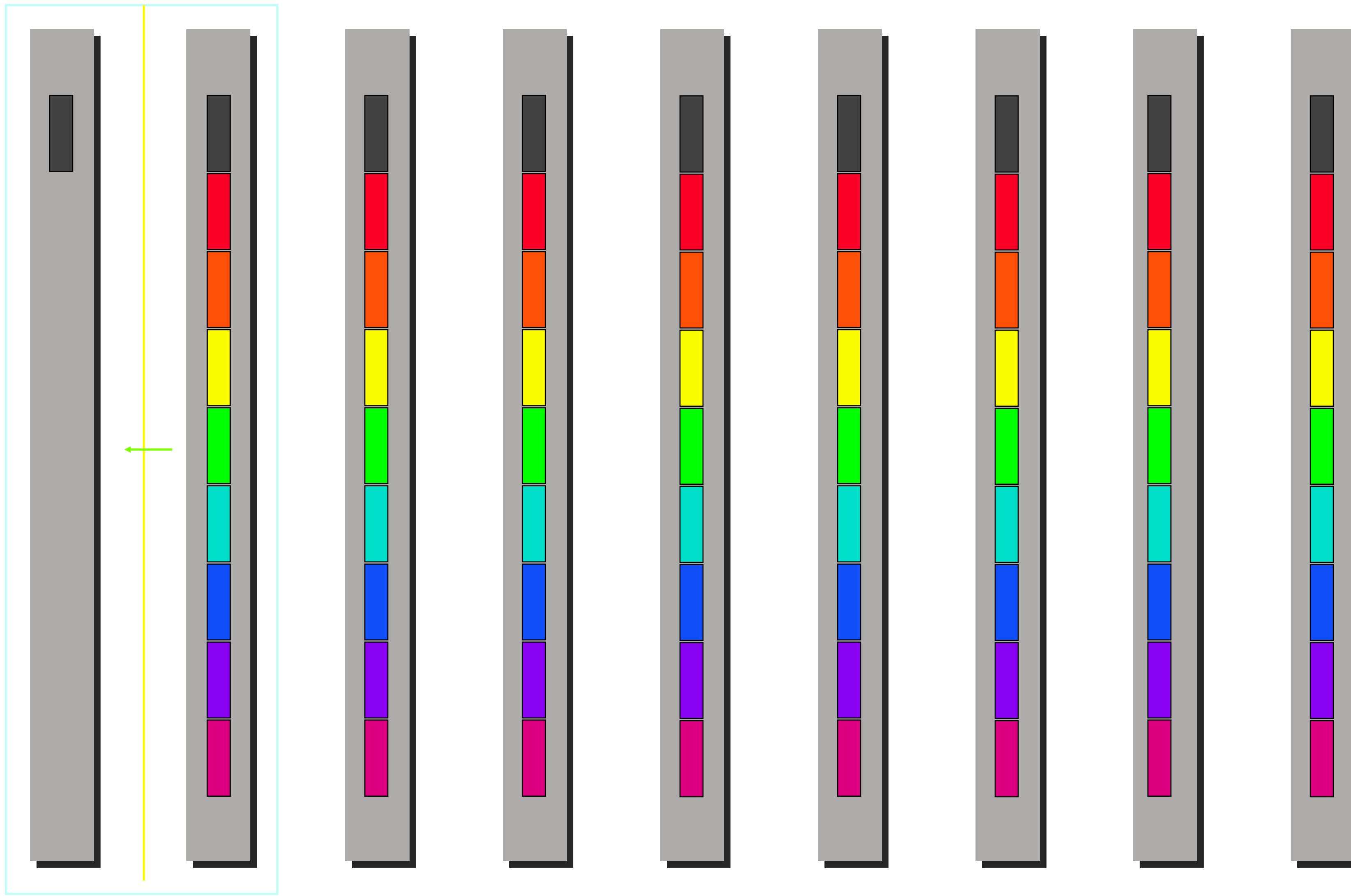


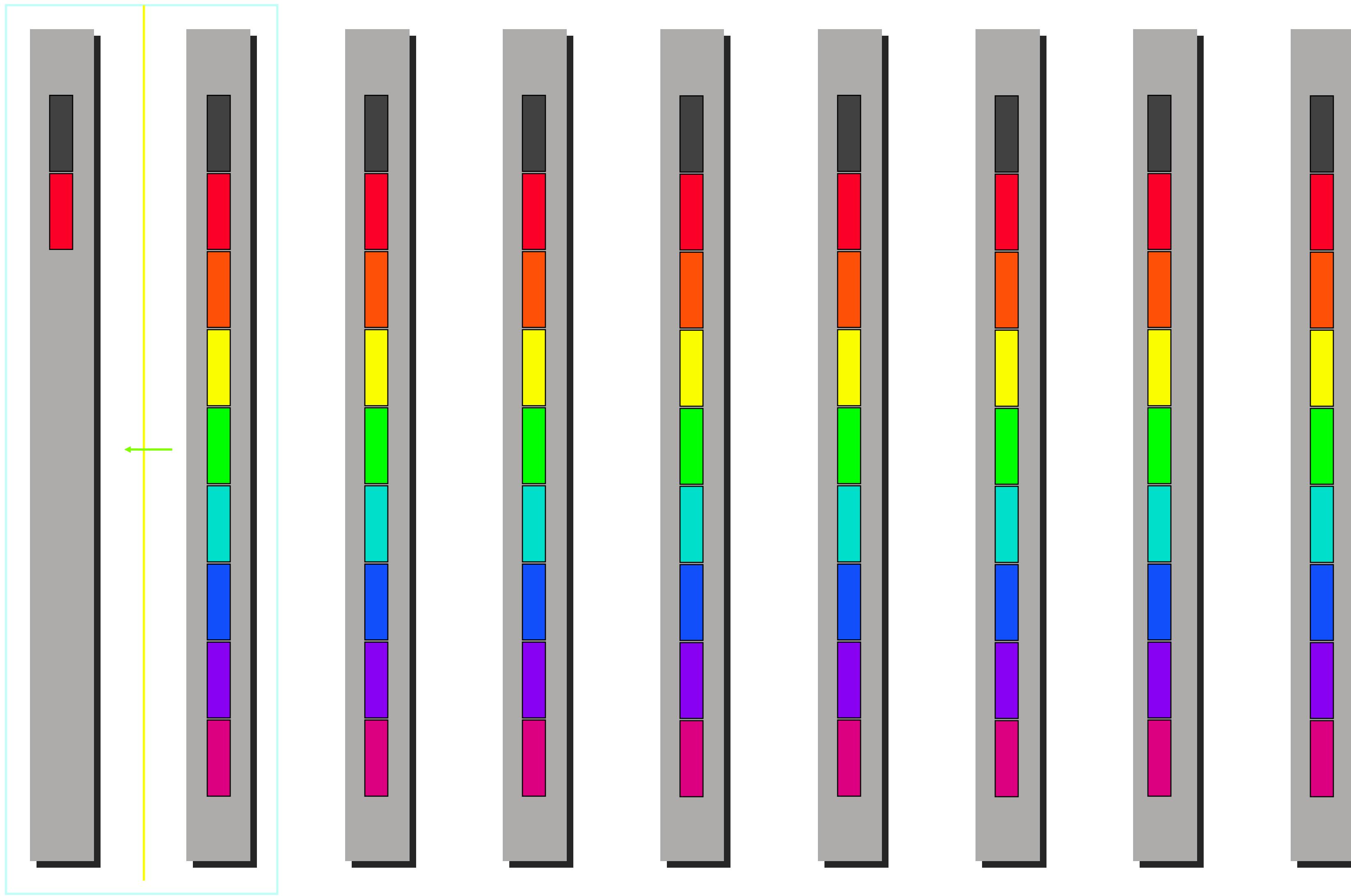


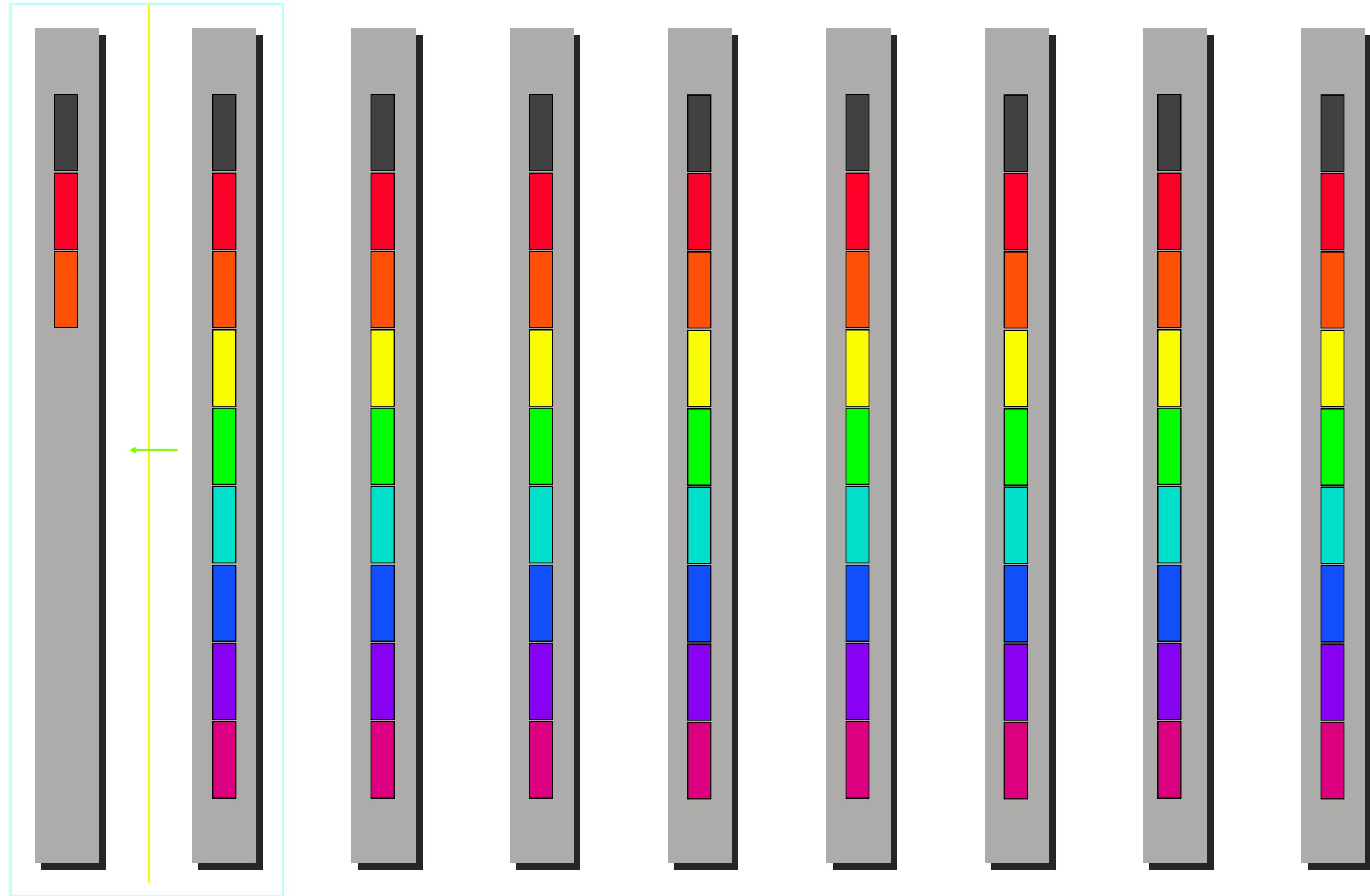


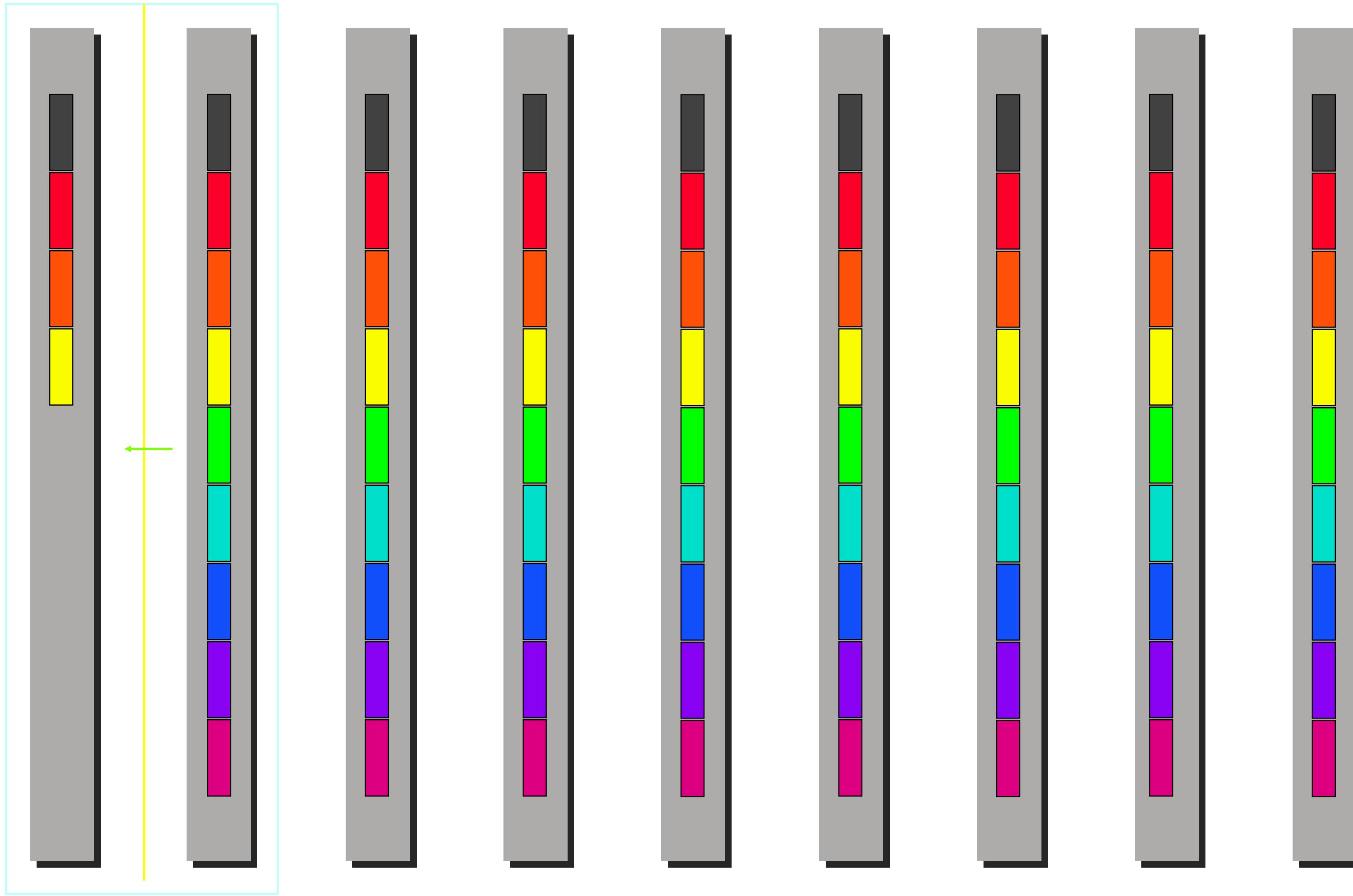


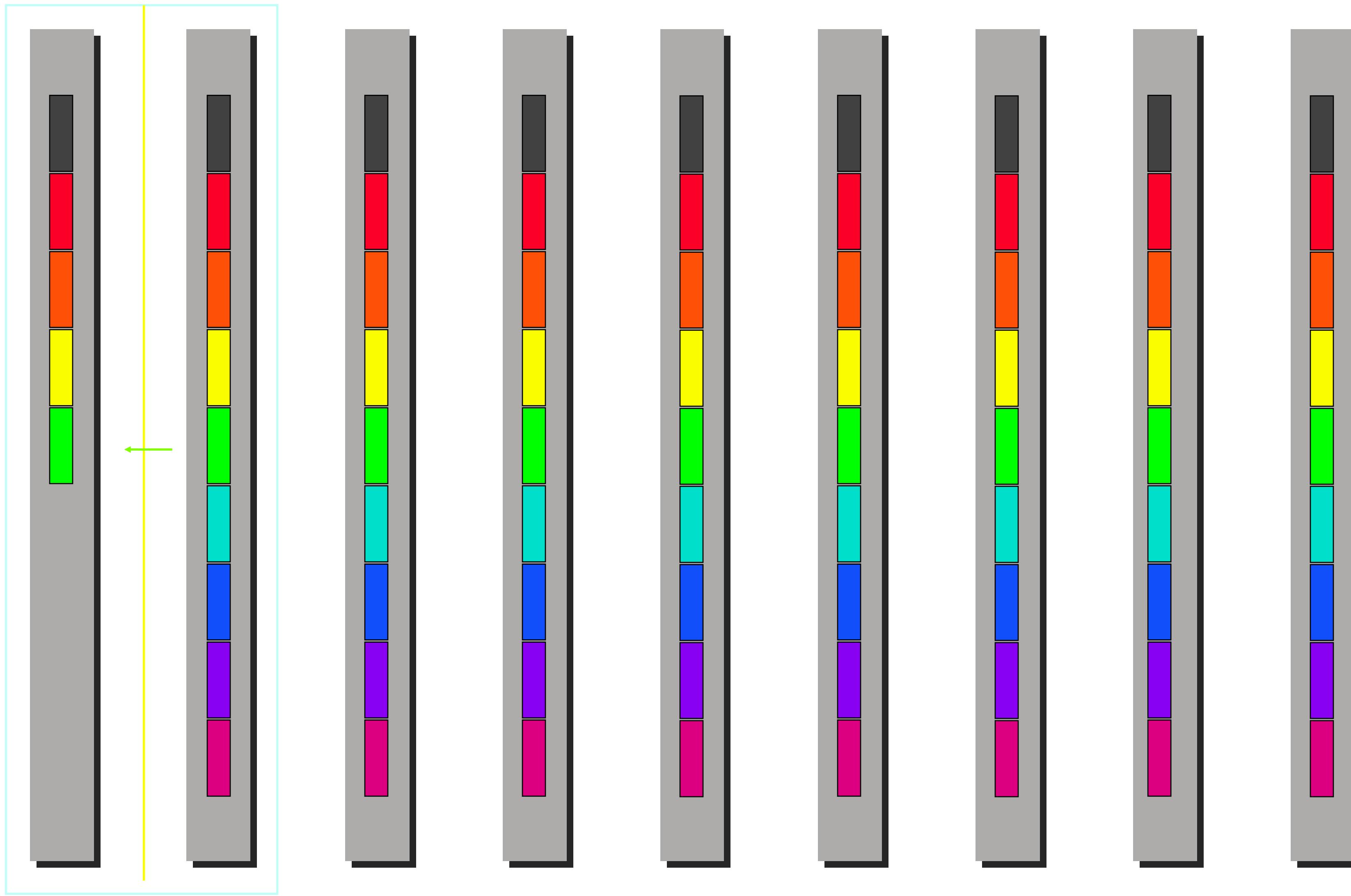


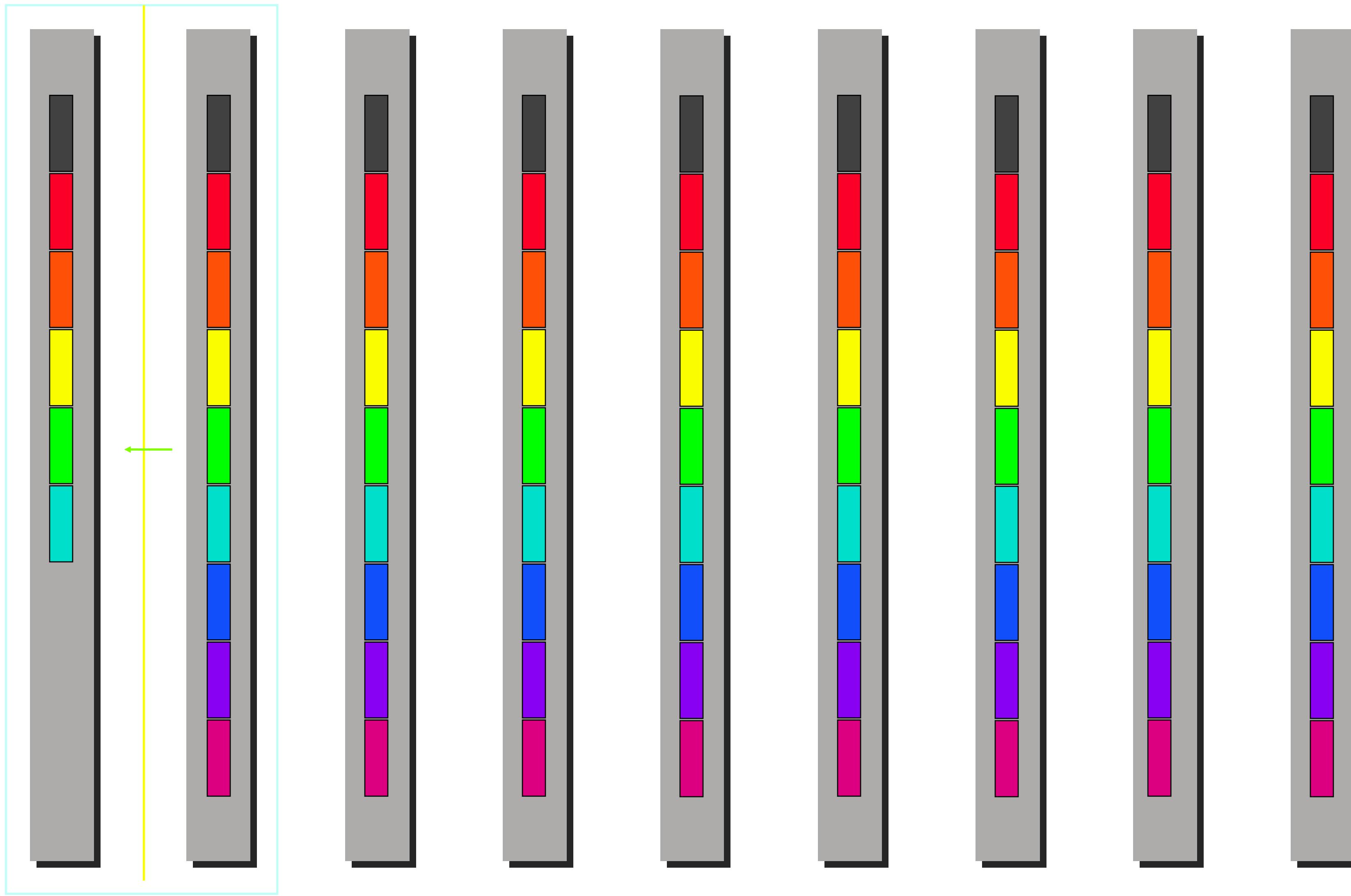


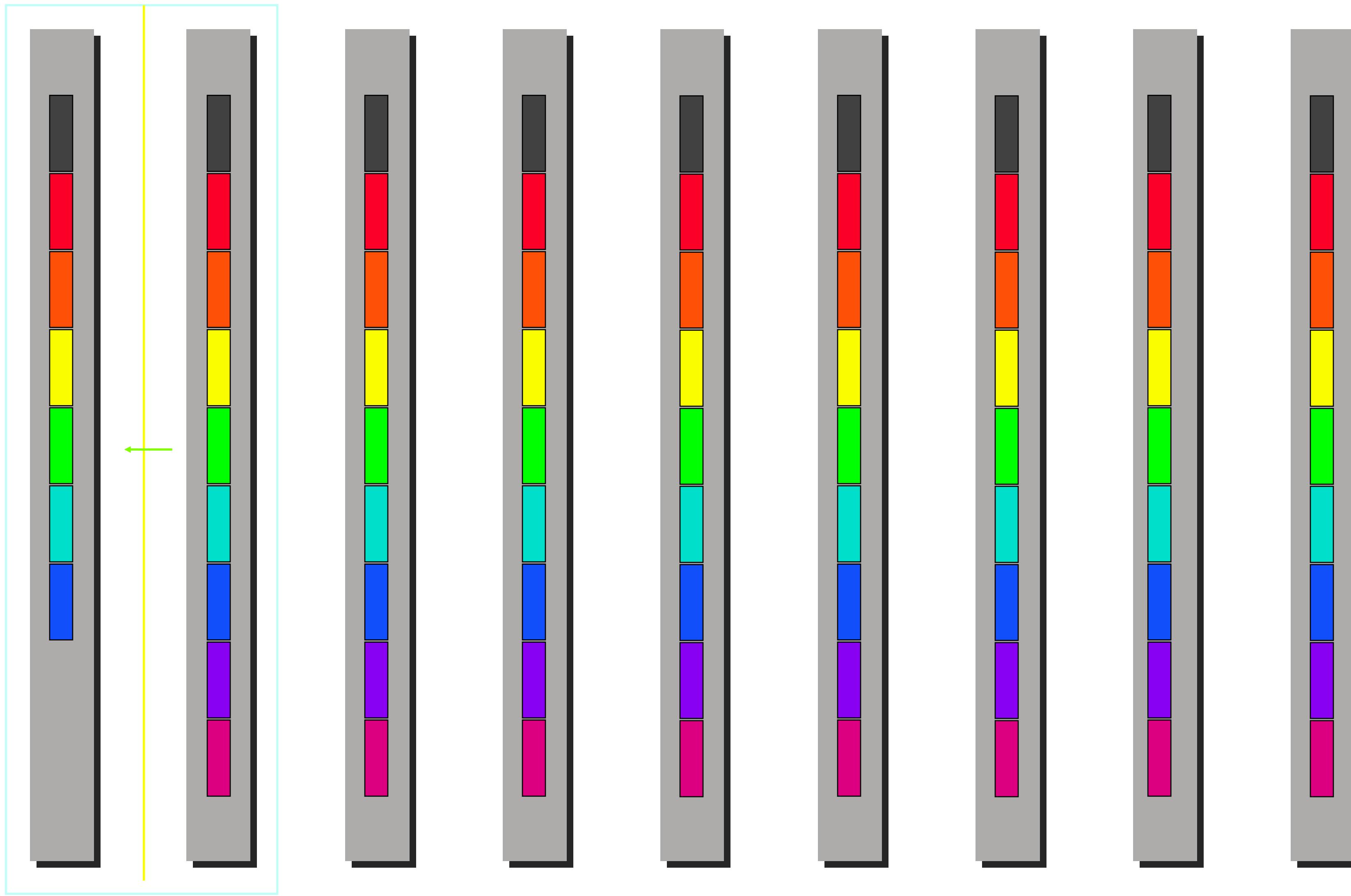


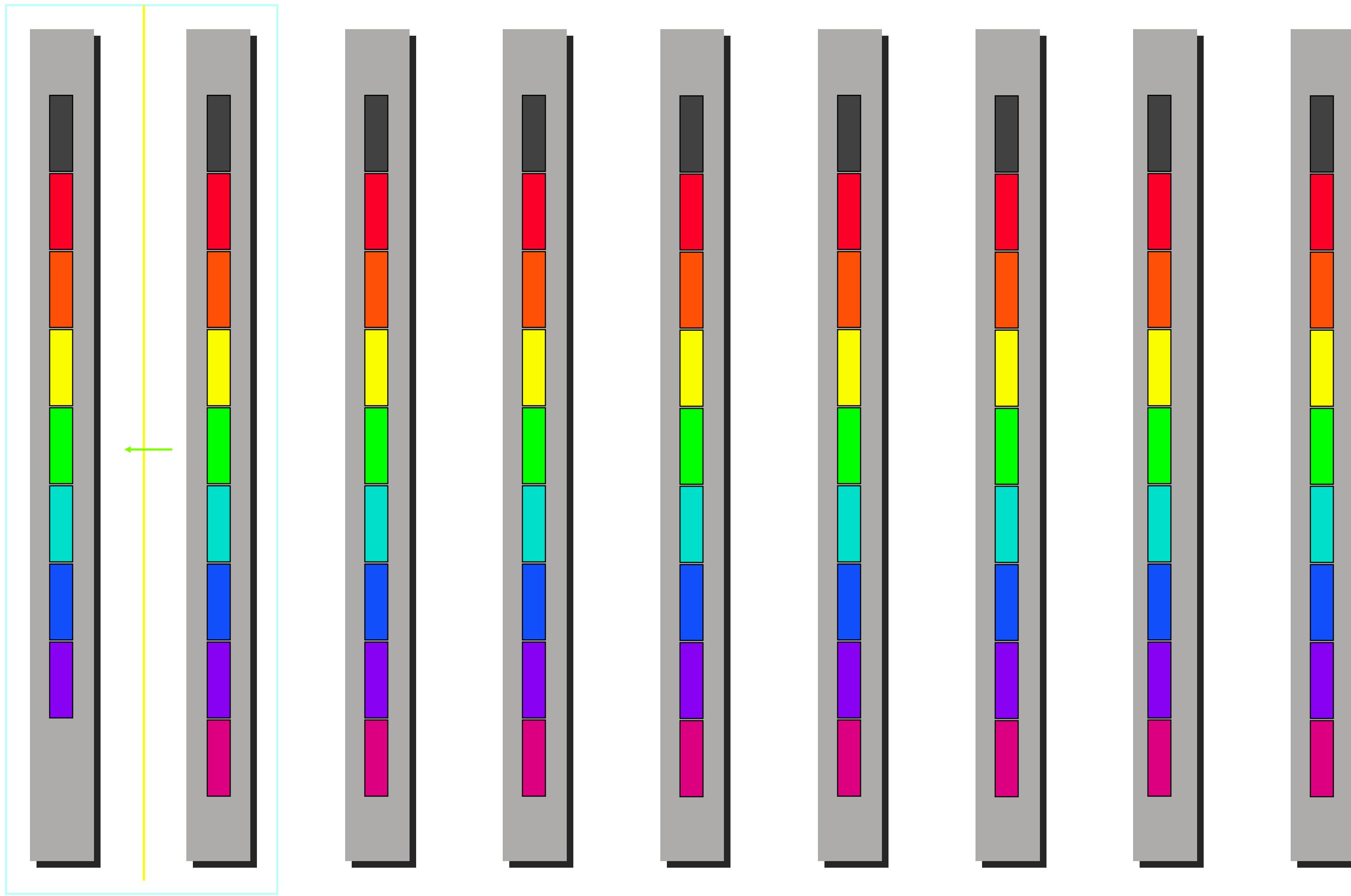


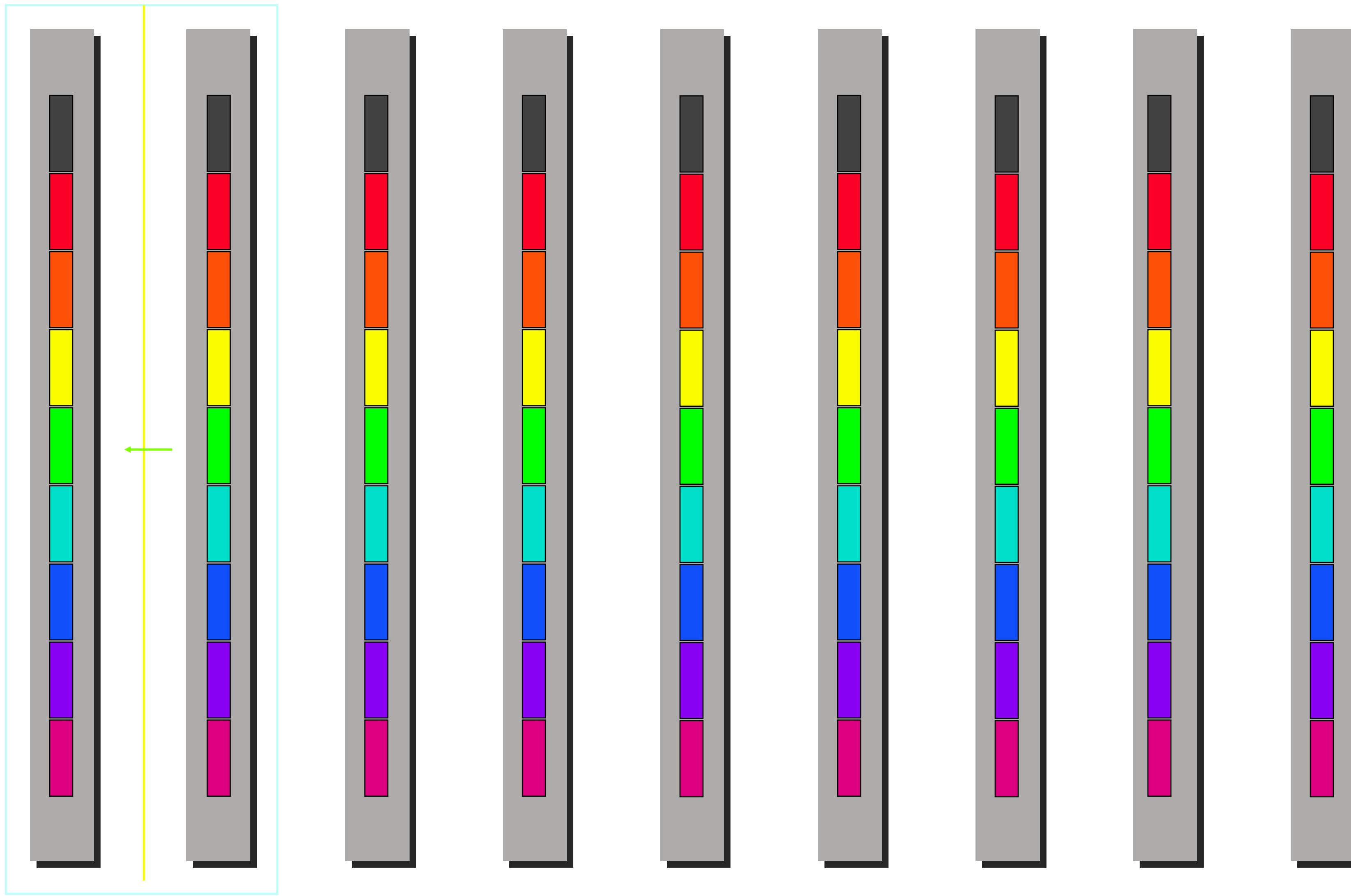


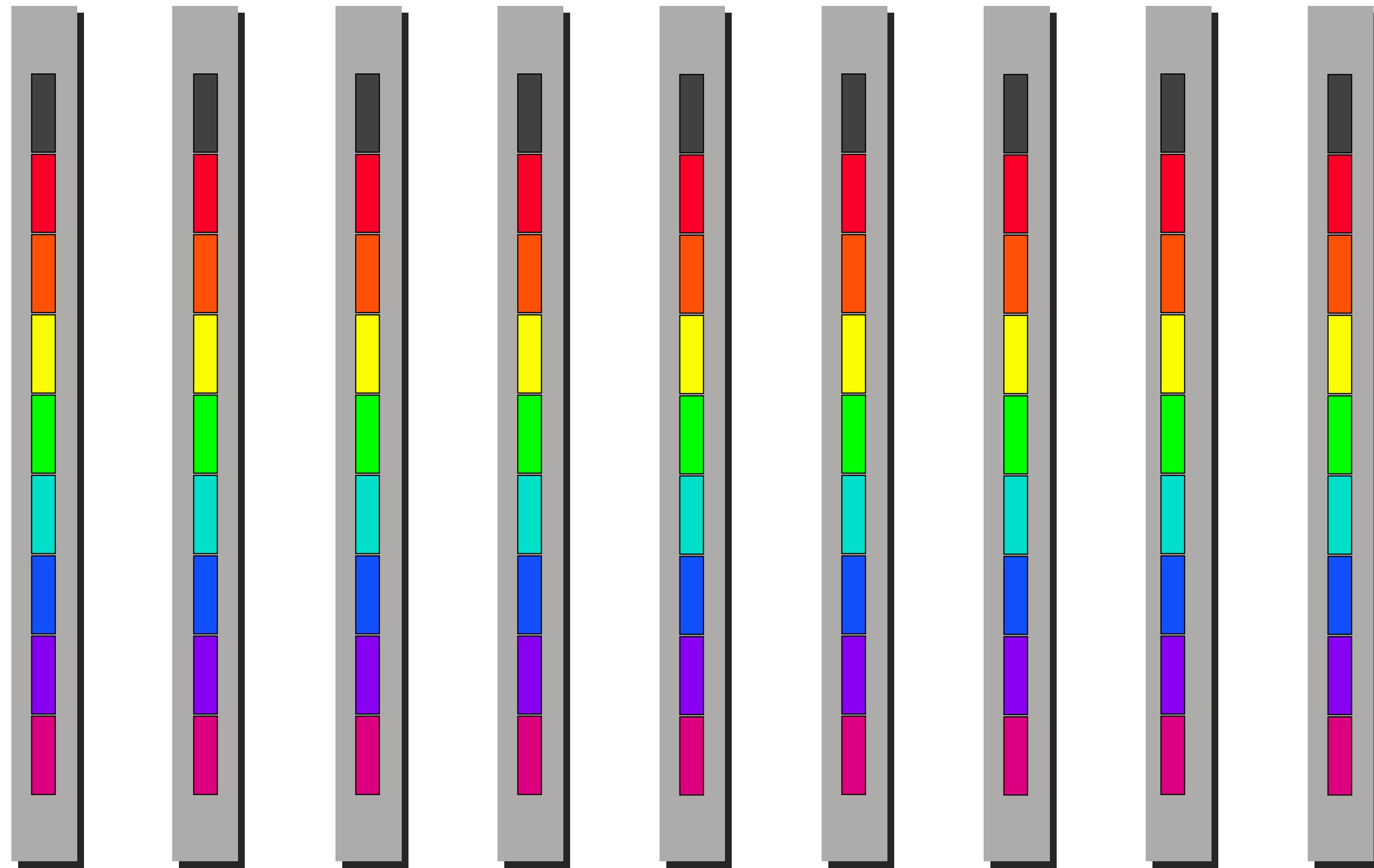












Cost of minimum spanning tree broadcast

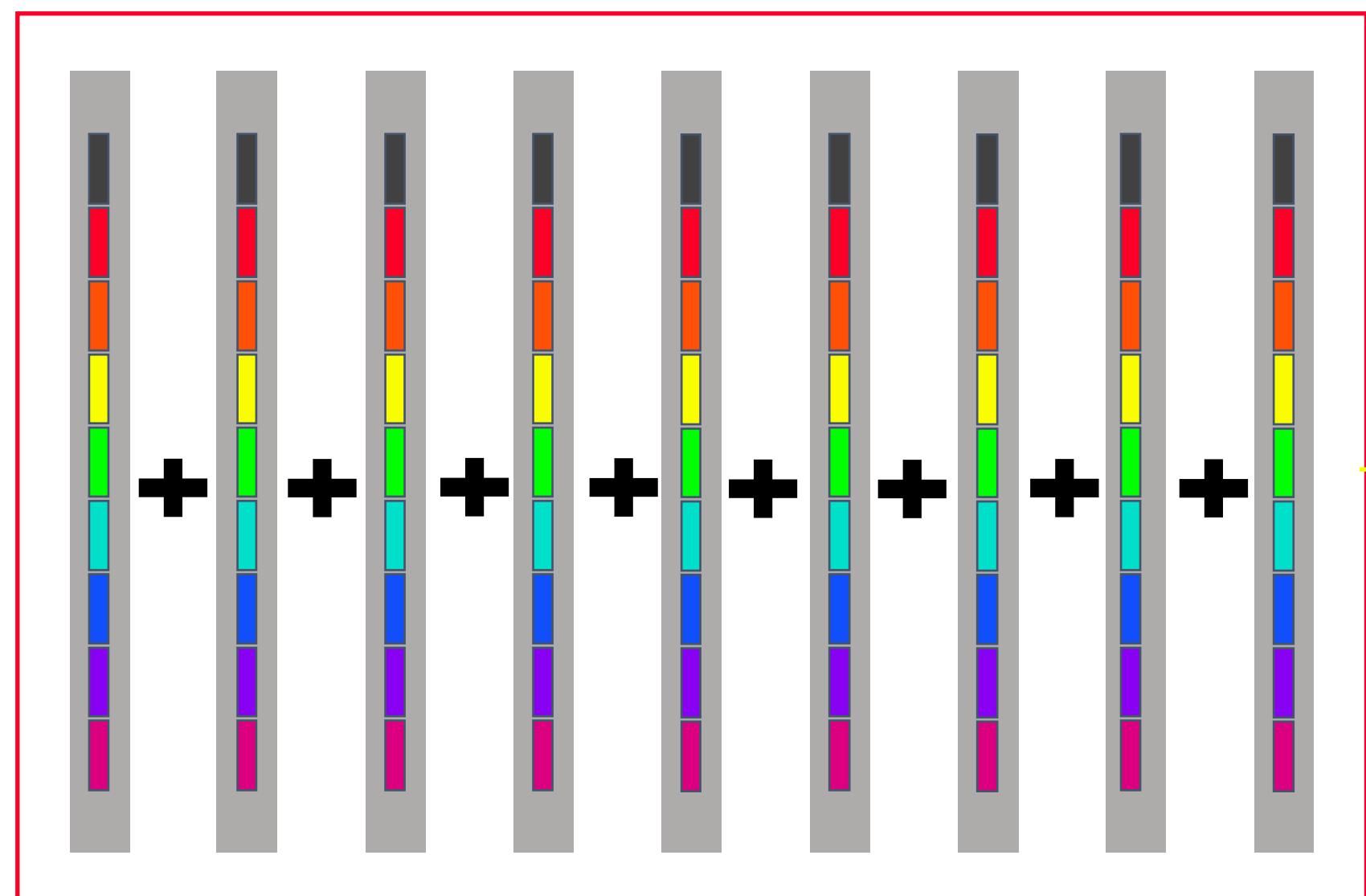
$$\left\lceil \log(p) \right\rceil (\alpha + n\beta)$$

number of steps

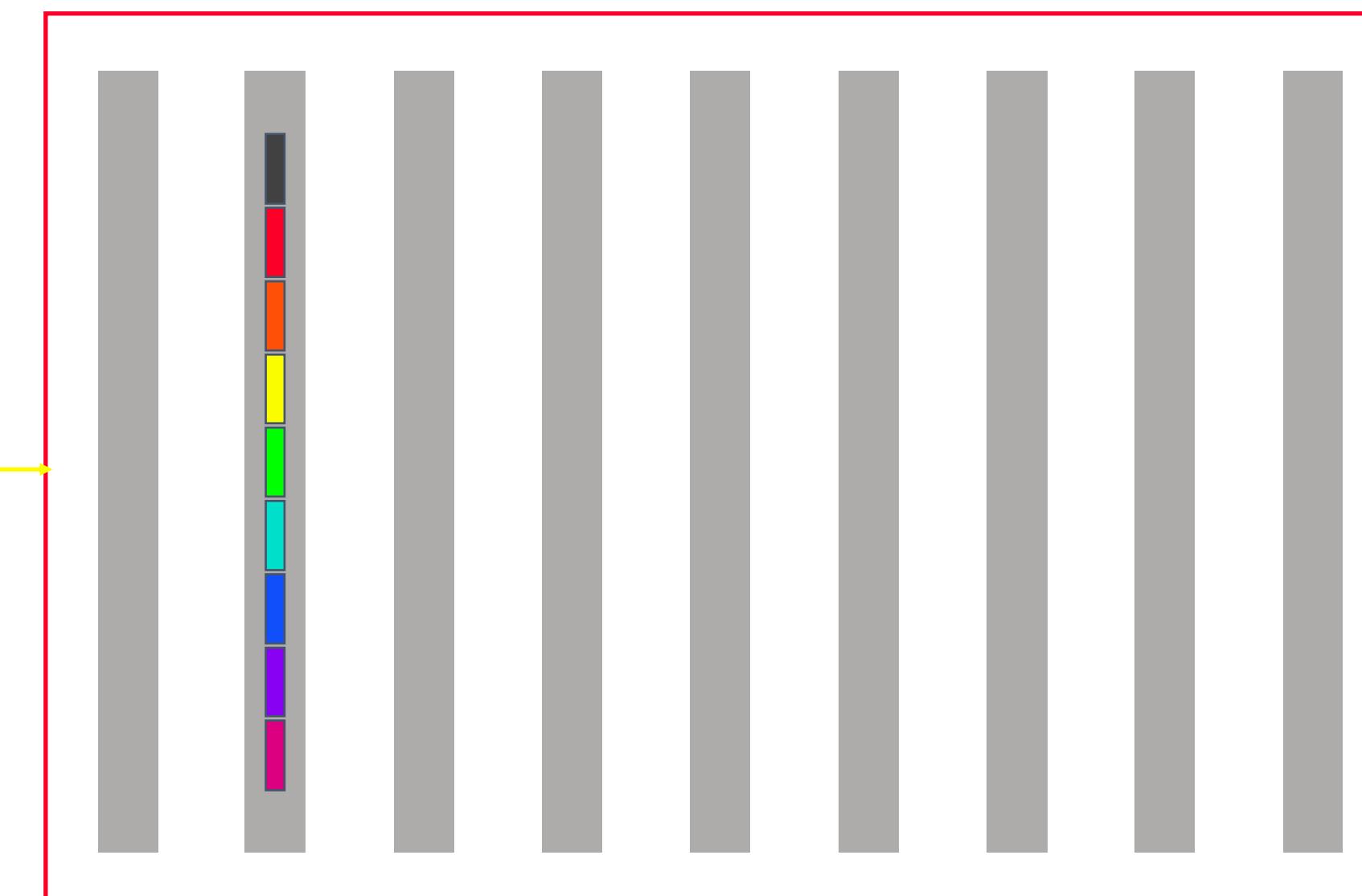
cost per steps

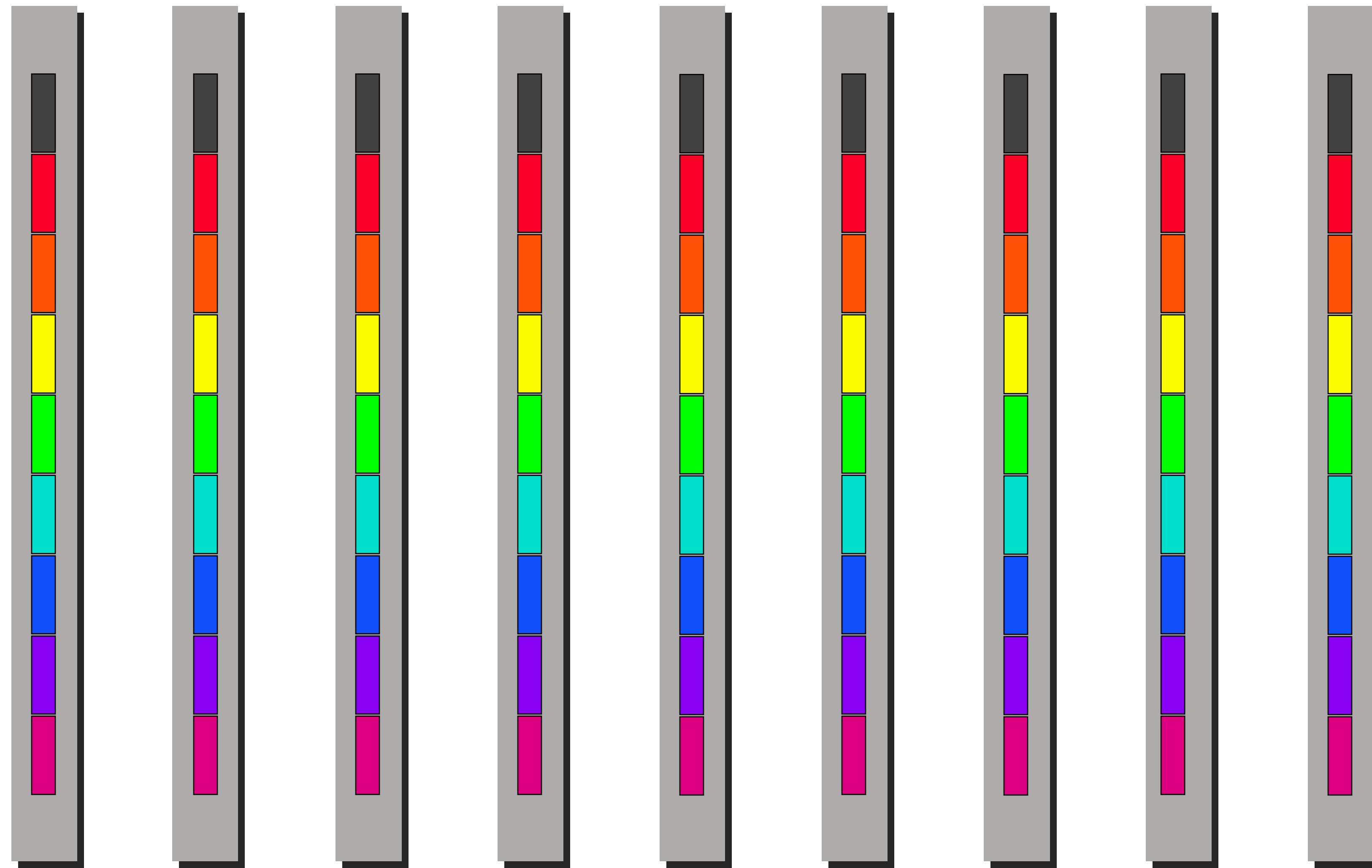
Reduce(-to-one)

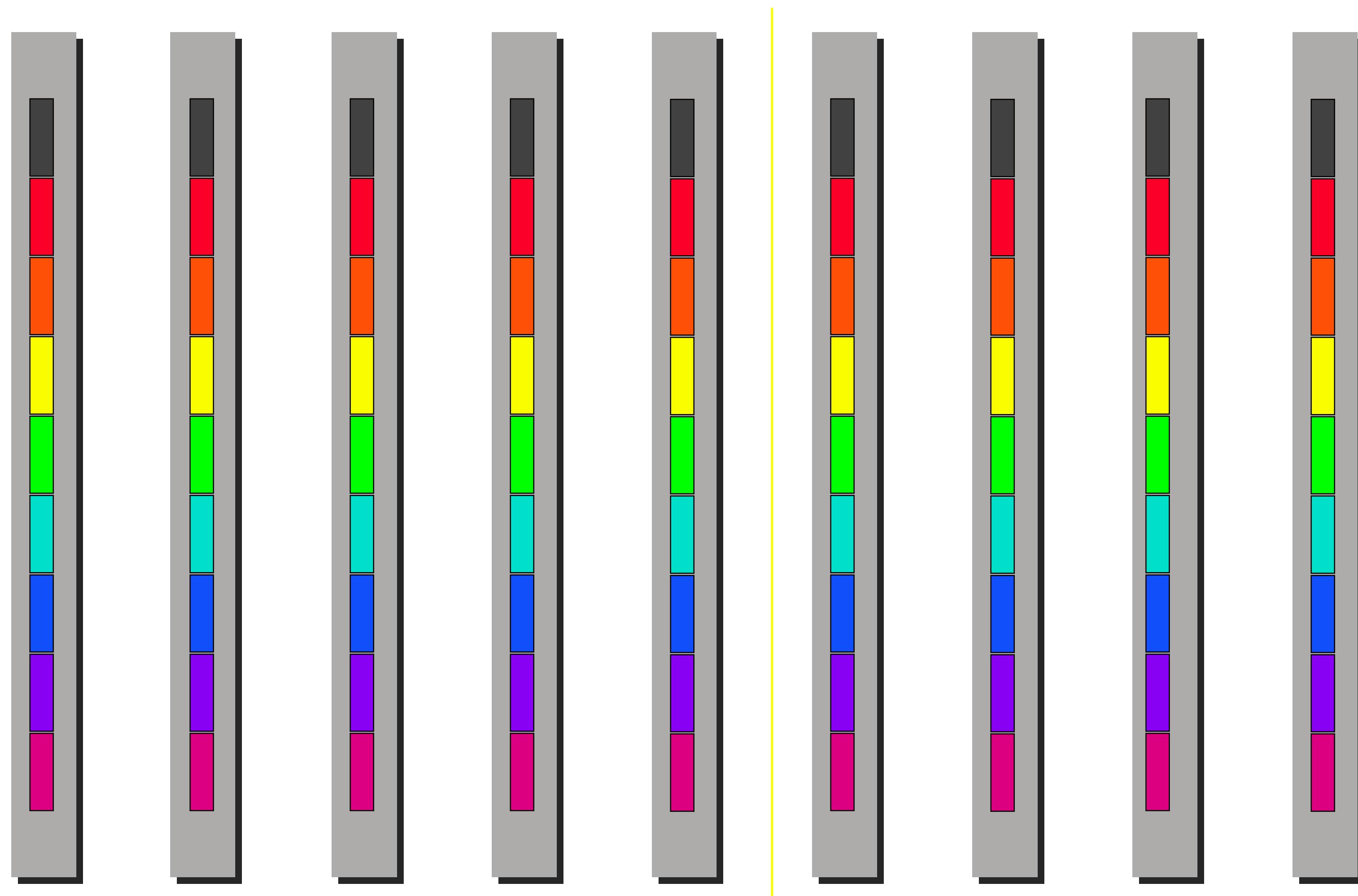
Before

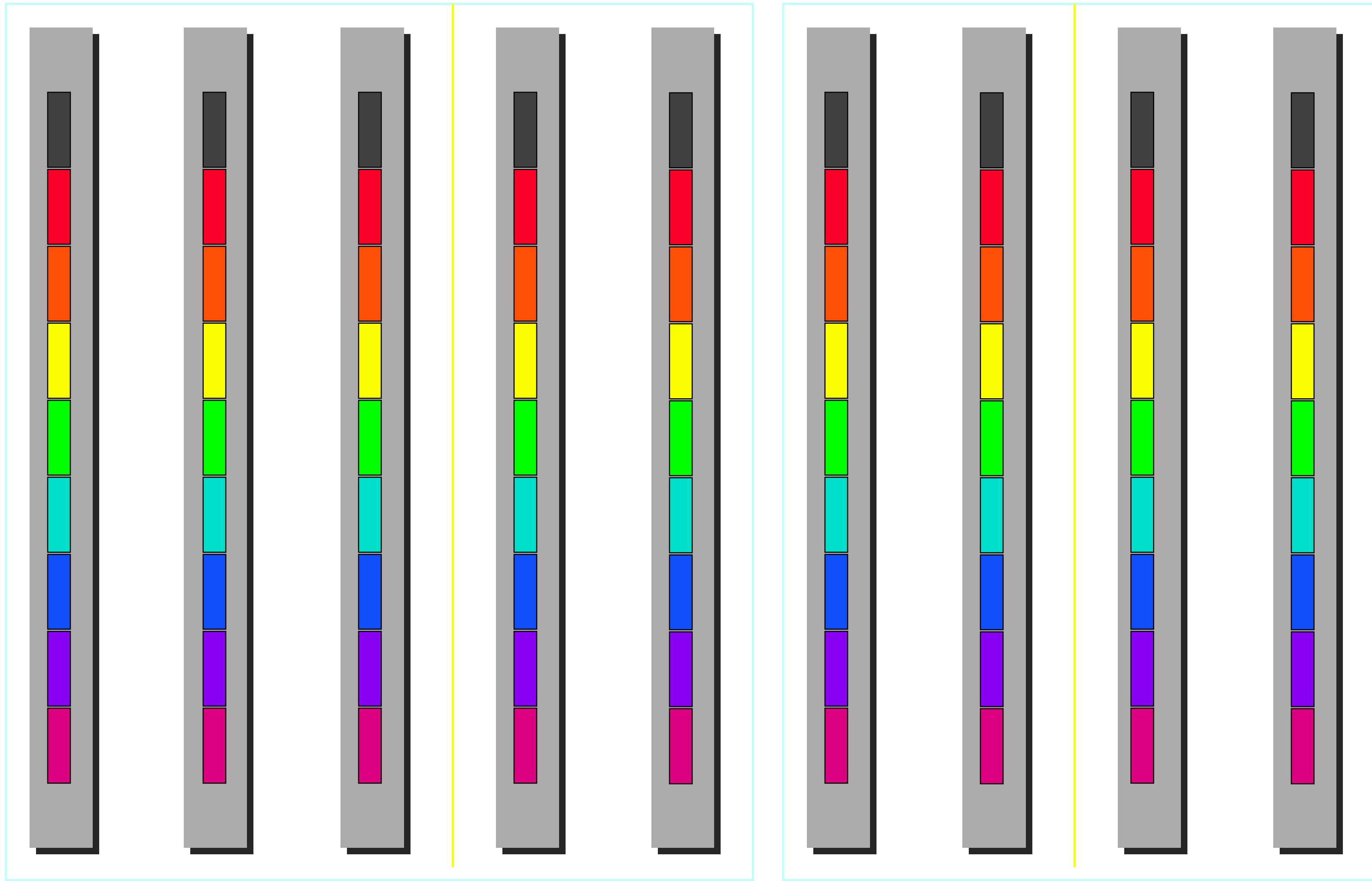


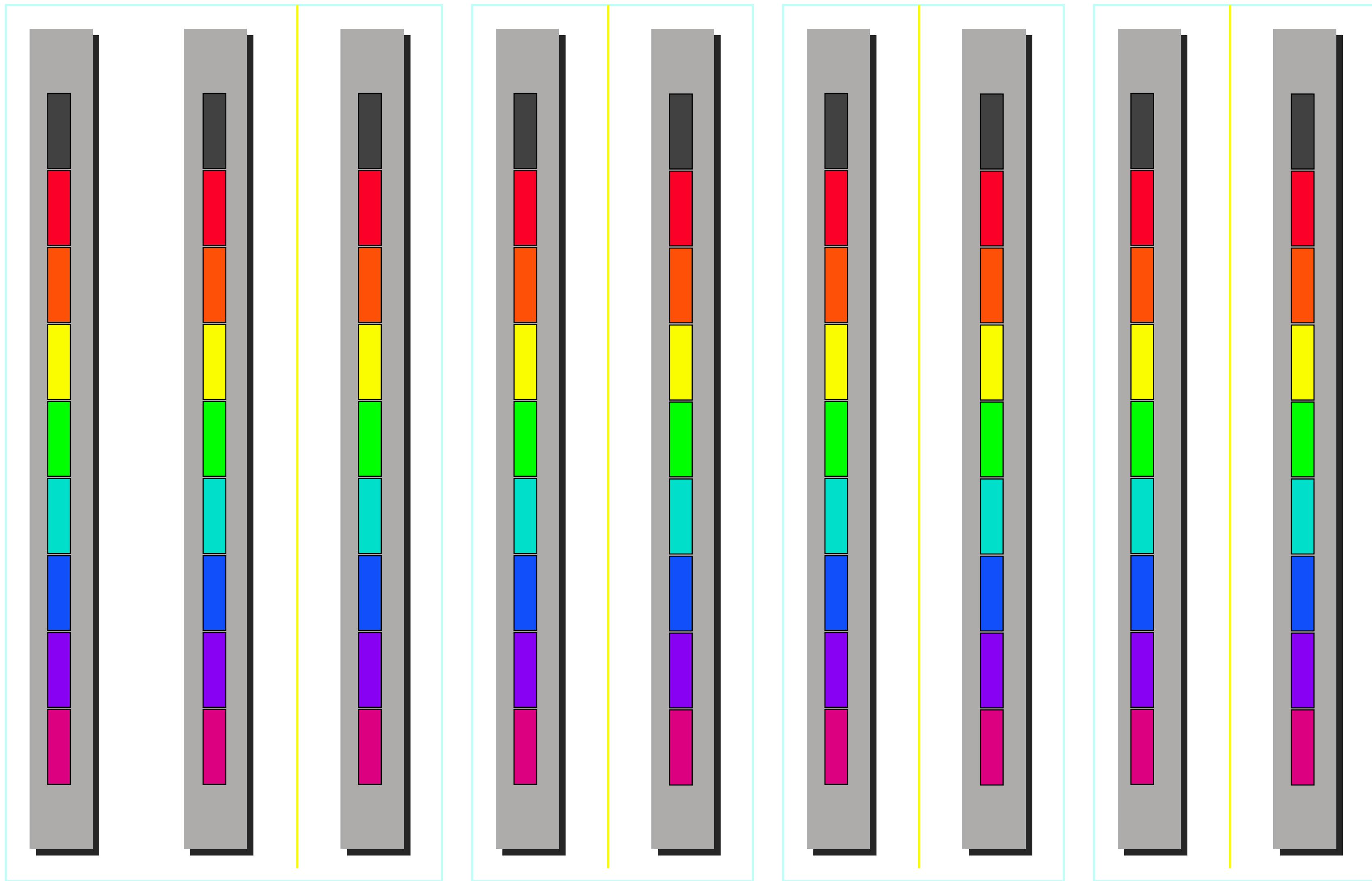
After

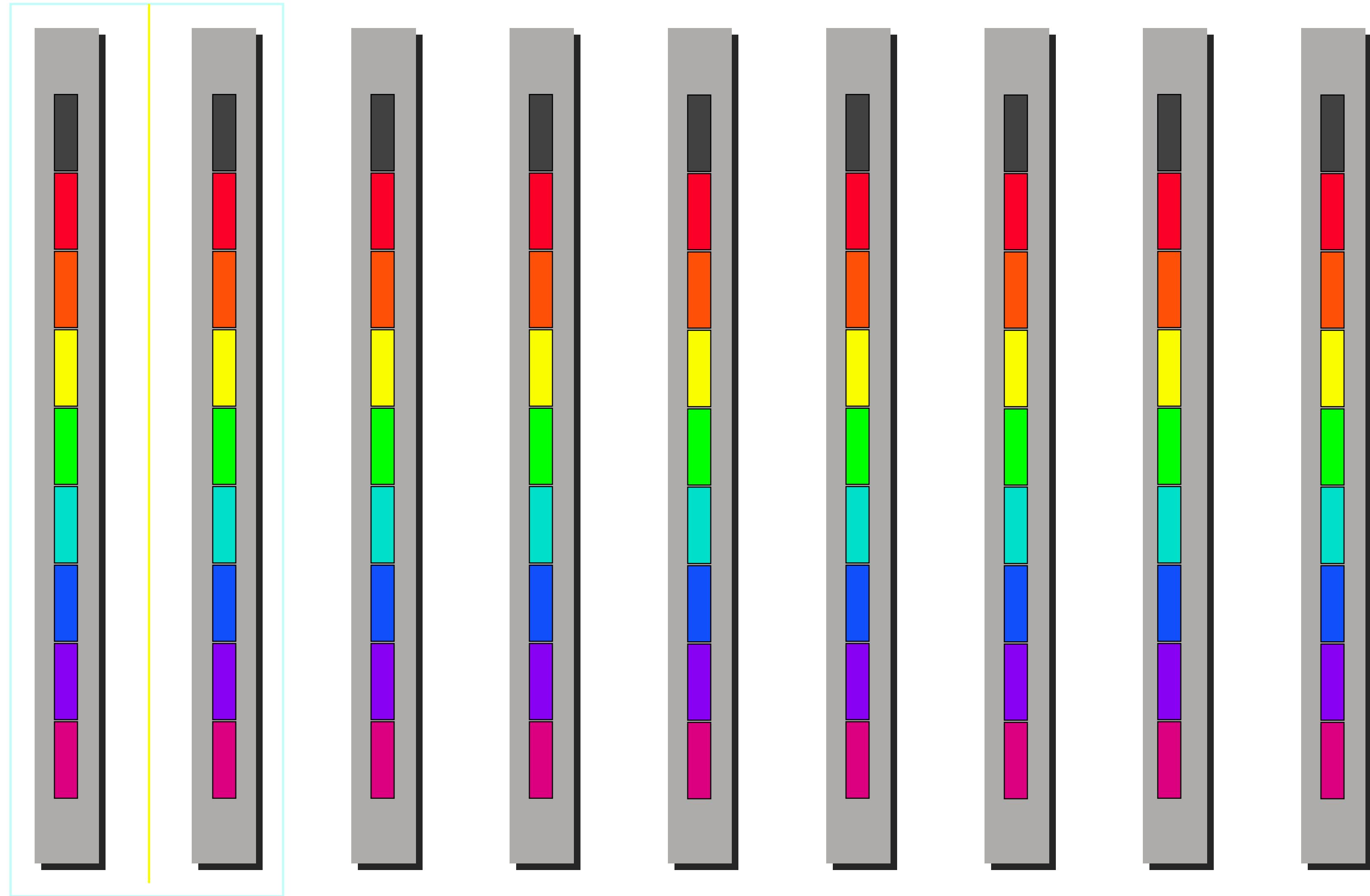


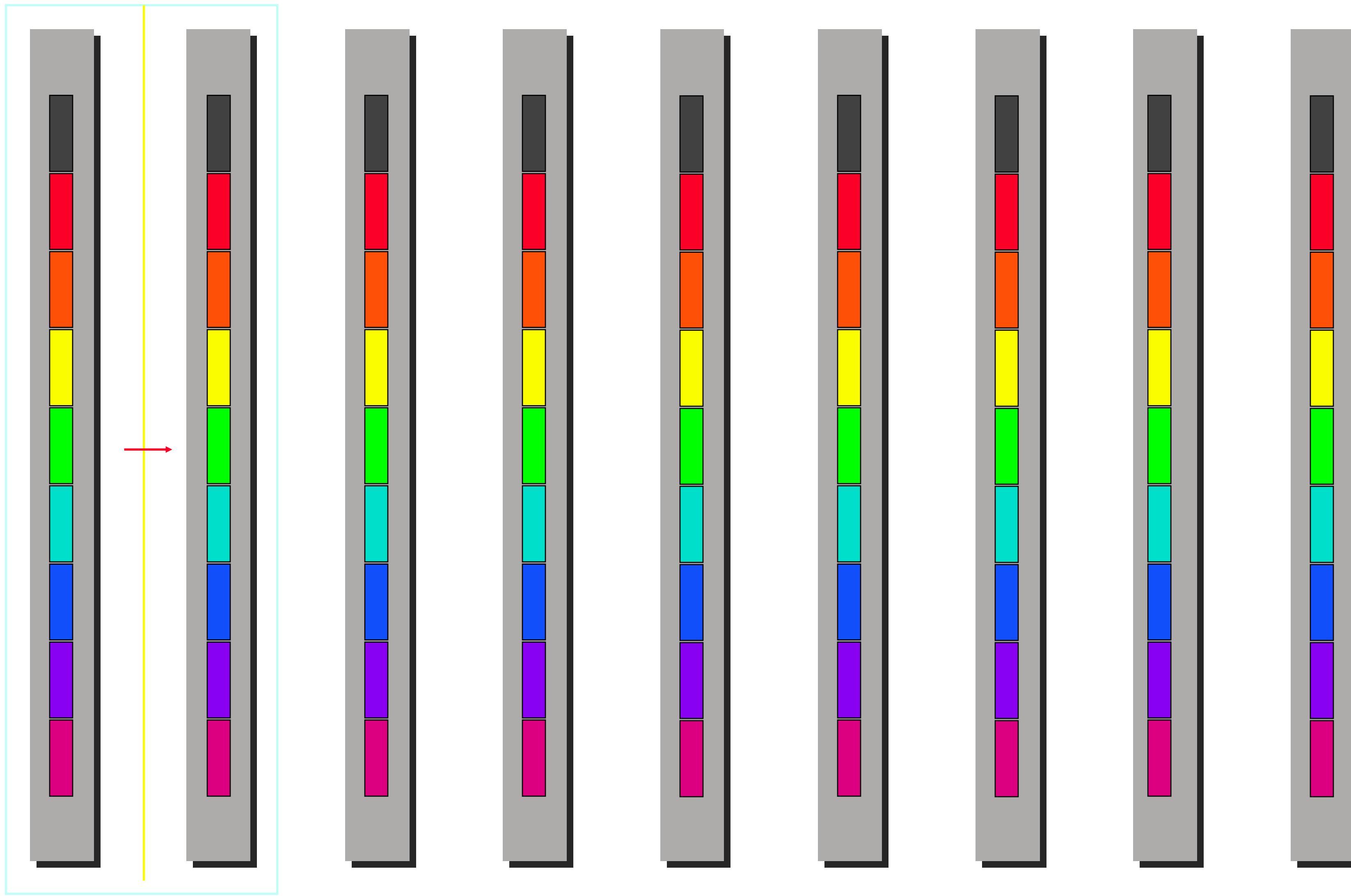


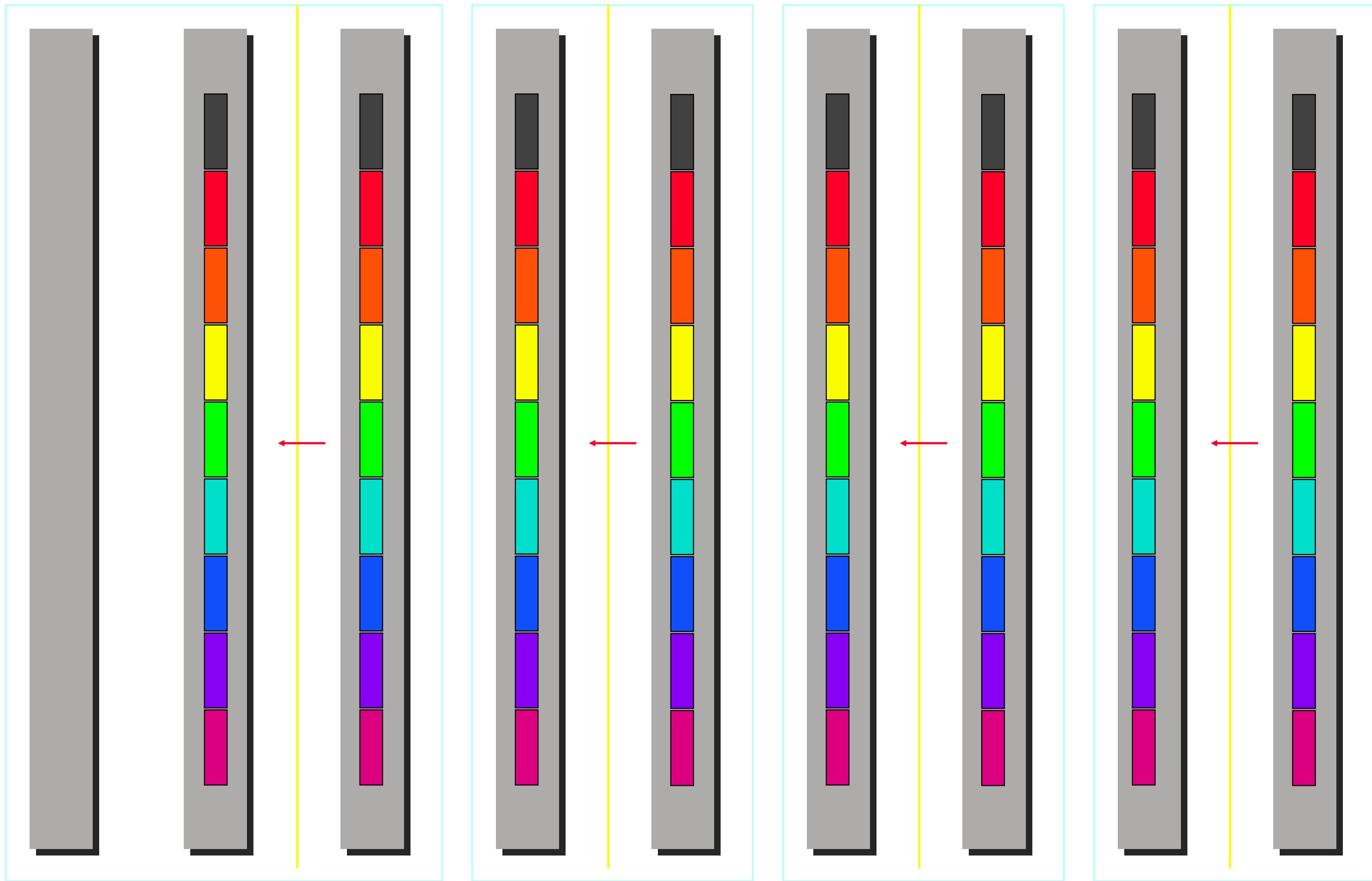


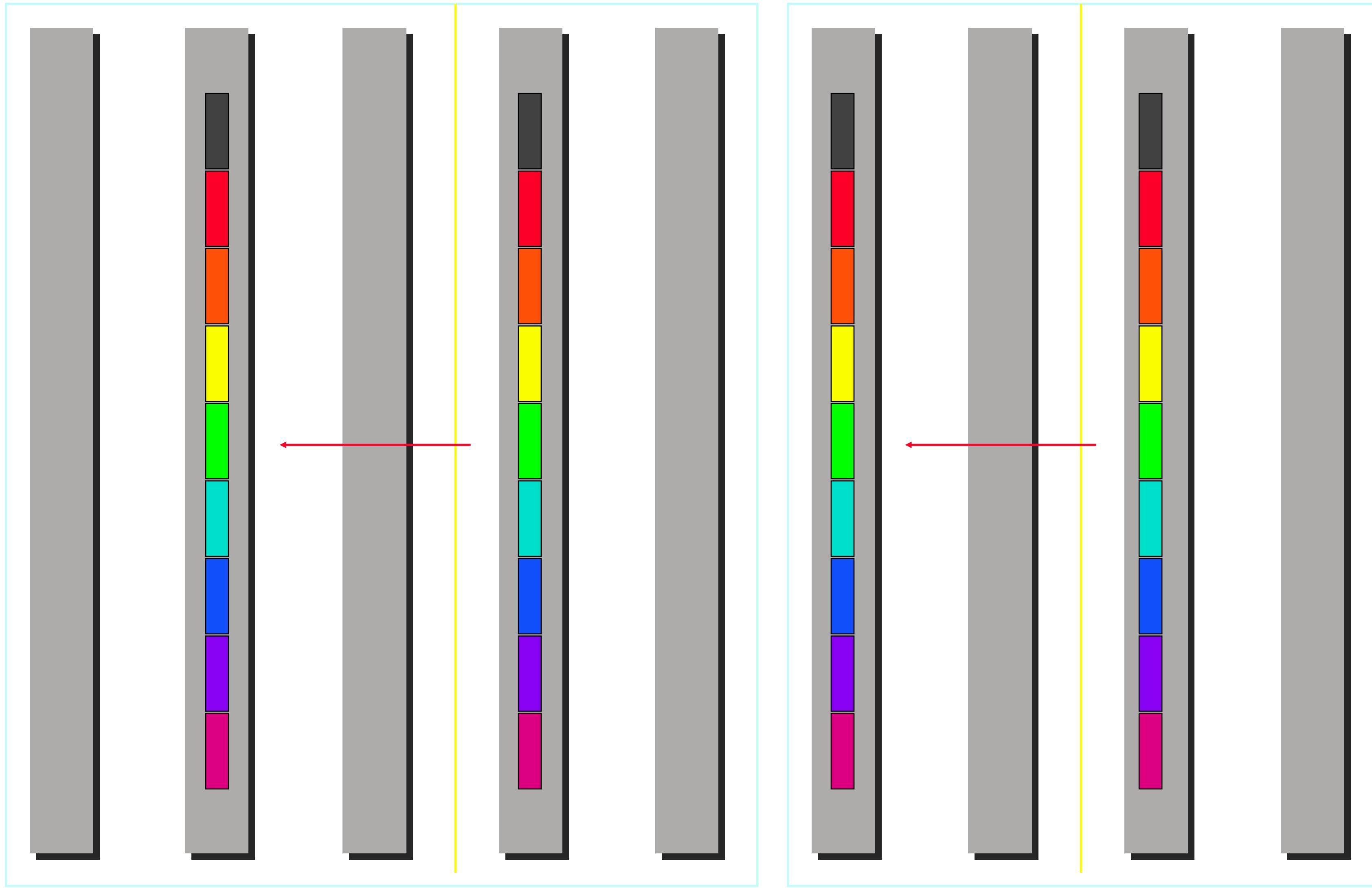


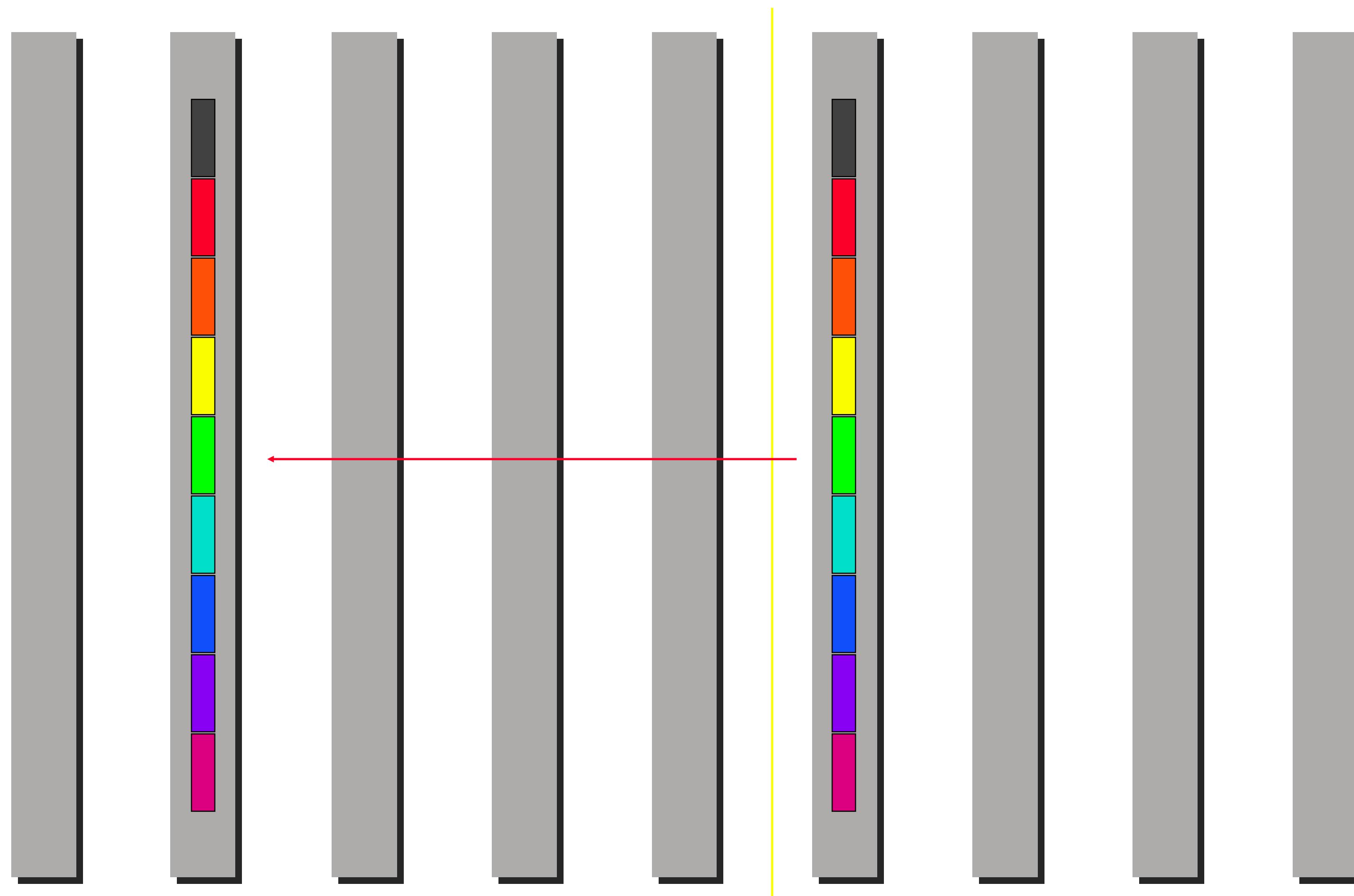










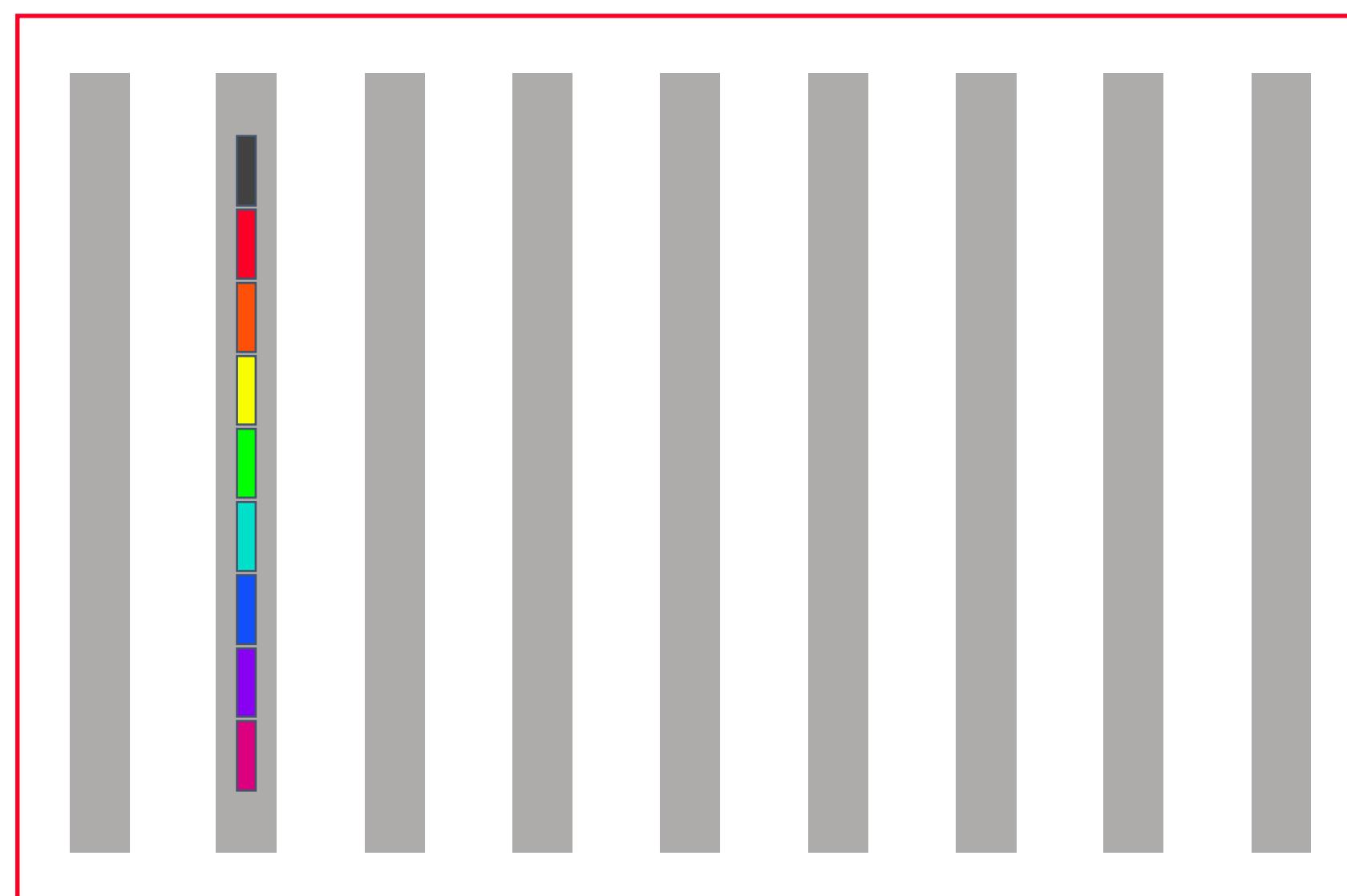


Cost of minimum spanning tree reduce(-to-one)

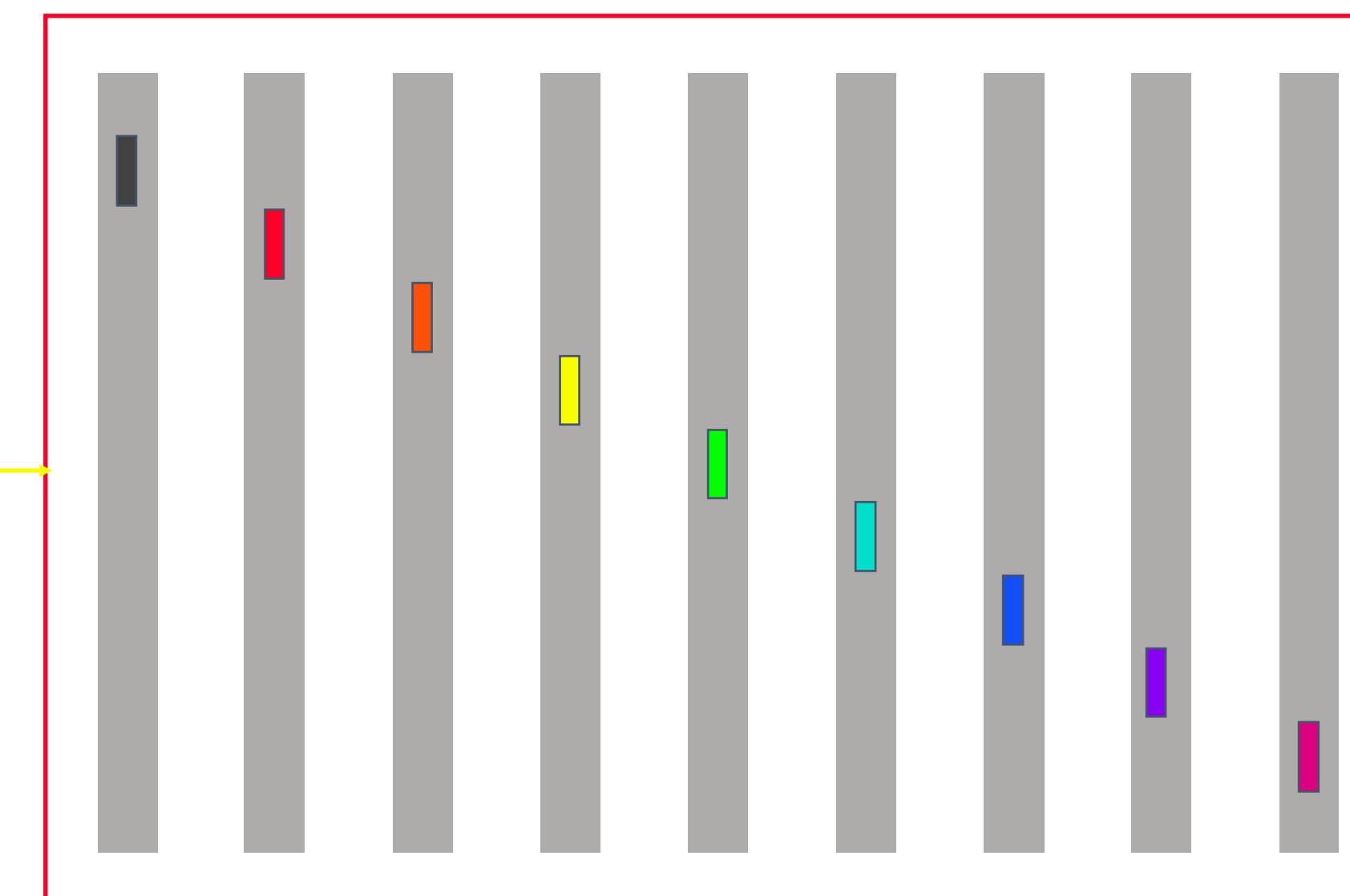
$$\lceil \log(p) \rceil (\alpha + n\beta + n\gamma)$$

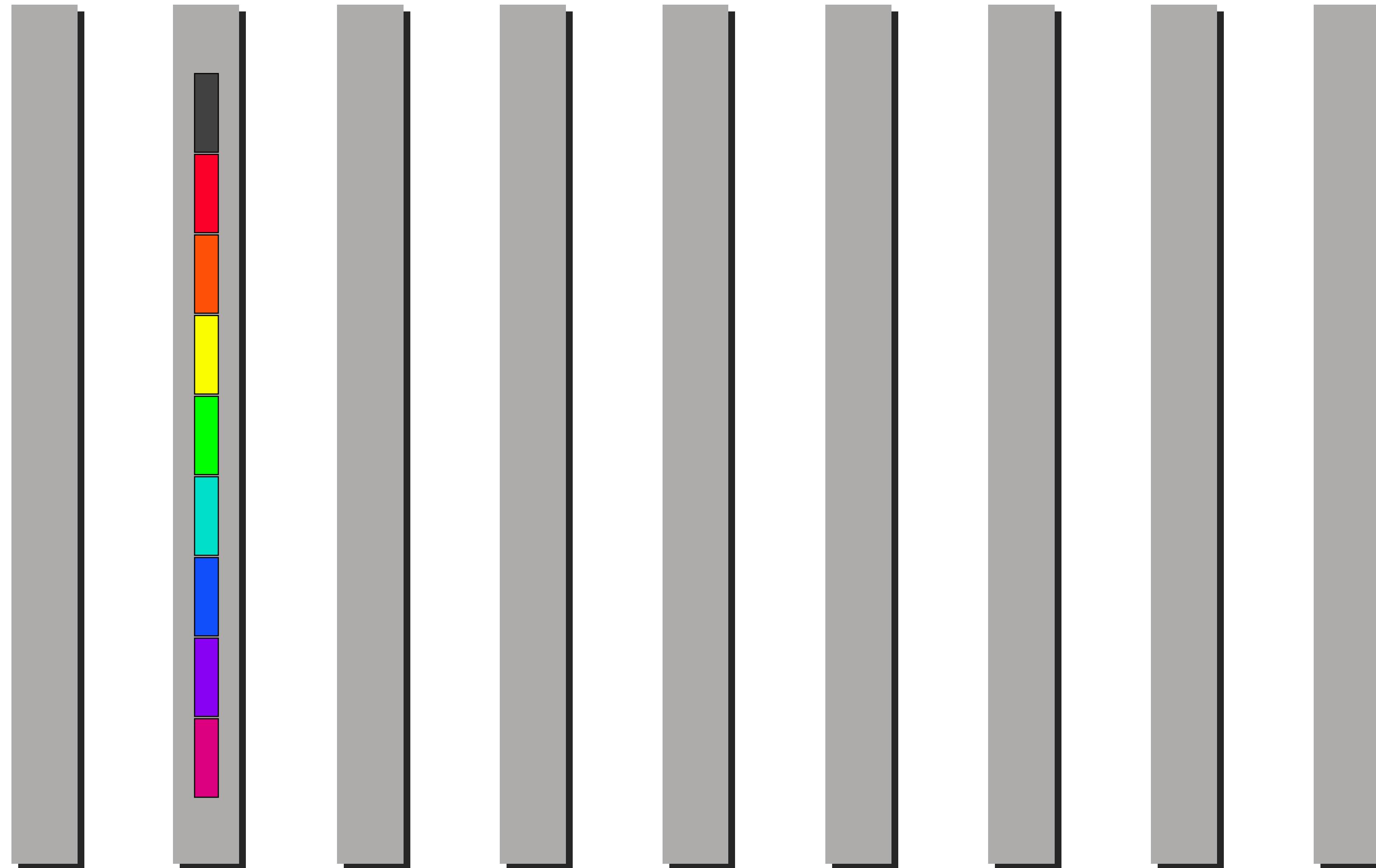
Scatter

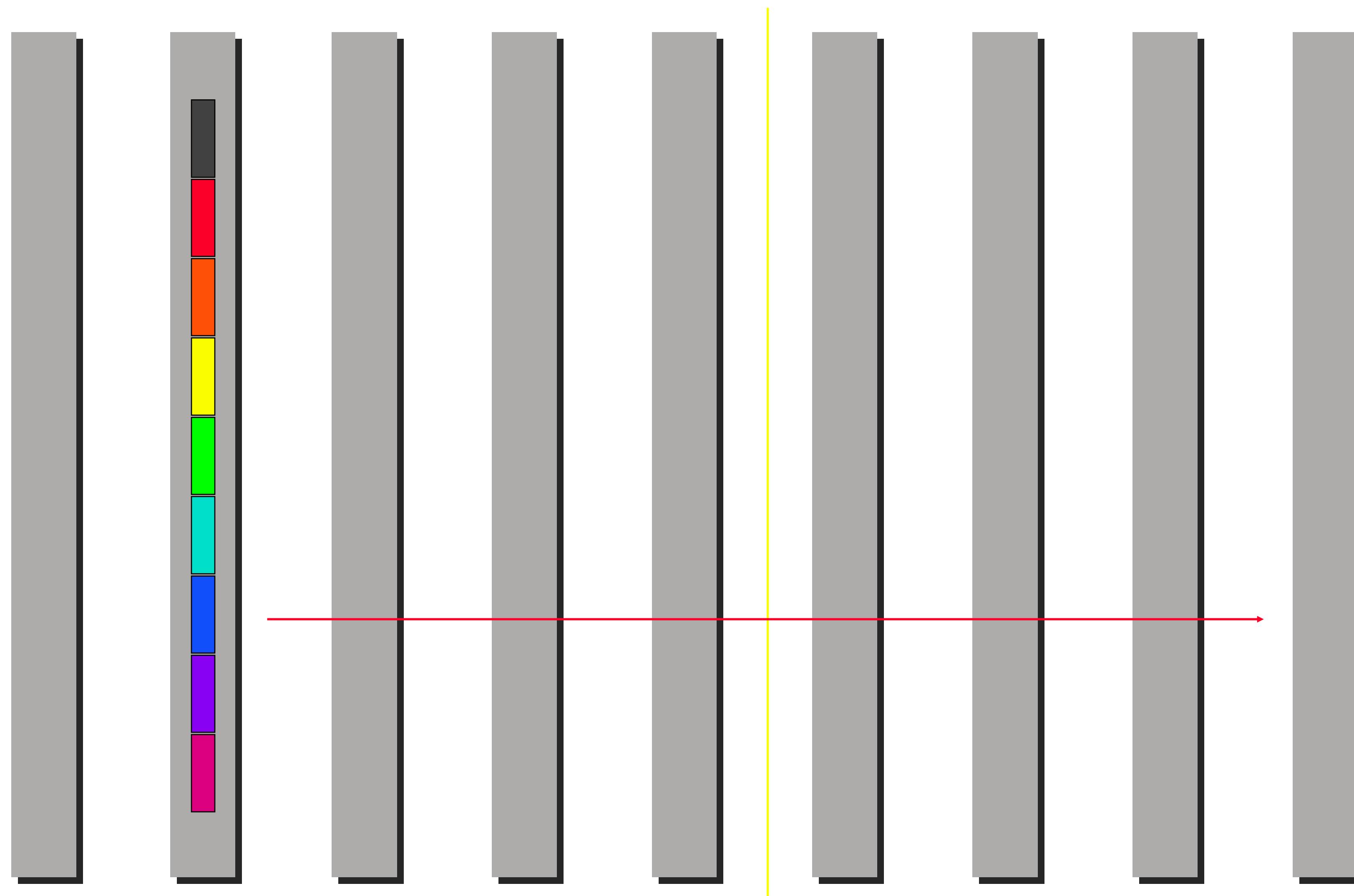
Before

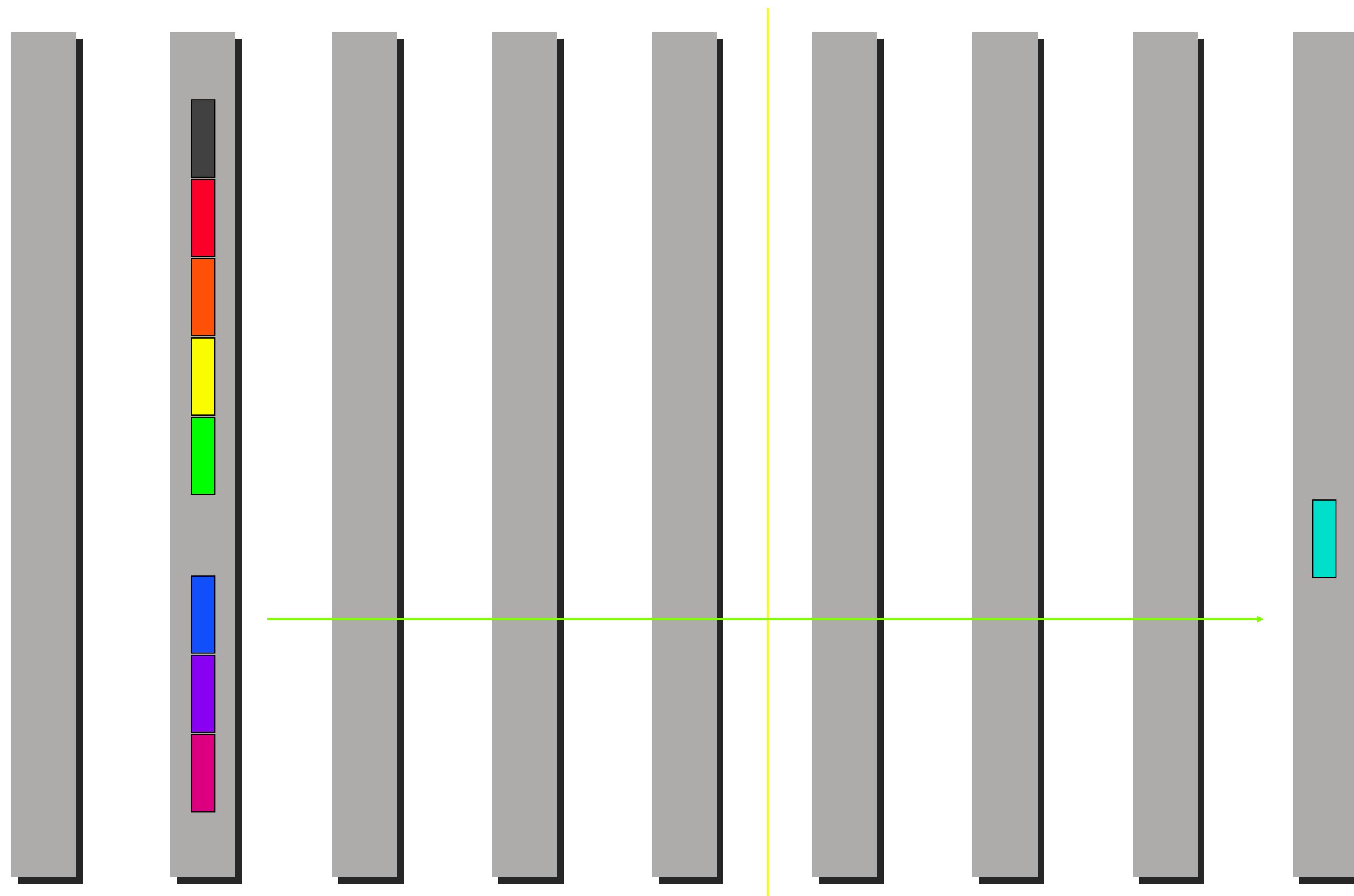


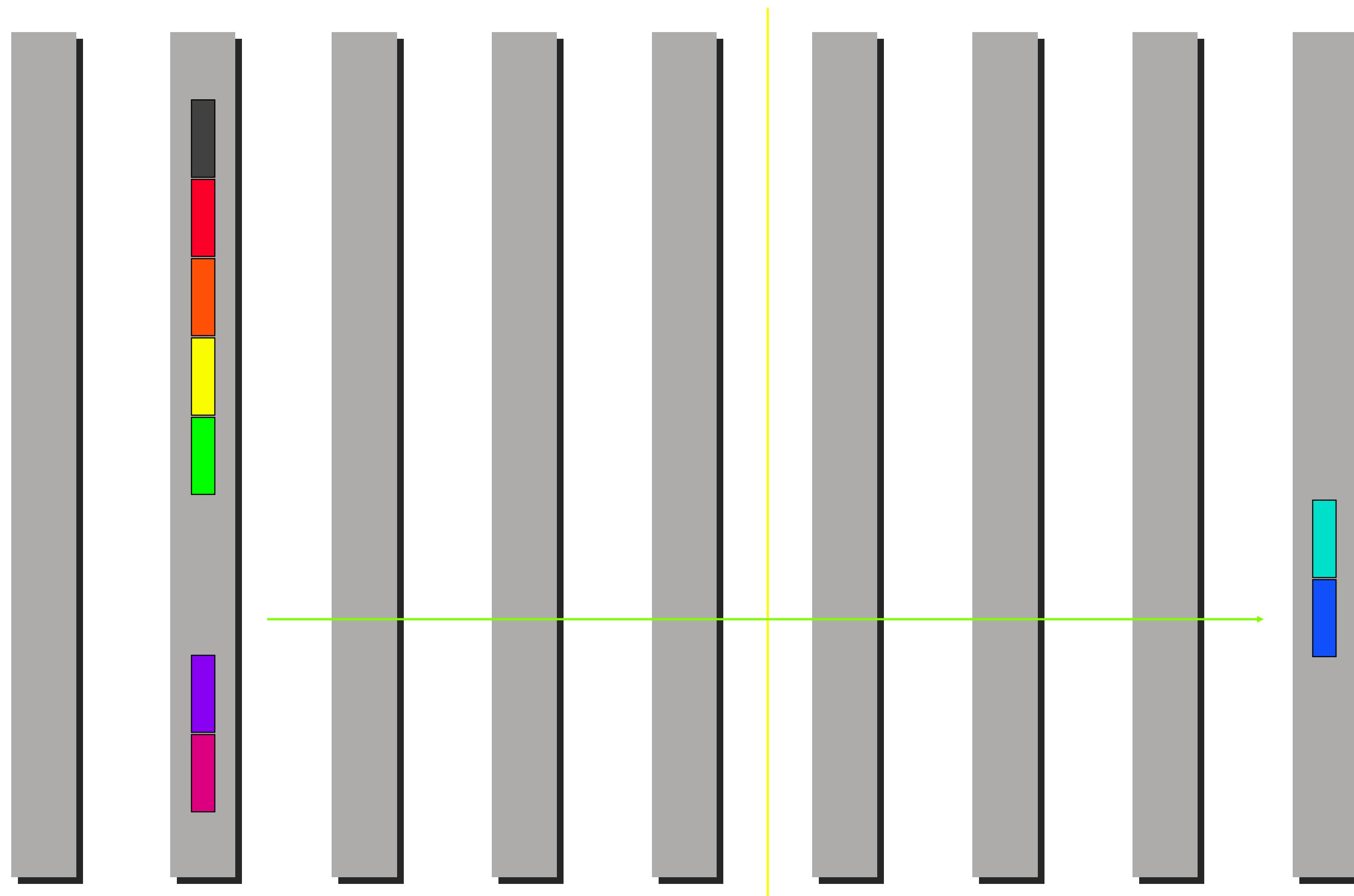
After

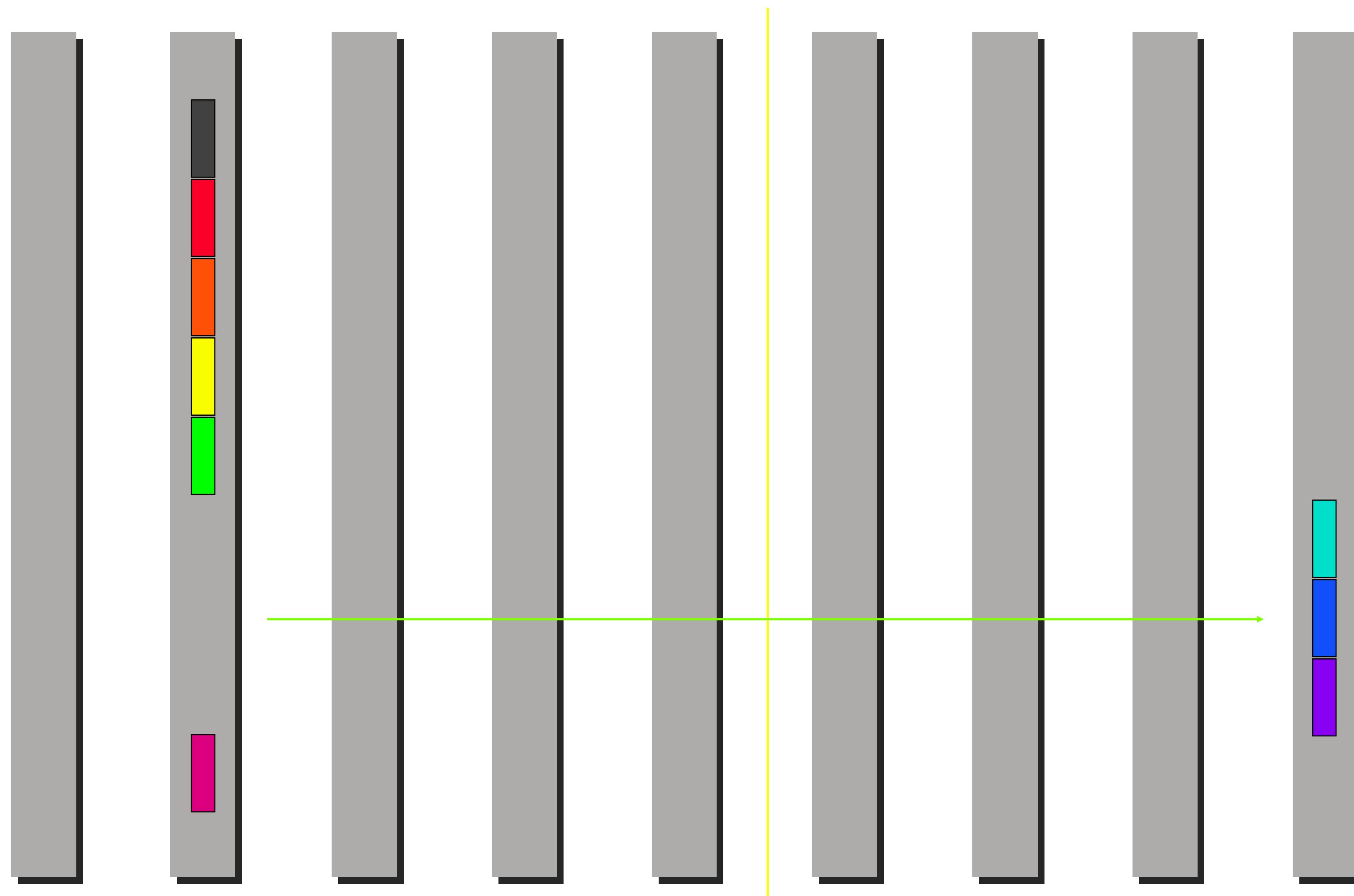


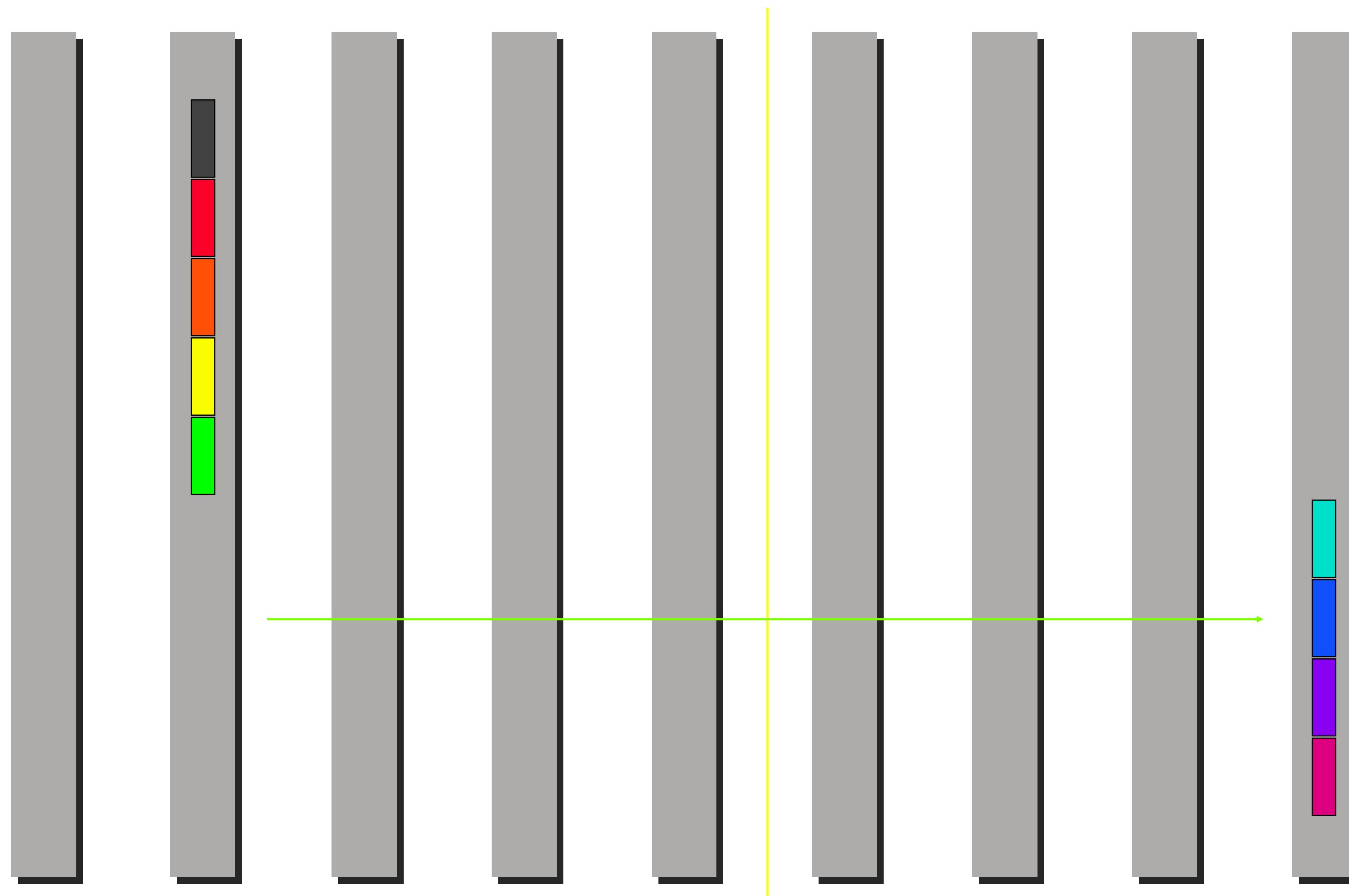


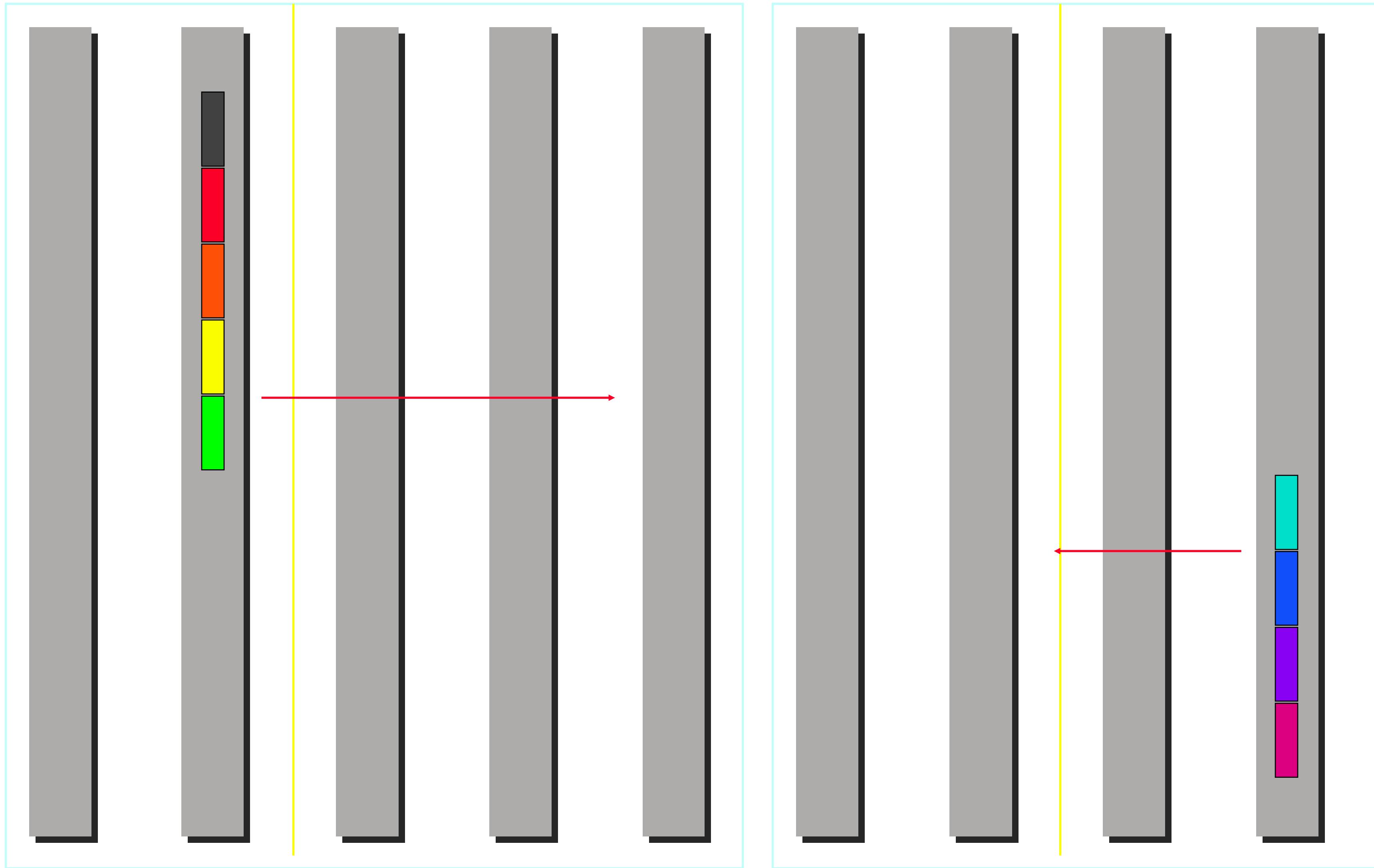


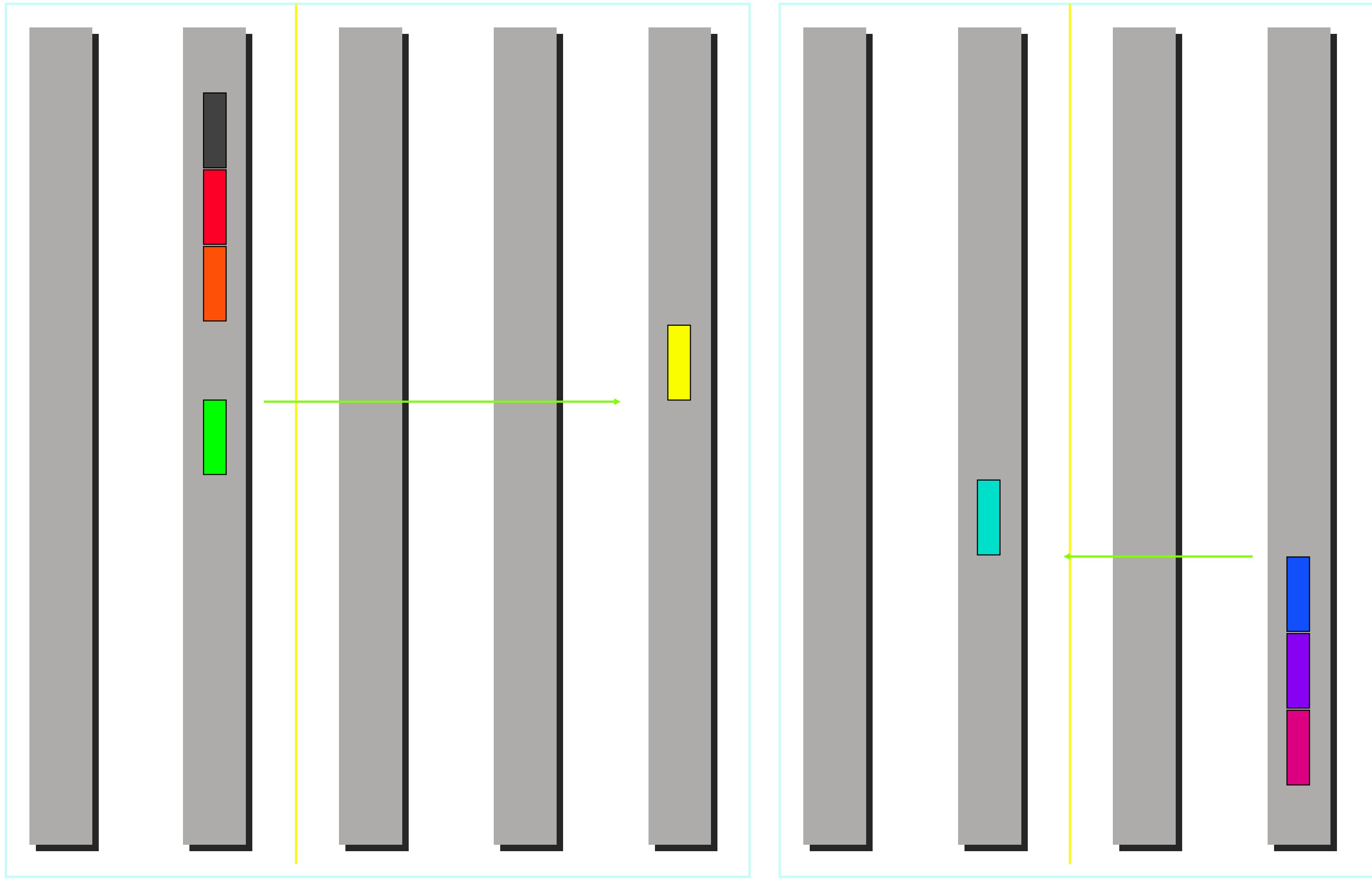


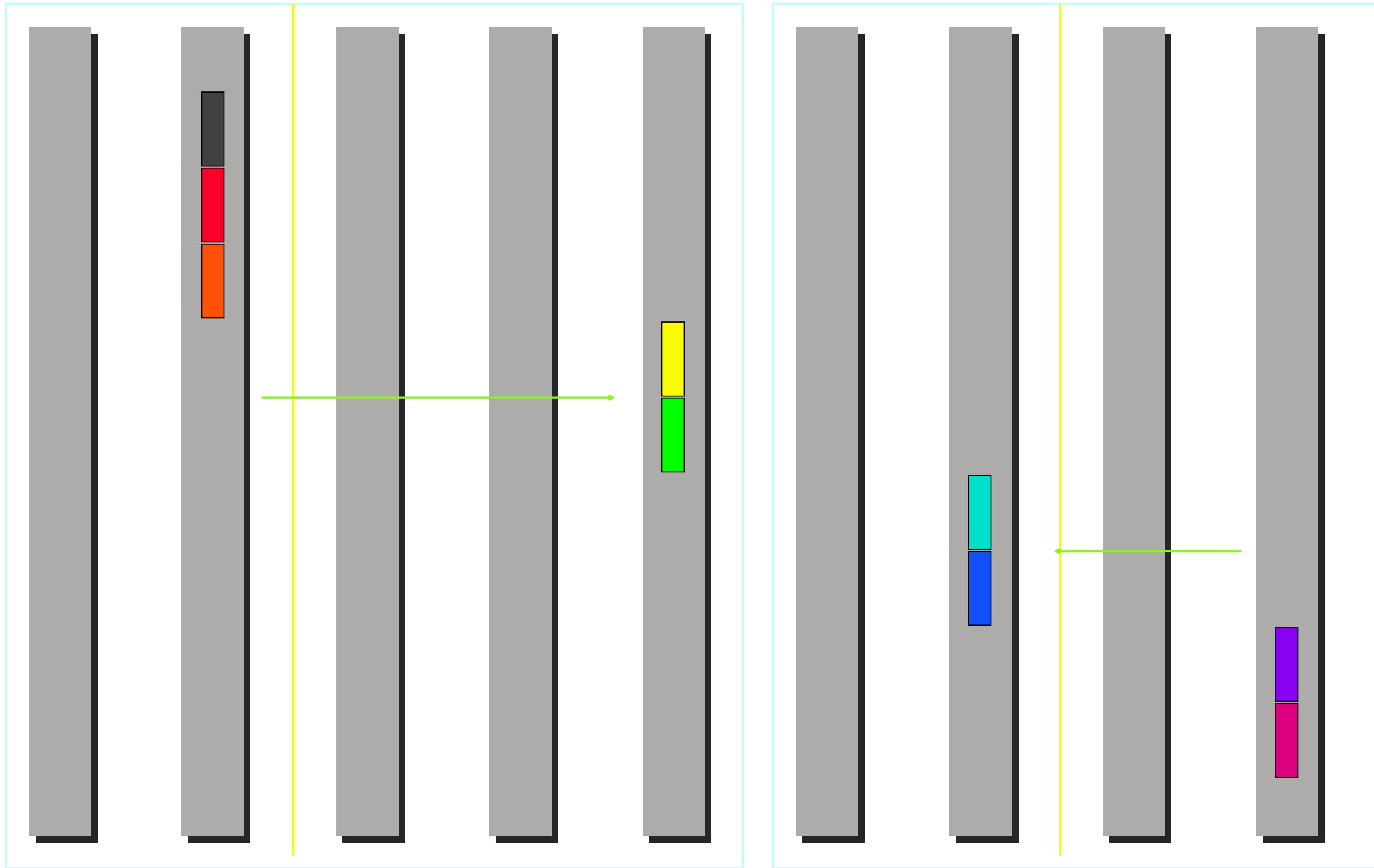


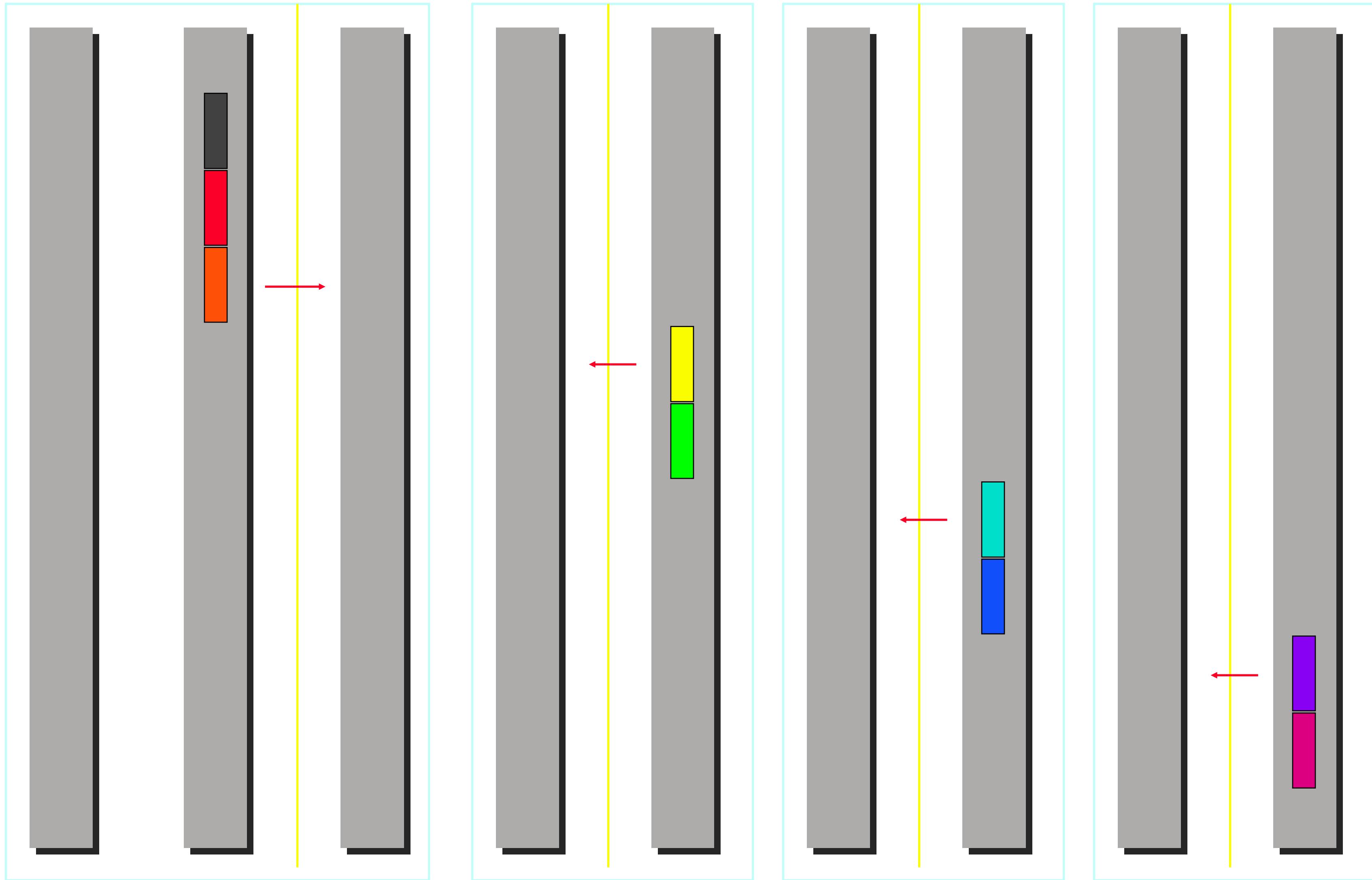


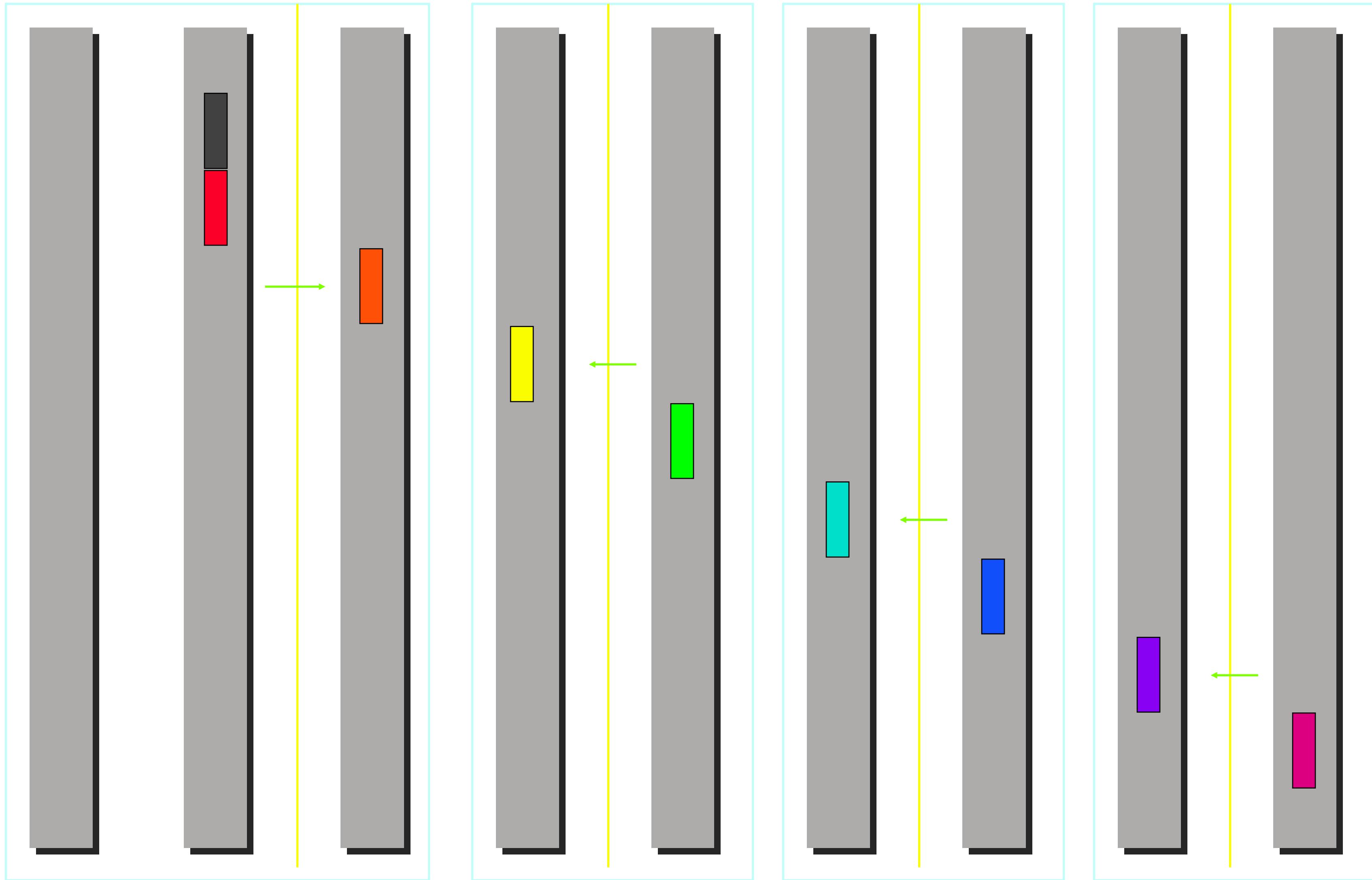


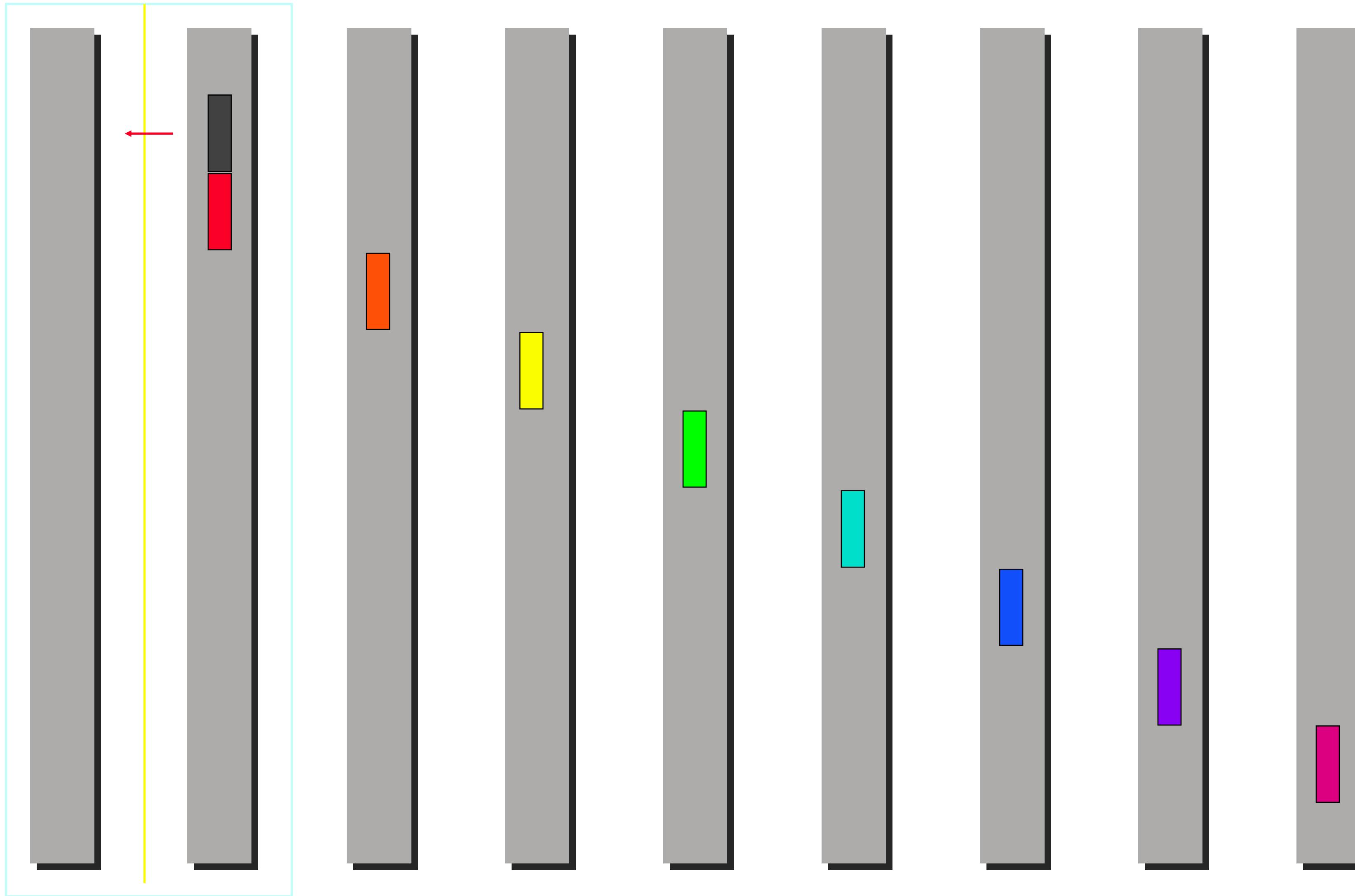


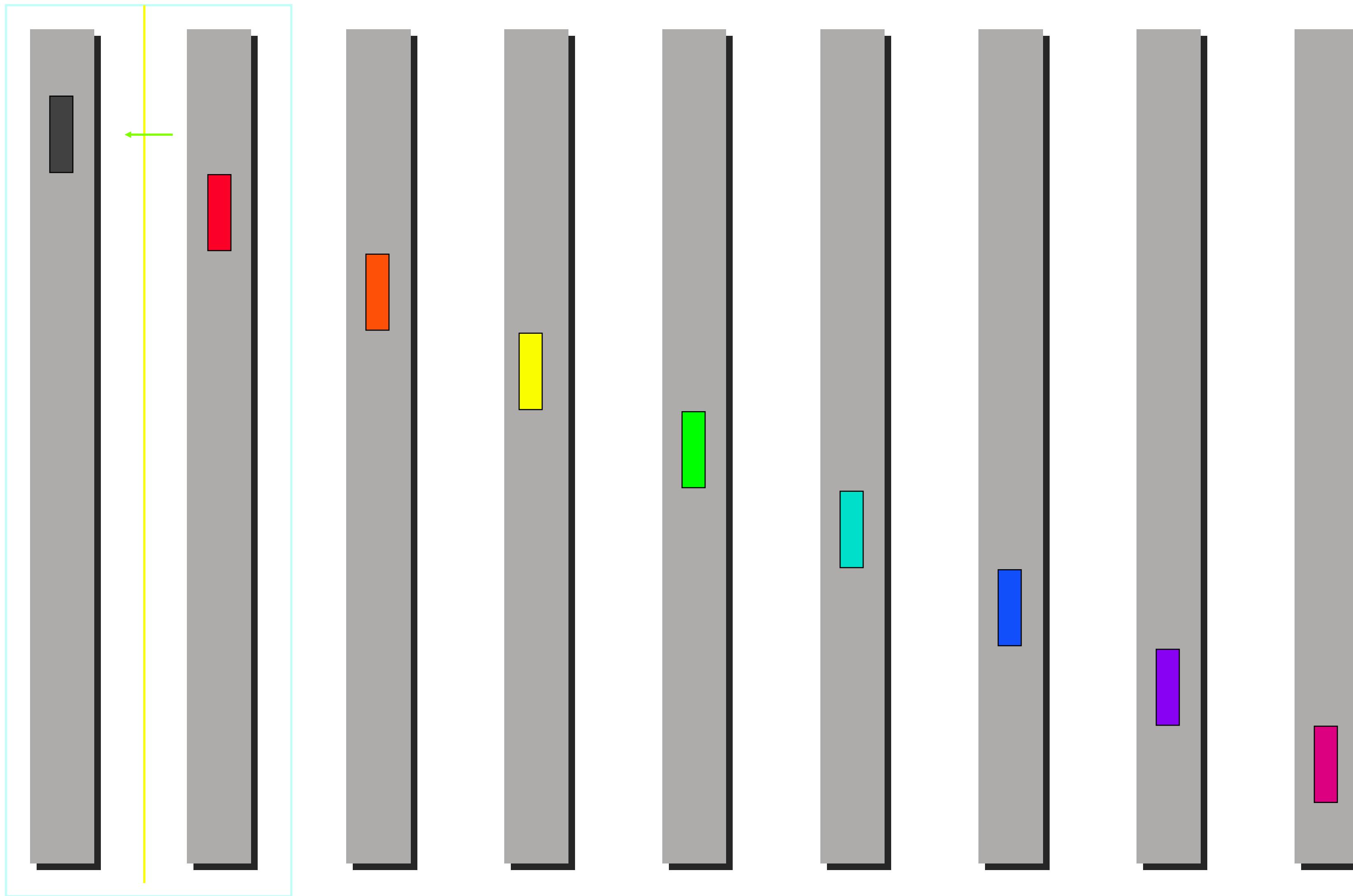


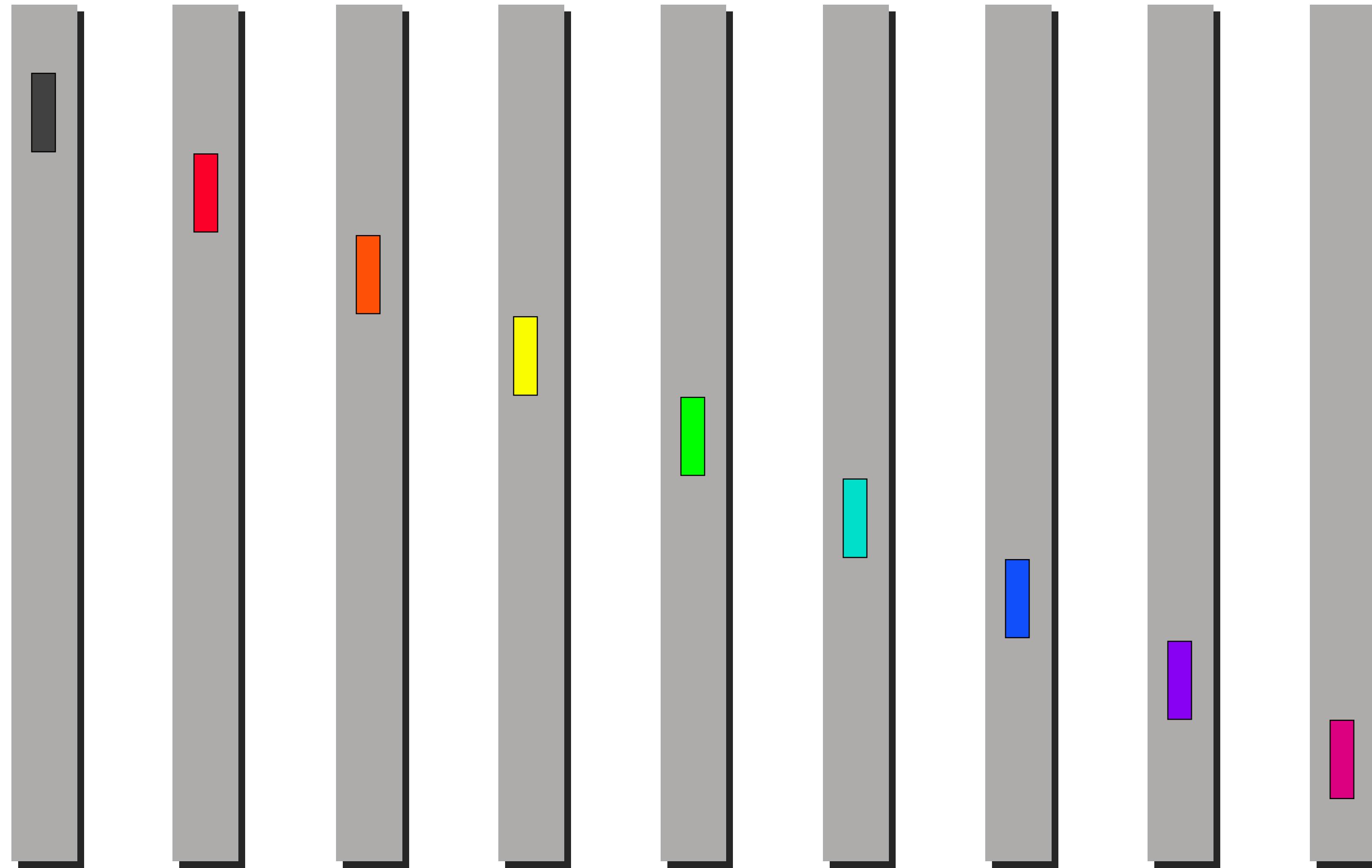












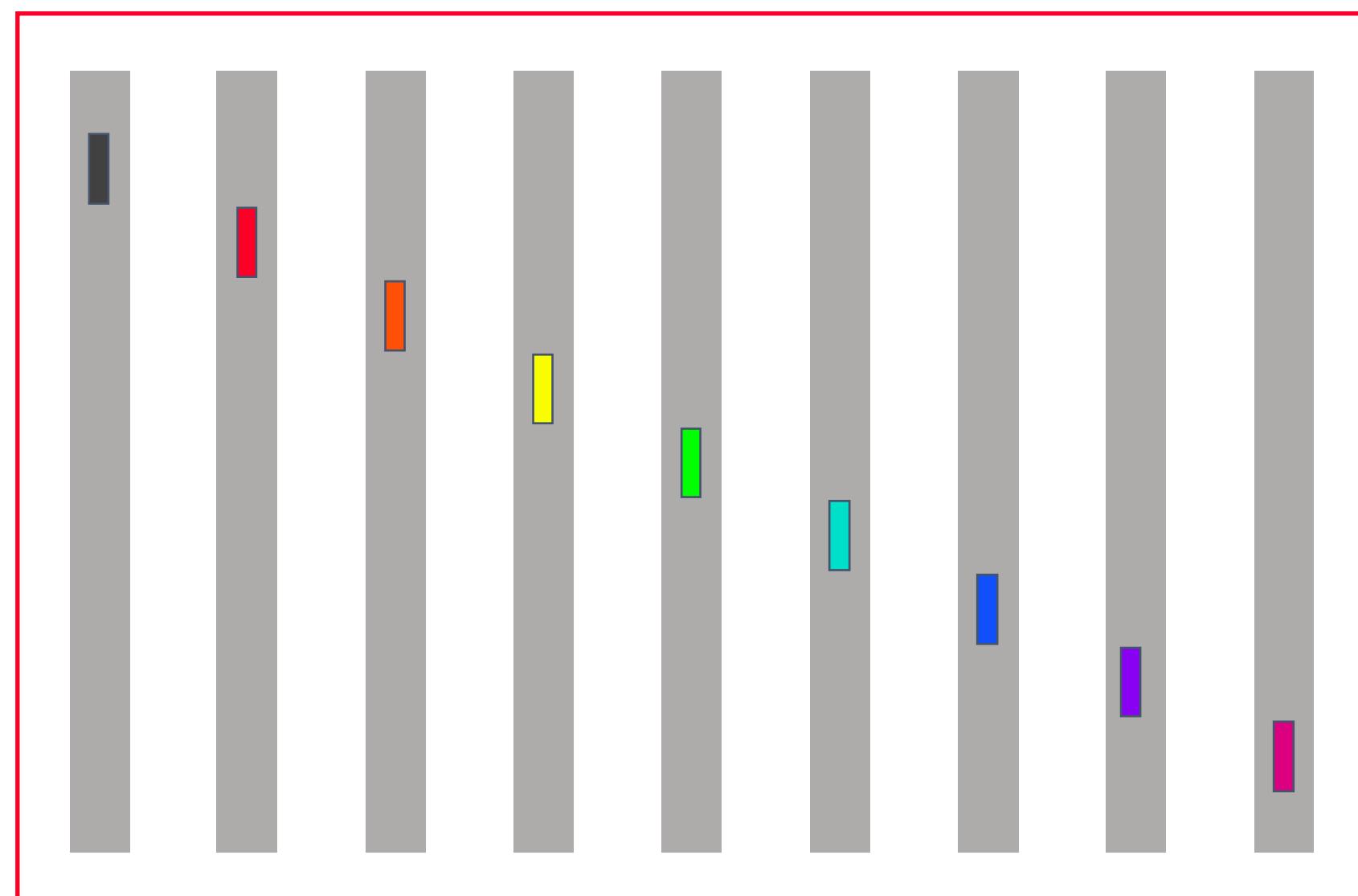
Cost of minimum spanning tree scatter

- Assumption: power of two number of nodes

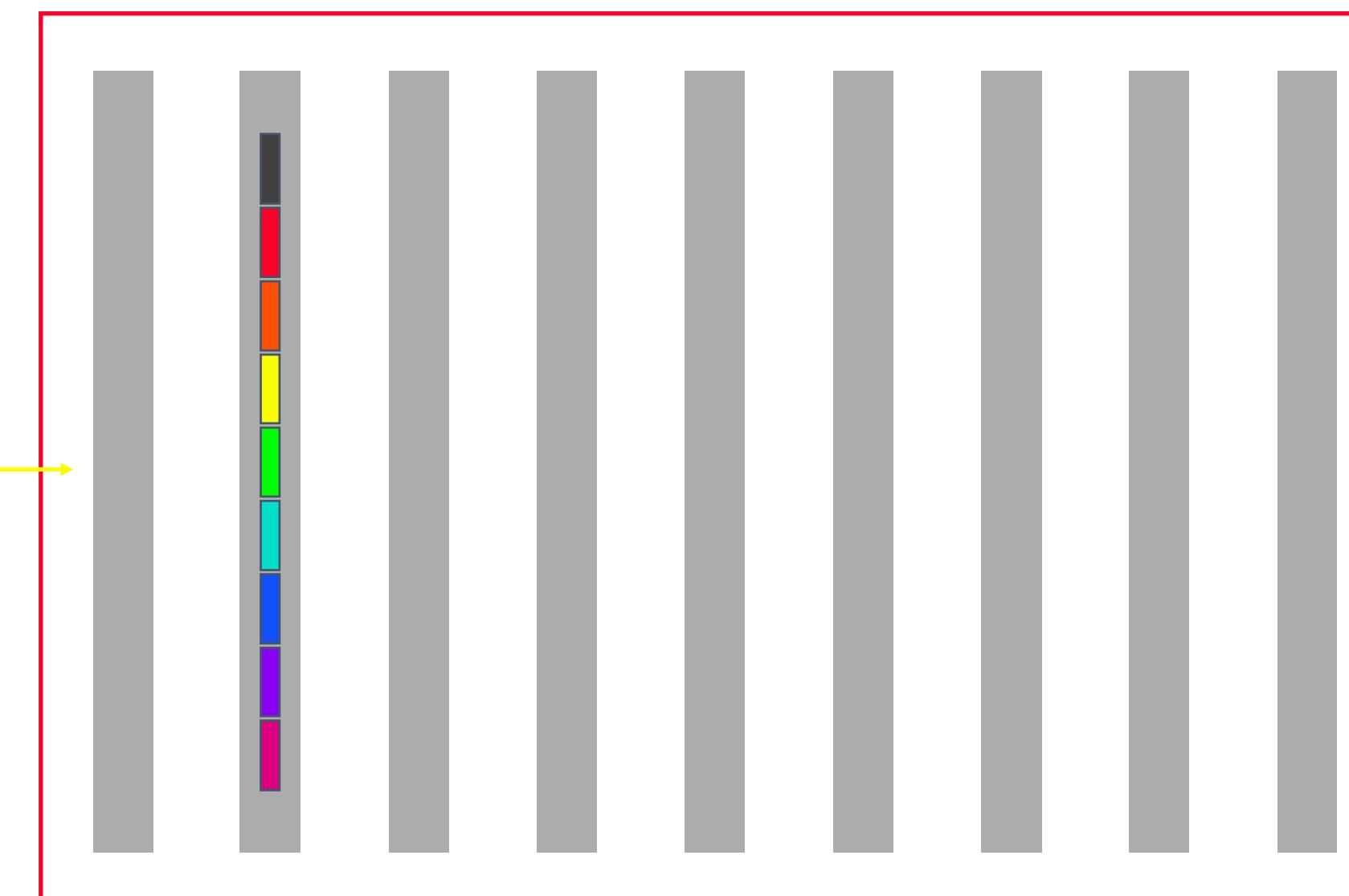
$$\begin{aligned} & \log(p) \sum_{k=1}^{\log(p)} \left(\alpha + \frac{n}{2^k} \beta \right) \\ & = \\ & \log(p) \quad \alpha + \frac{p-1}{p} n \beta \end{aligned}$$

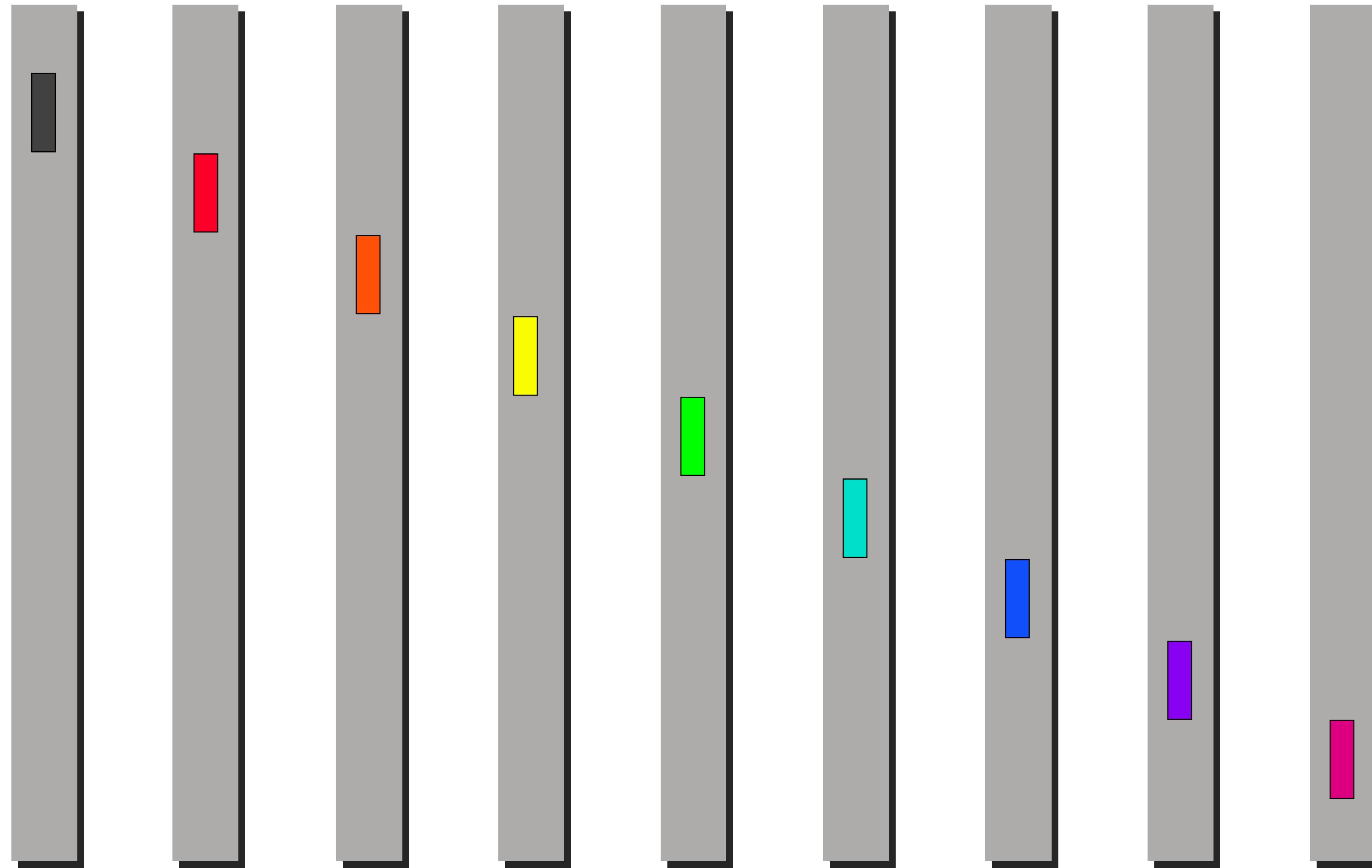
Gather

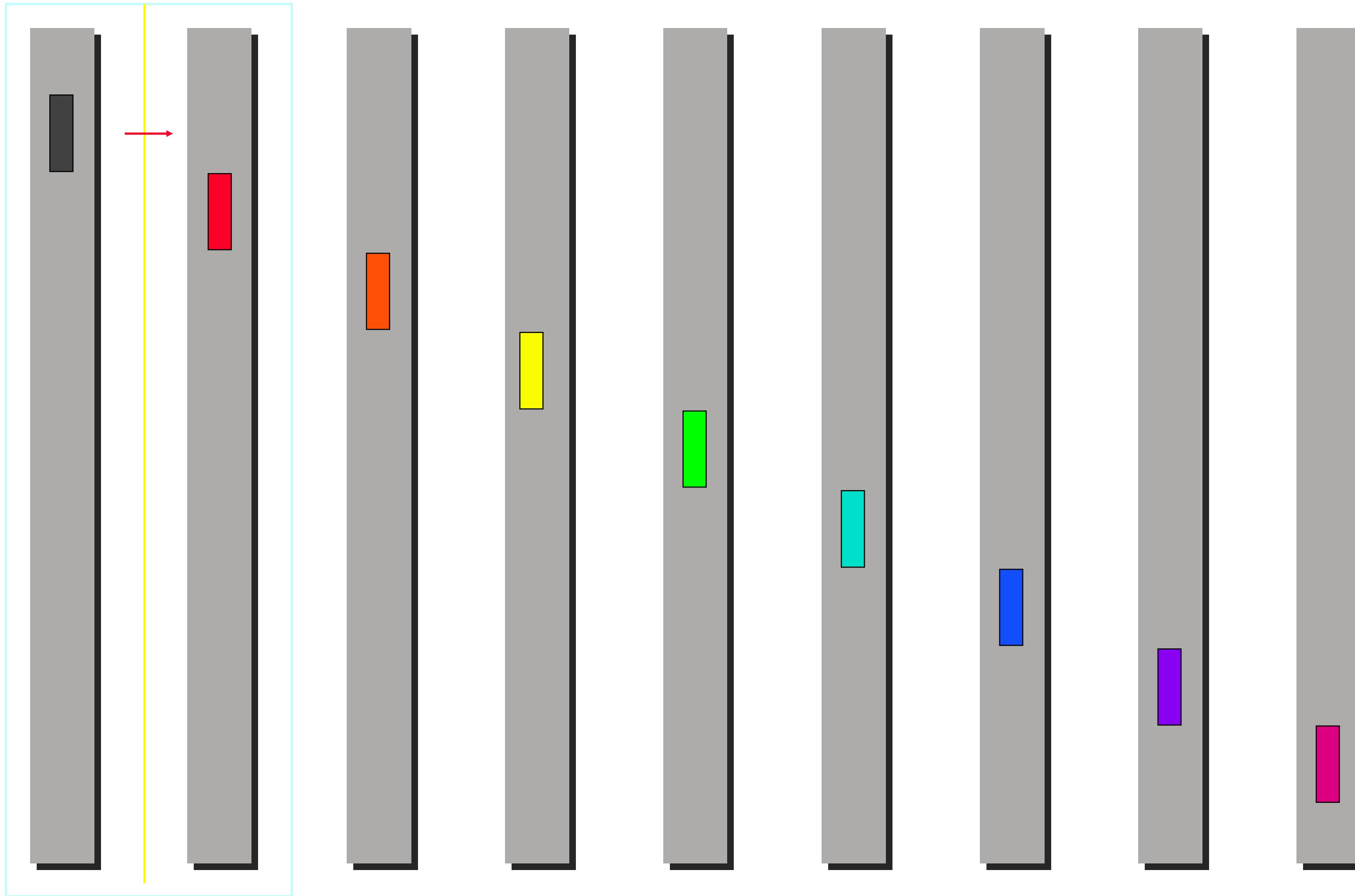
Before

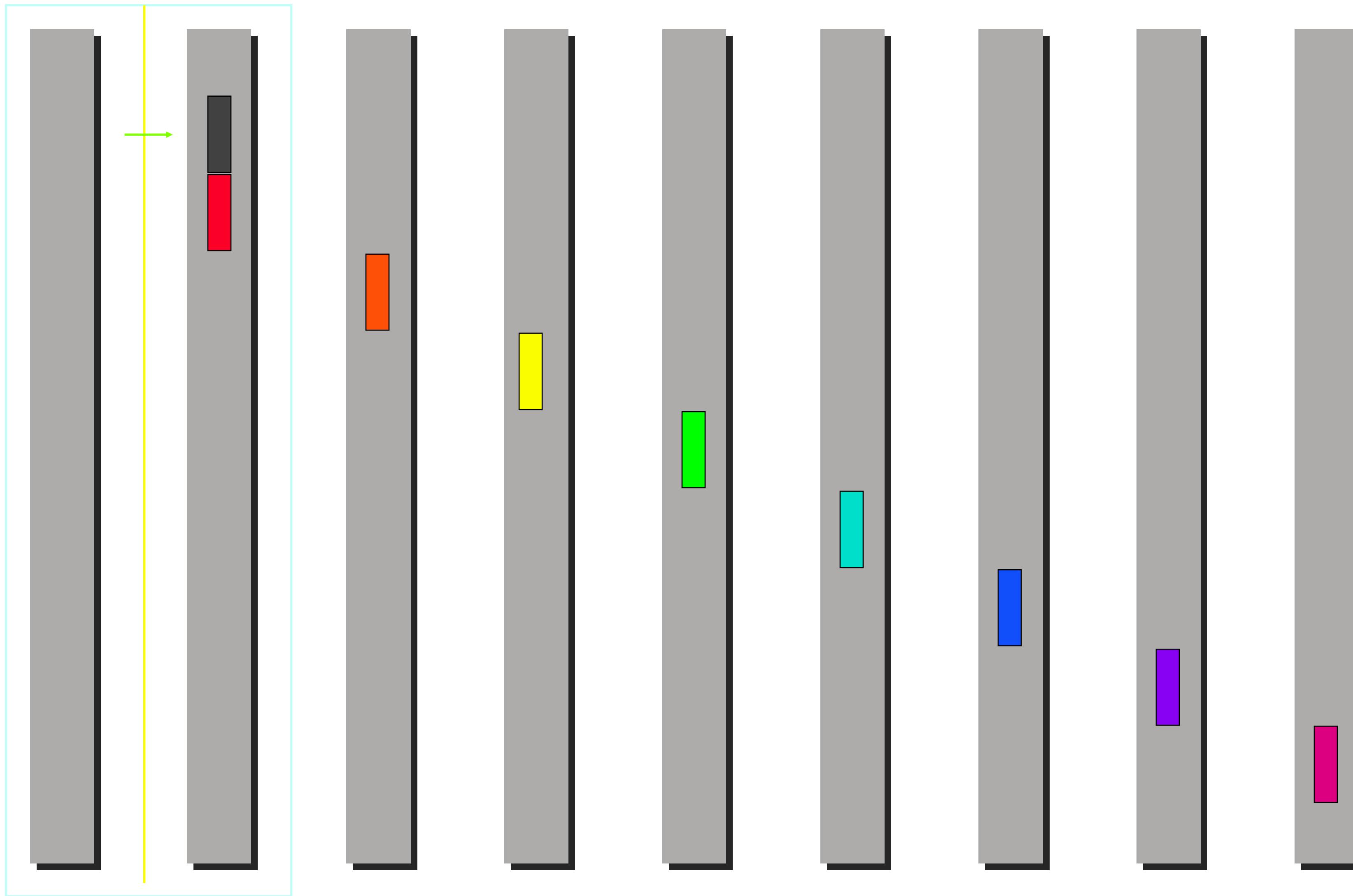


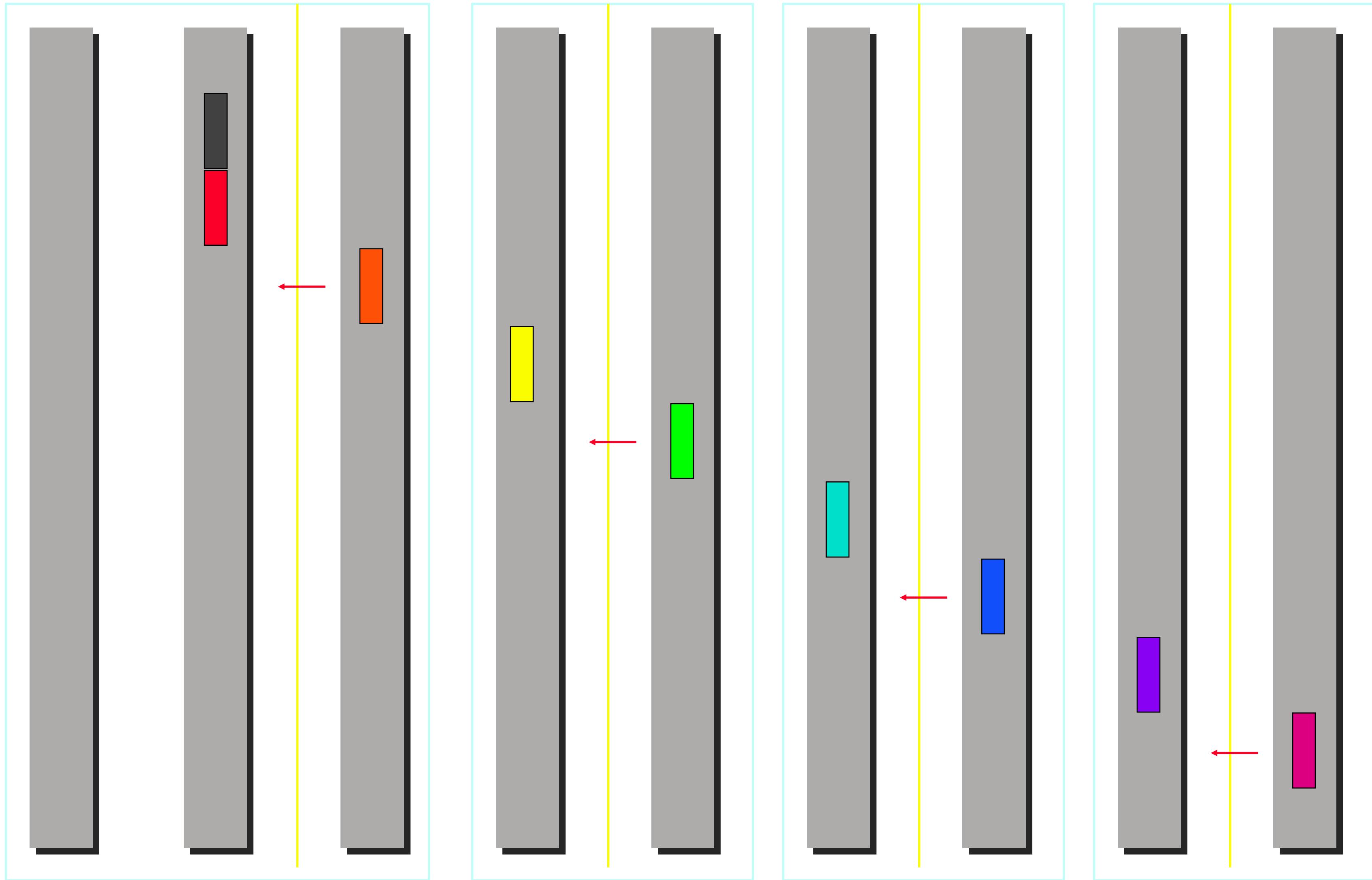
After

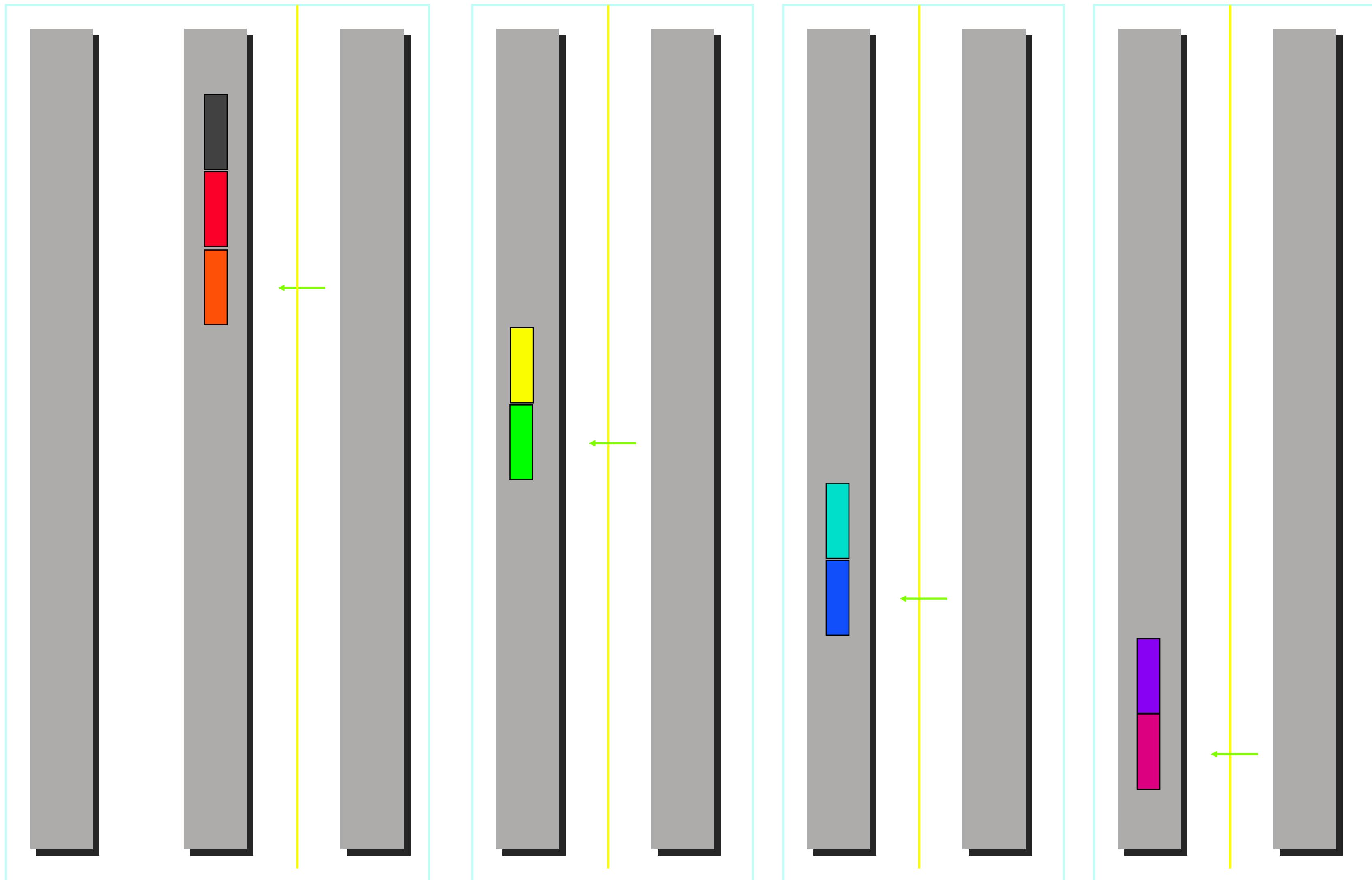


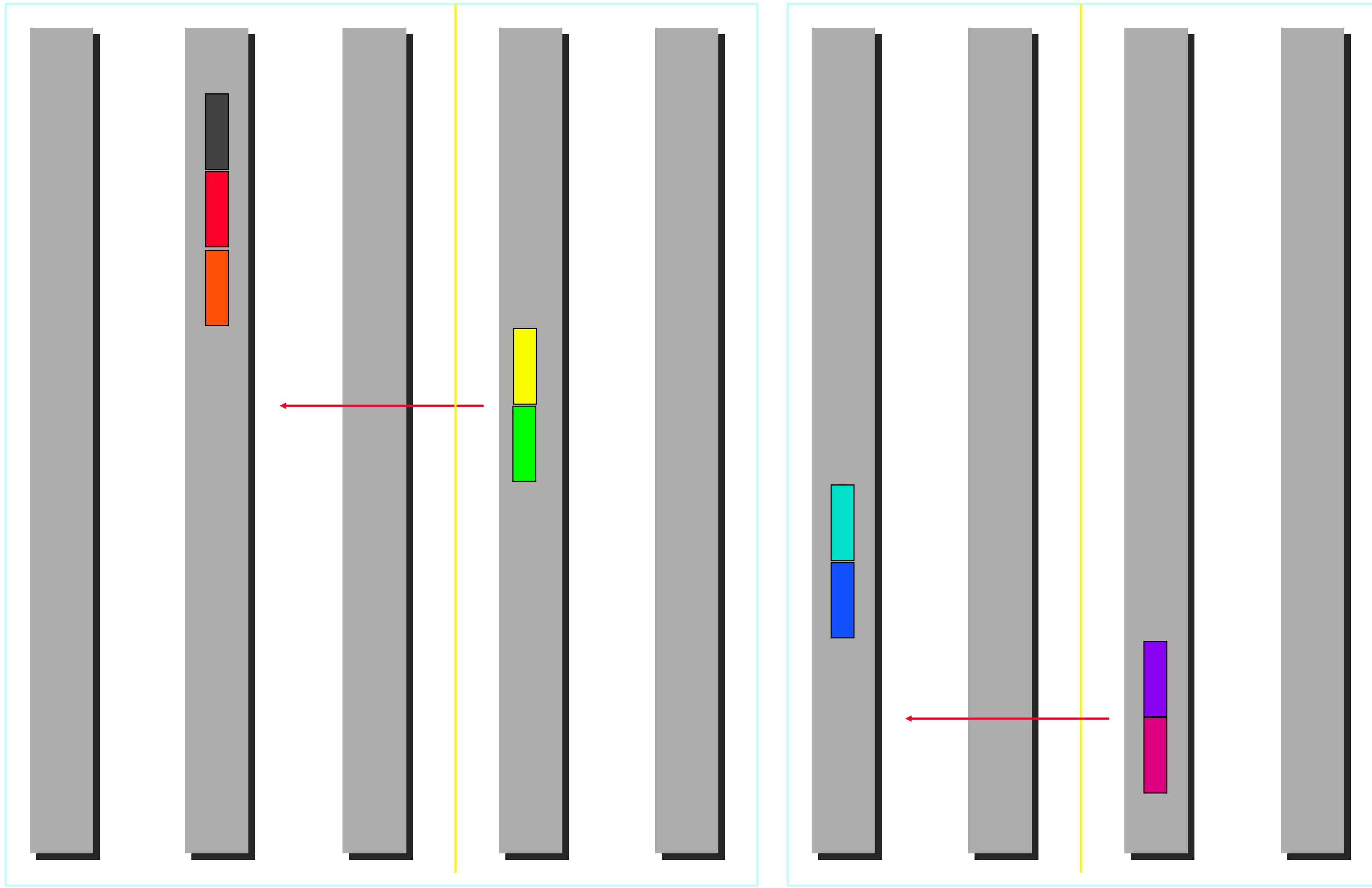


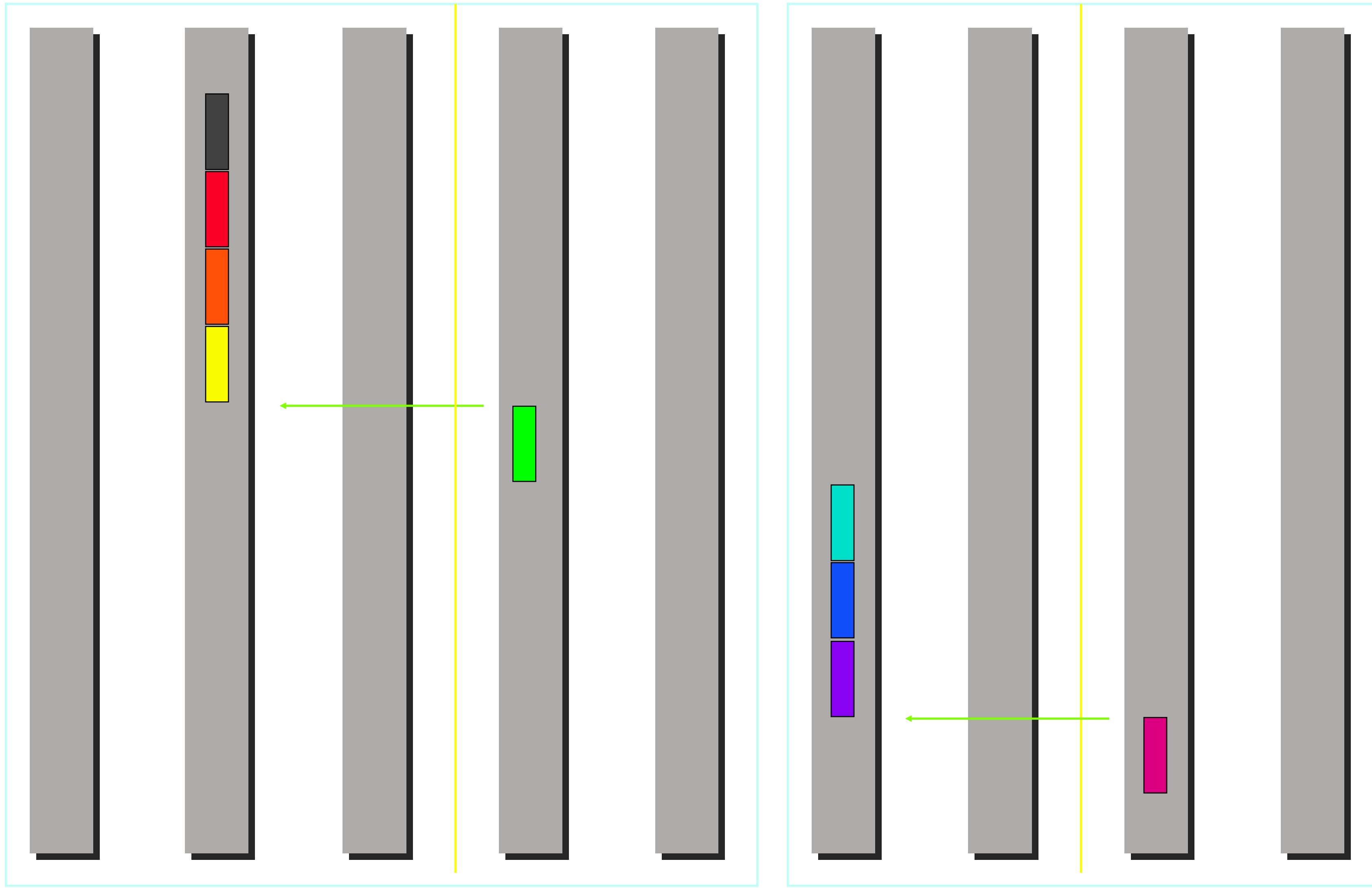


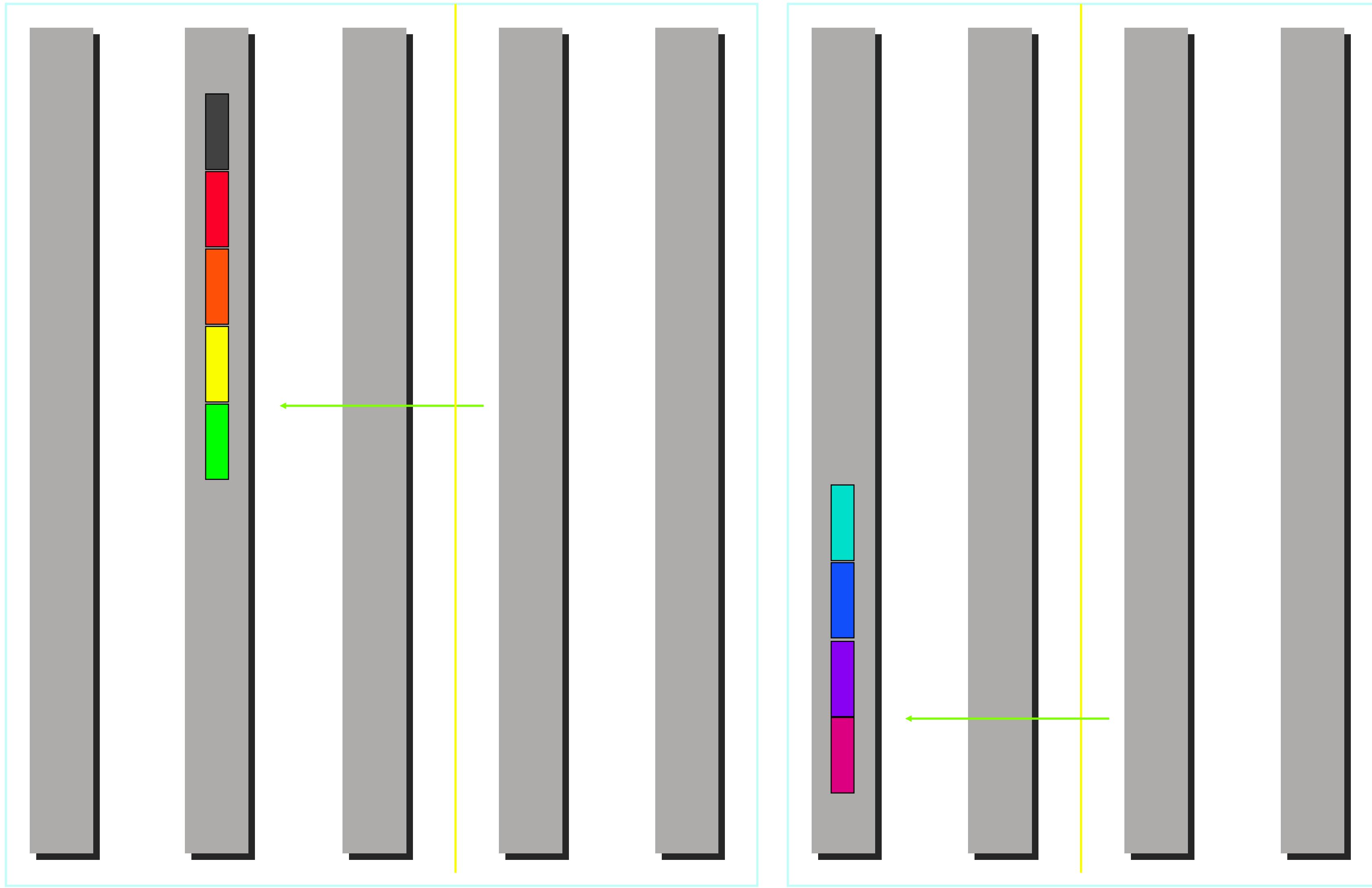


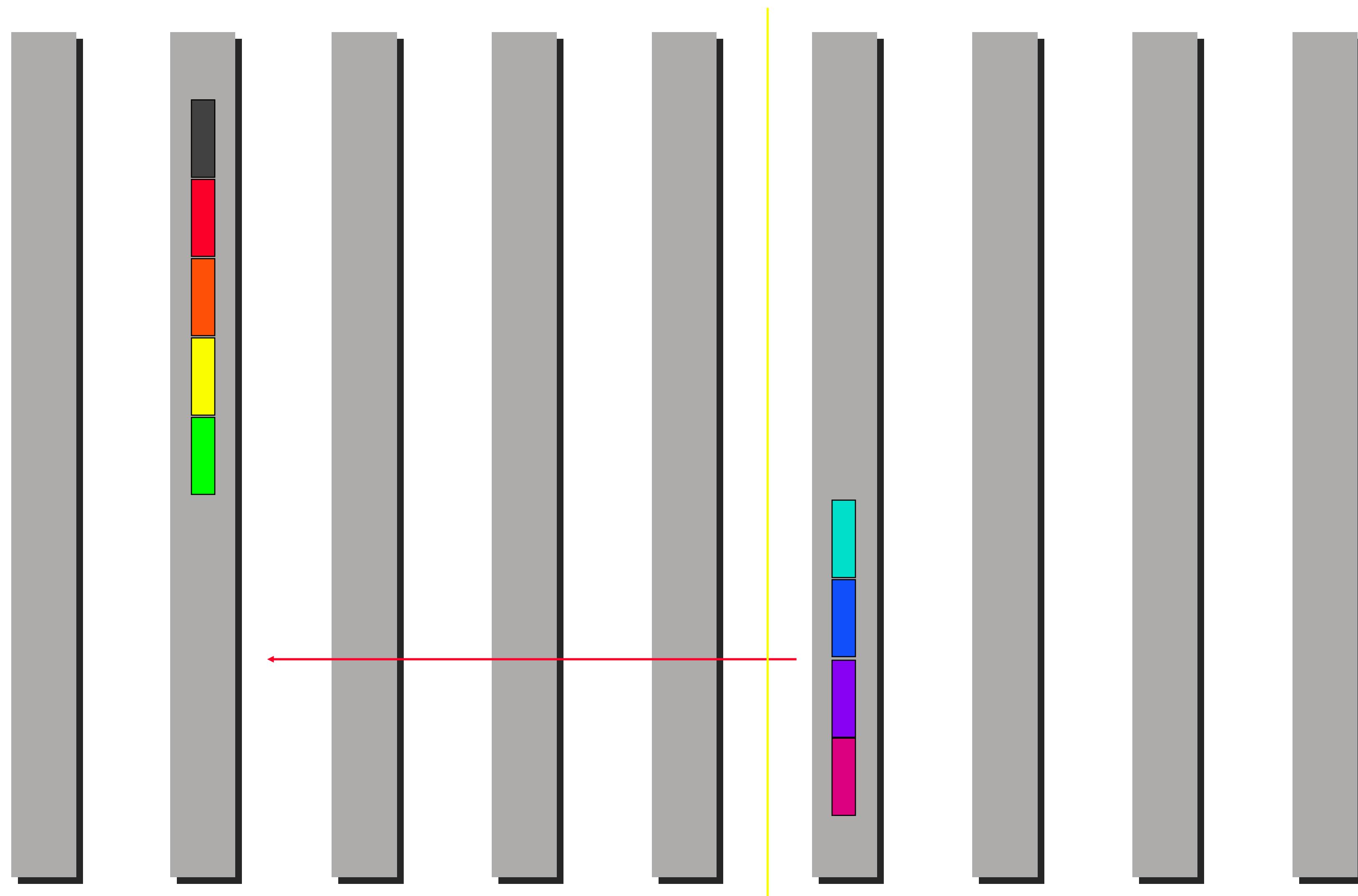


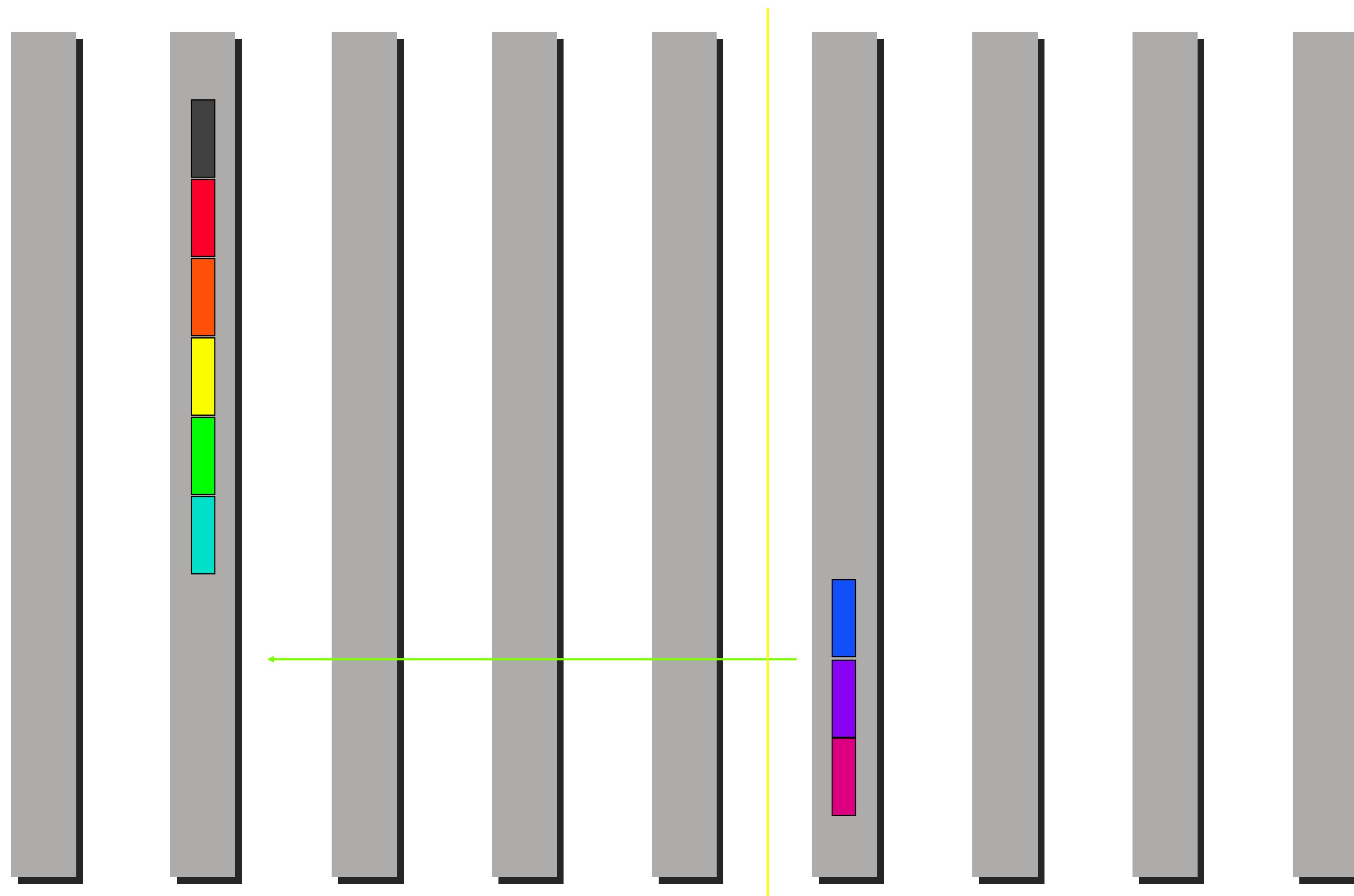


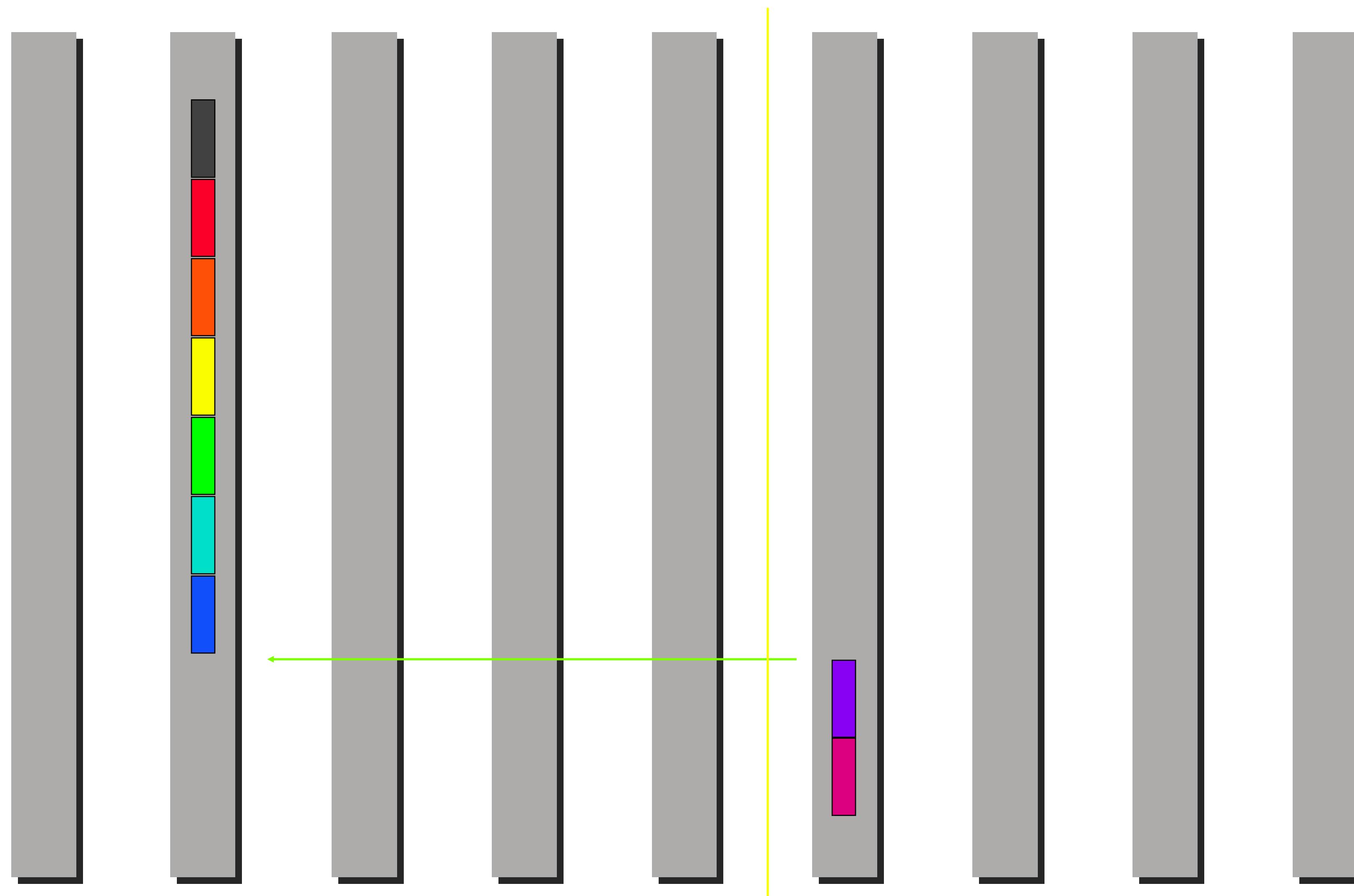


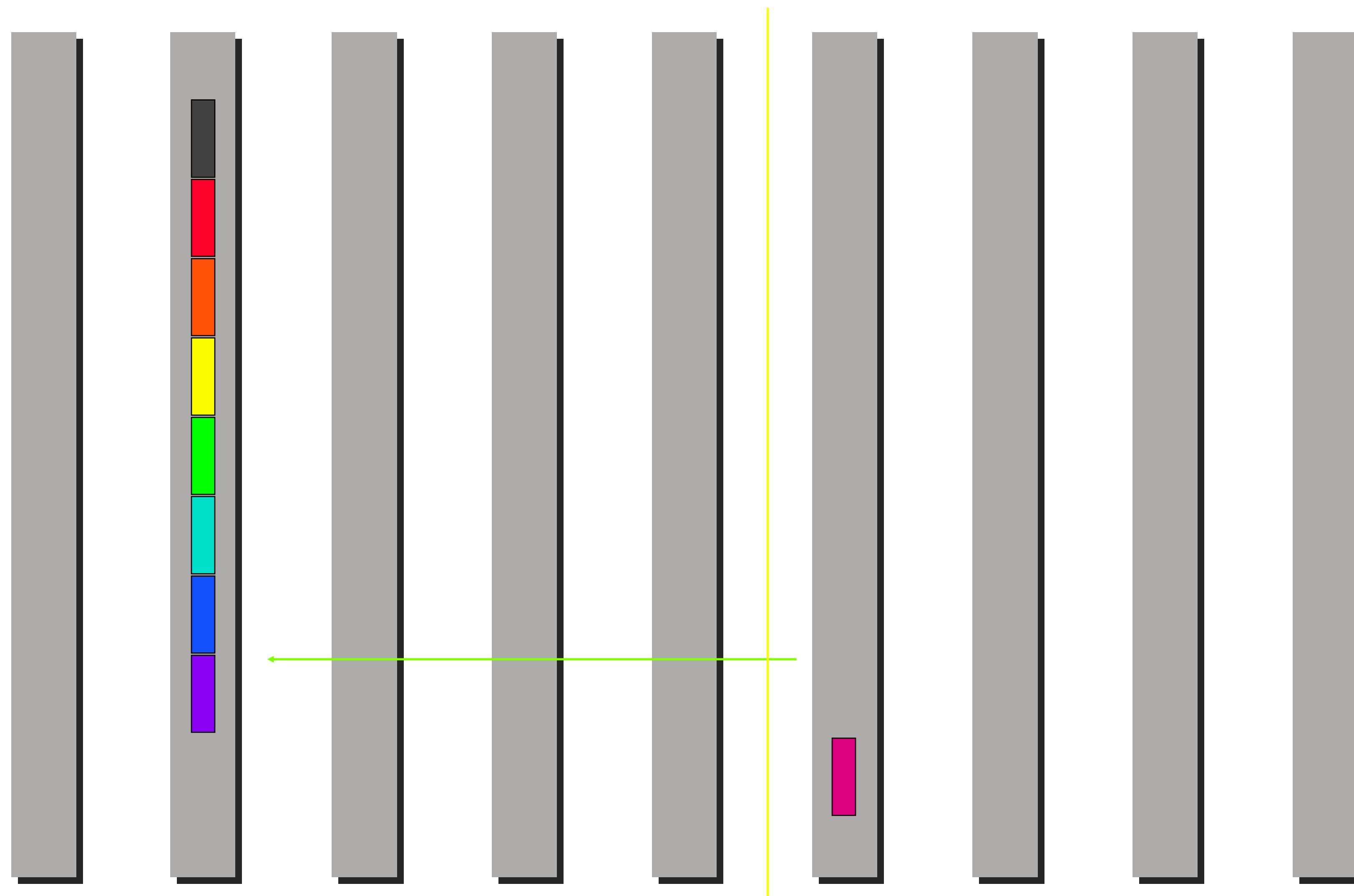


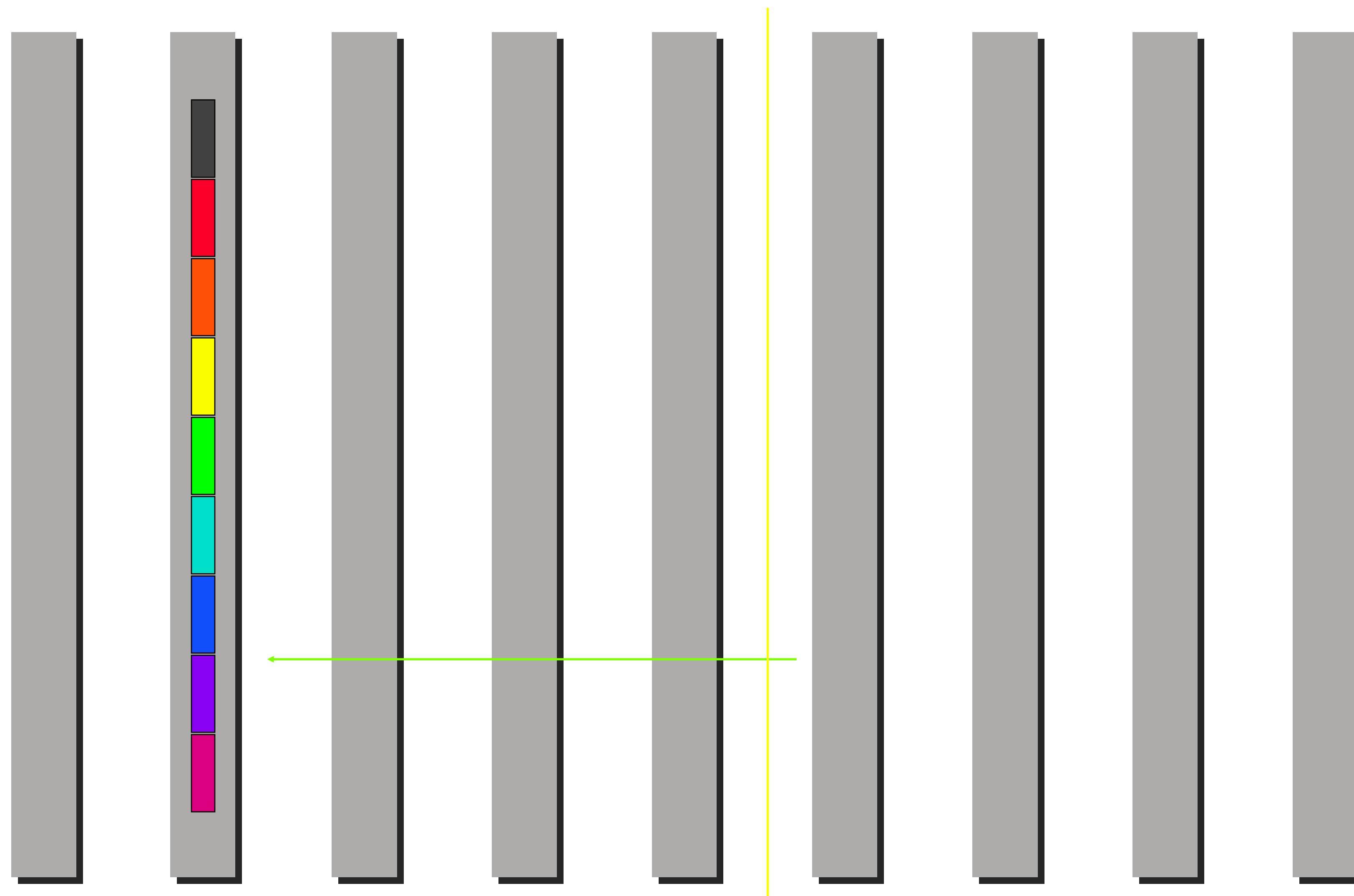


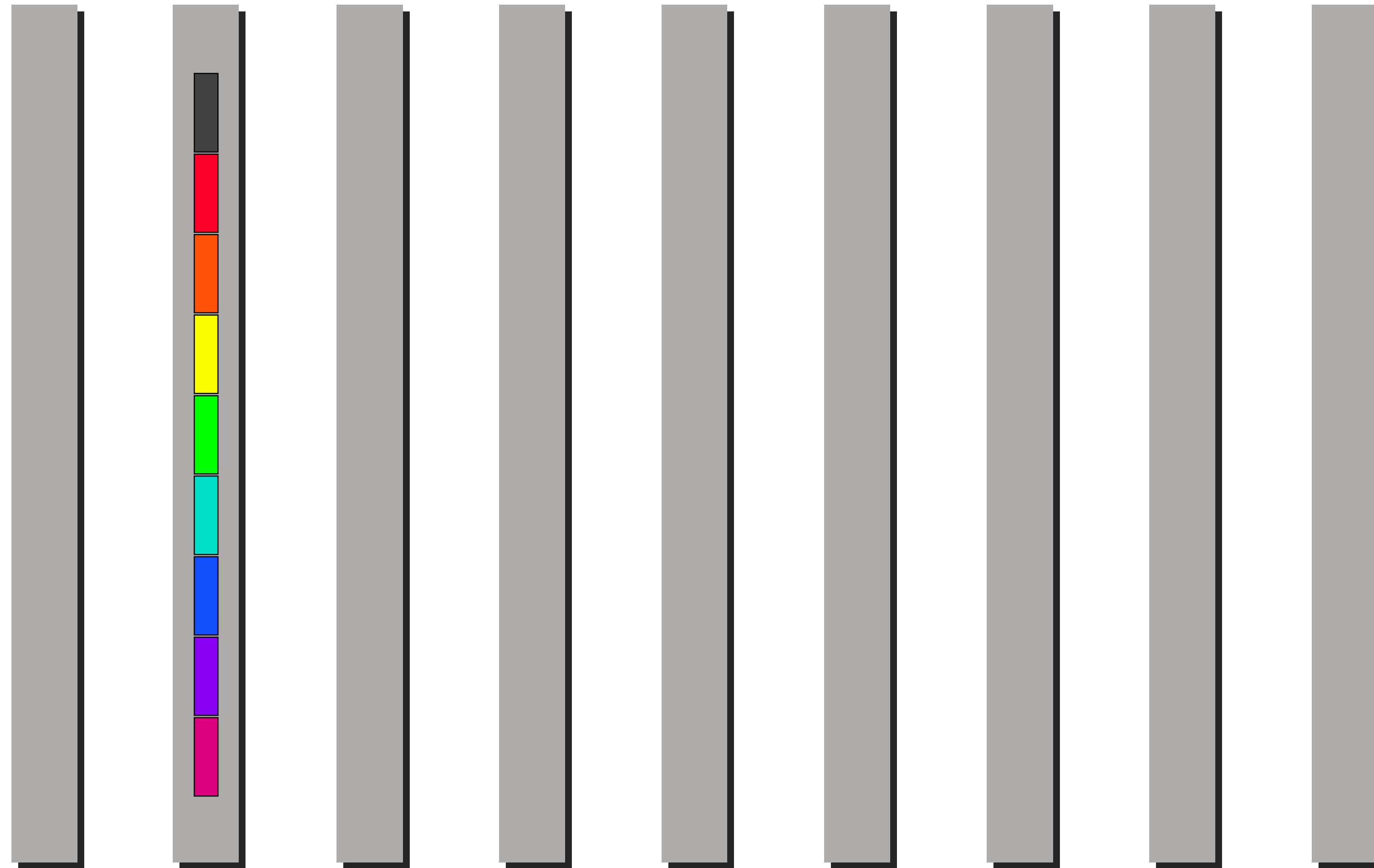












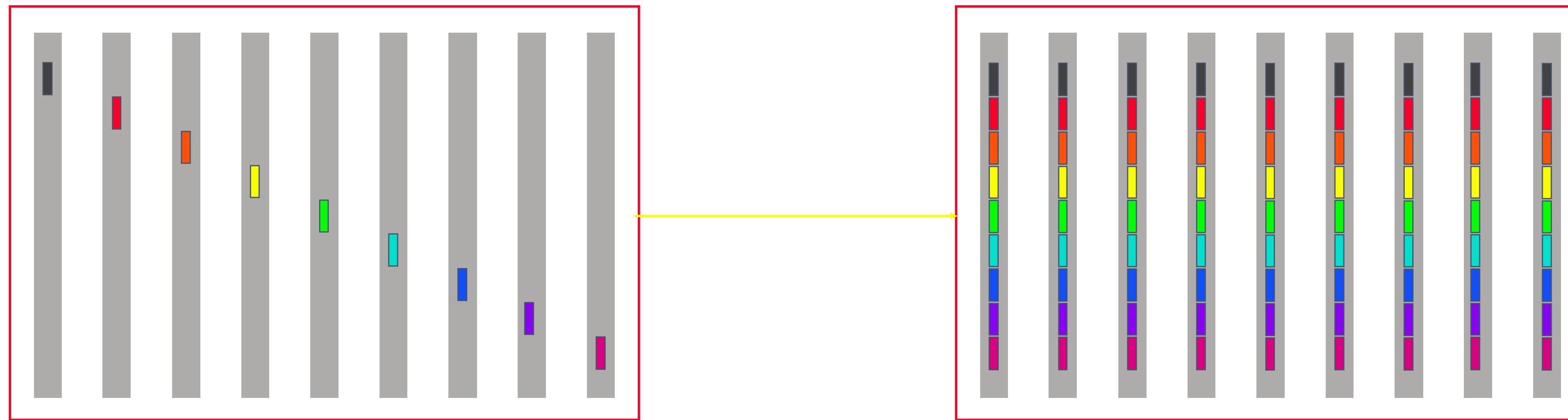
Cost of minimum spanning tree gather

- Assumption: power of two number of nodes

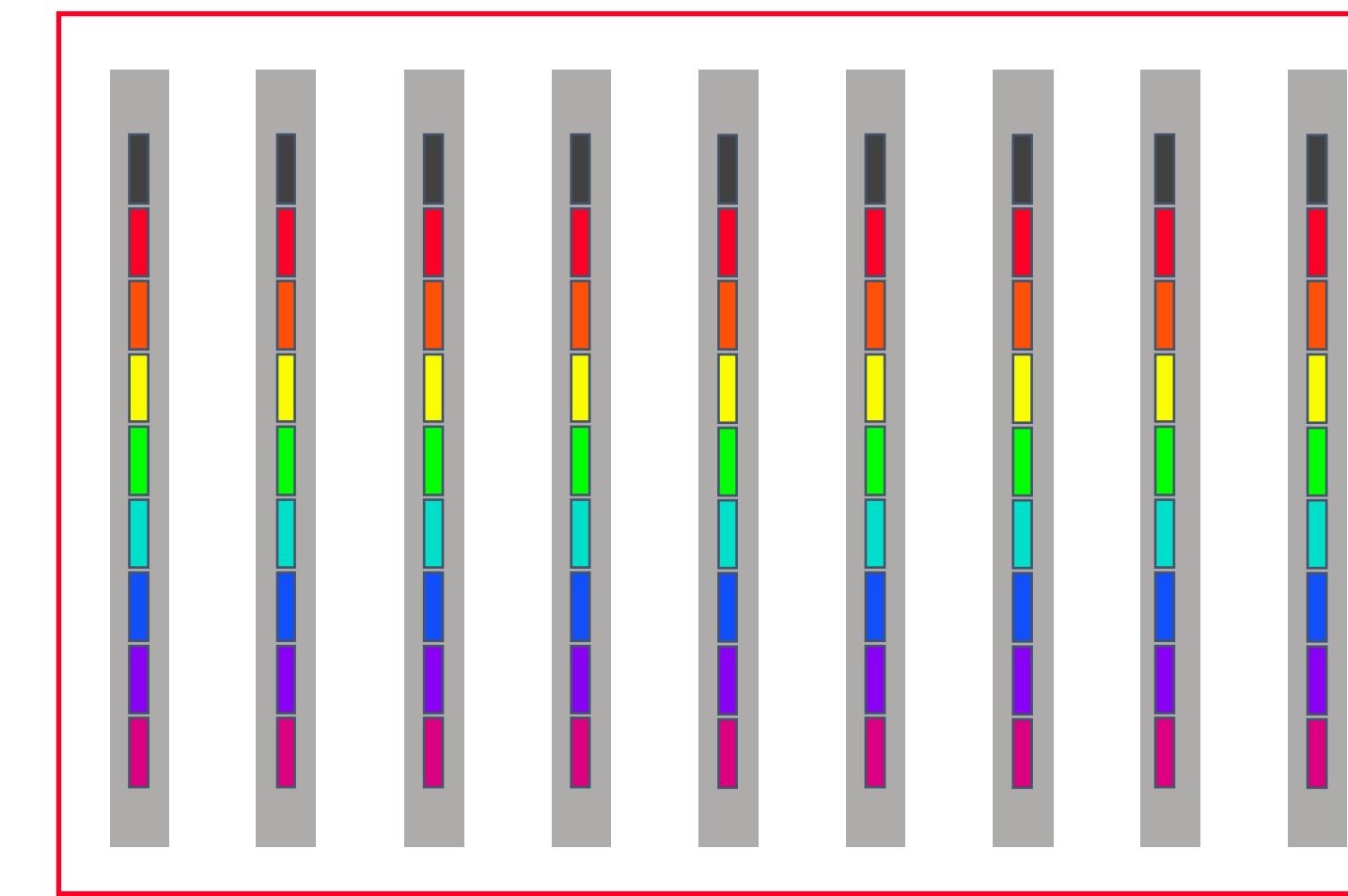
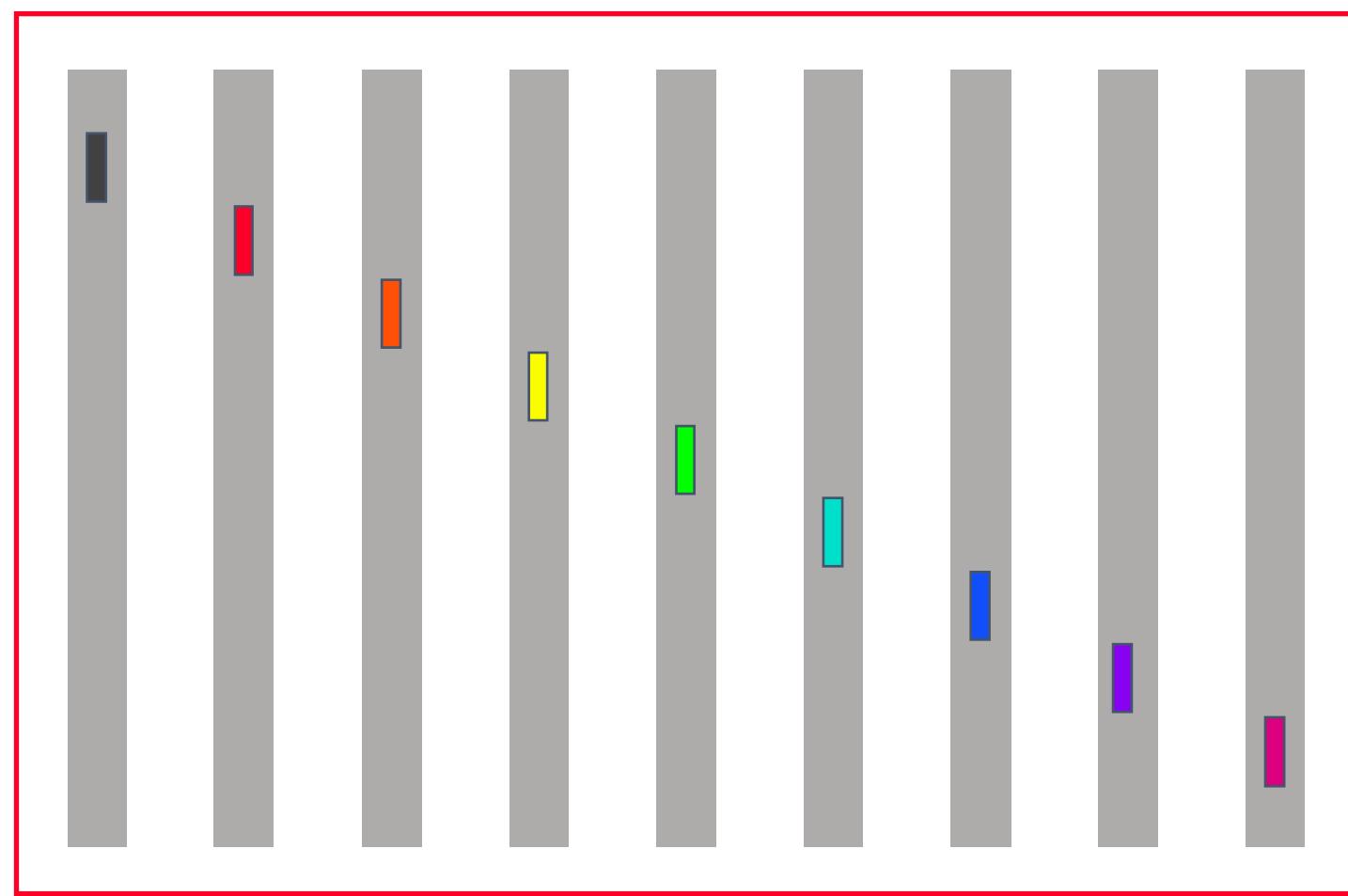
$$\begin{aligned} & \log(p) \left(\sum_{k=1}^{\log(p)} \left(\alpha + \frac{n}{2^k} \beta \right) \right) \\ & = \\ & \log(p) \quad \alpha + \frac{p-1}{p} n \beta \end{aligned}$$

Using the building blocks

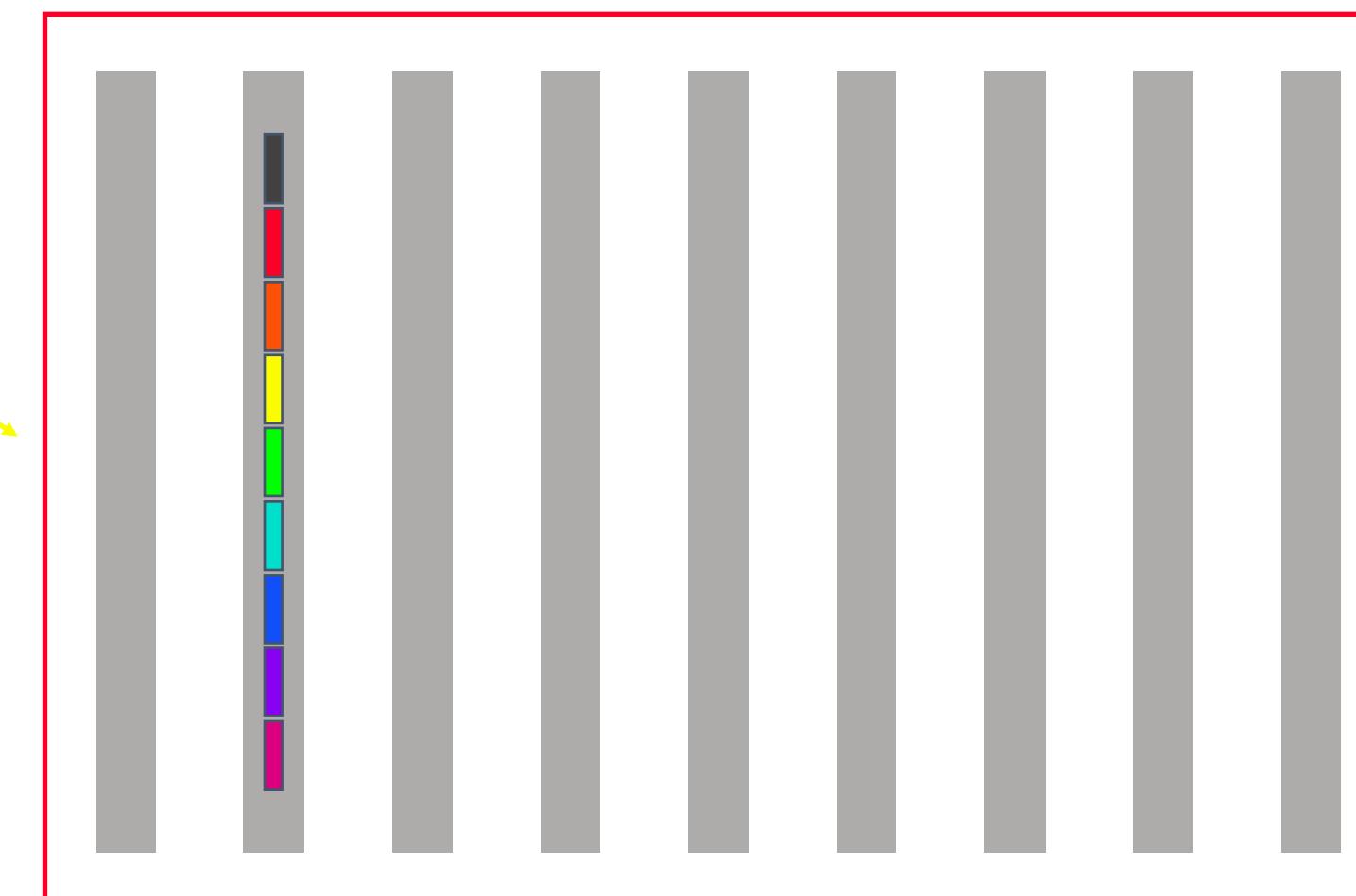
Allgather



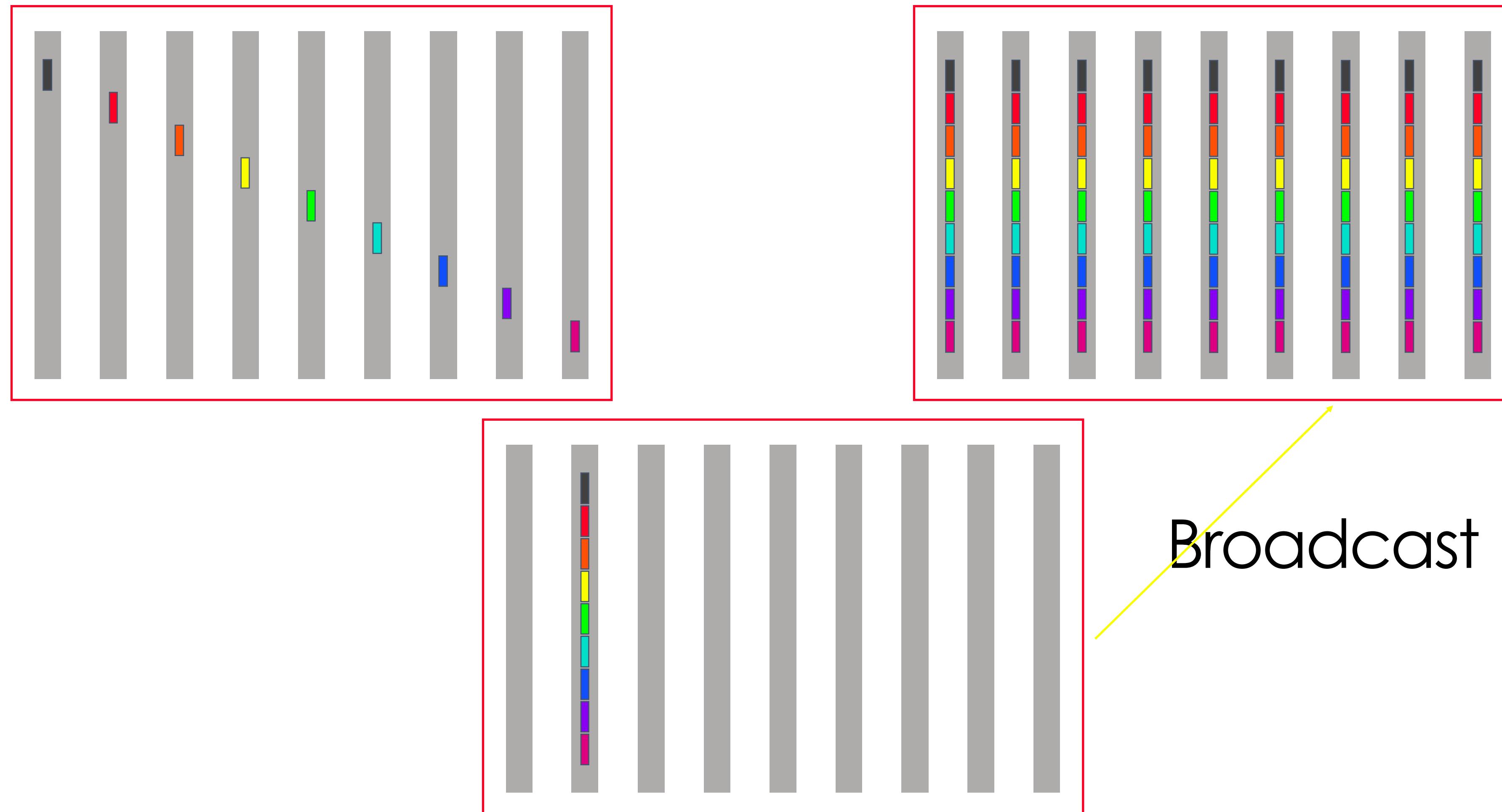
Allgather



Gather



Allgather (short vector)

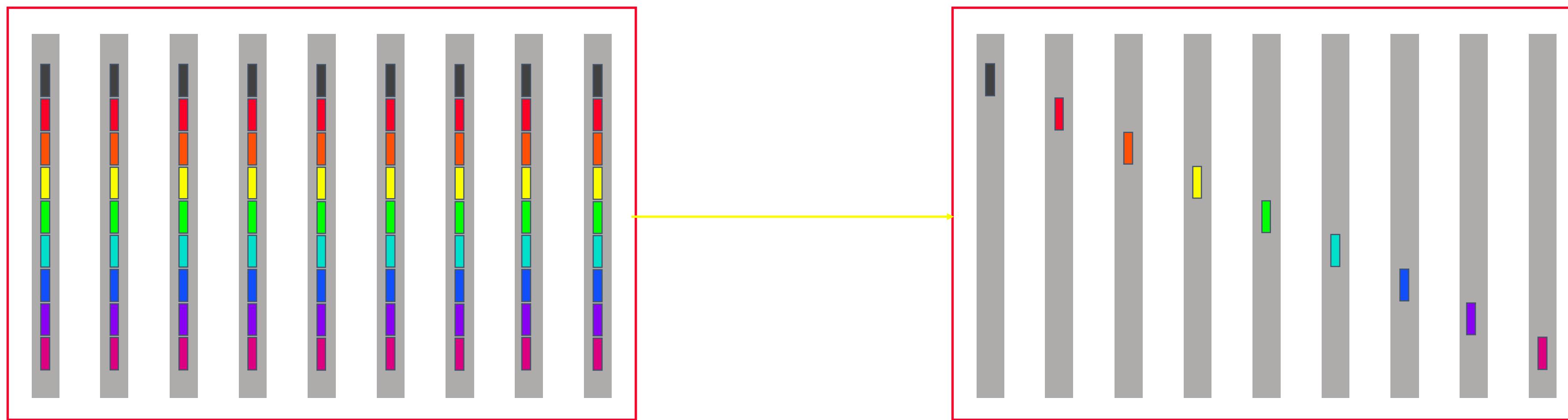


Cost of gather/broadcast allgather

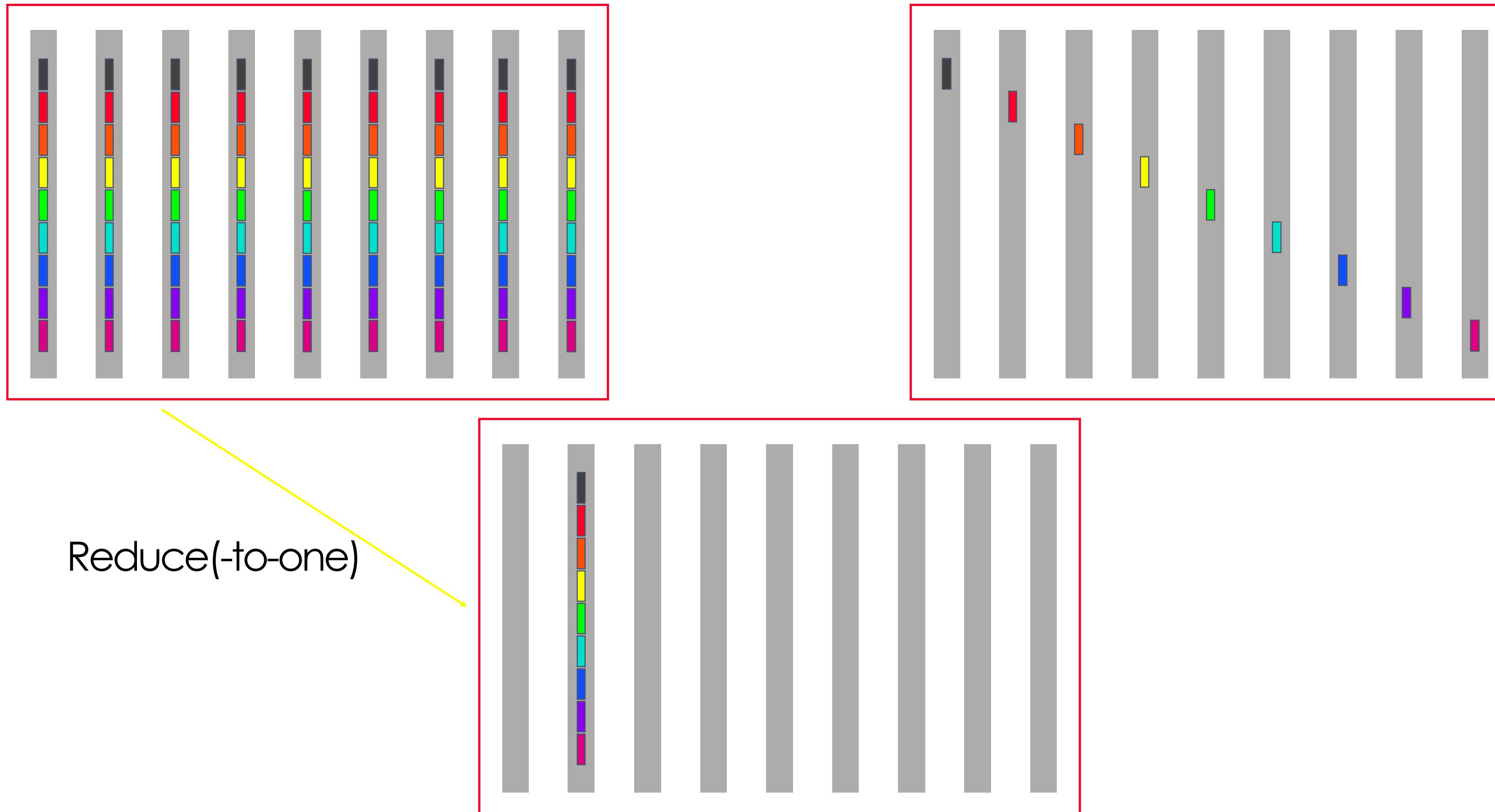
- Assumption: power of two number of nodes

gather	$\log(p)\alpha + \frac{p-1}{p}n\beta$
broadcast	$\frac{\log(p)(\alpha + n\beta)}{2\log(p)\alpha + \left(\frac{p-1}{p} + \log(p)\right)n\beta}$

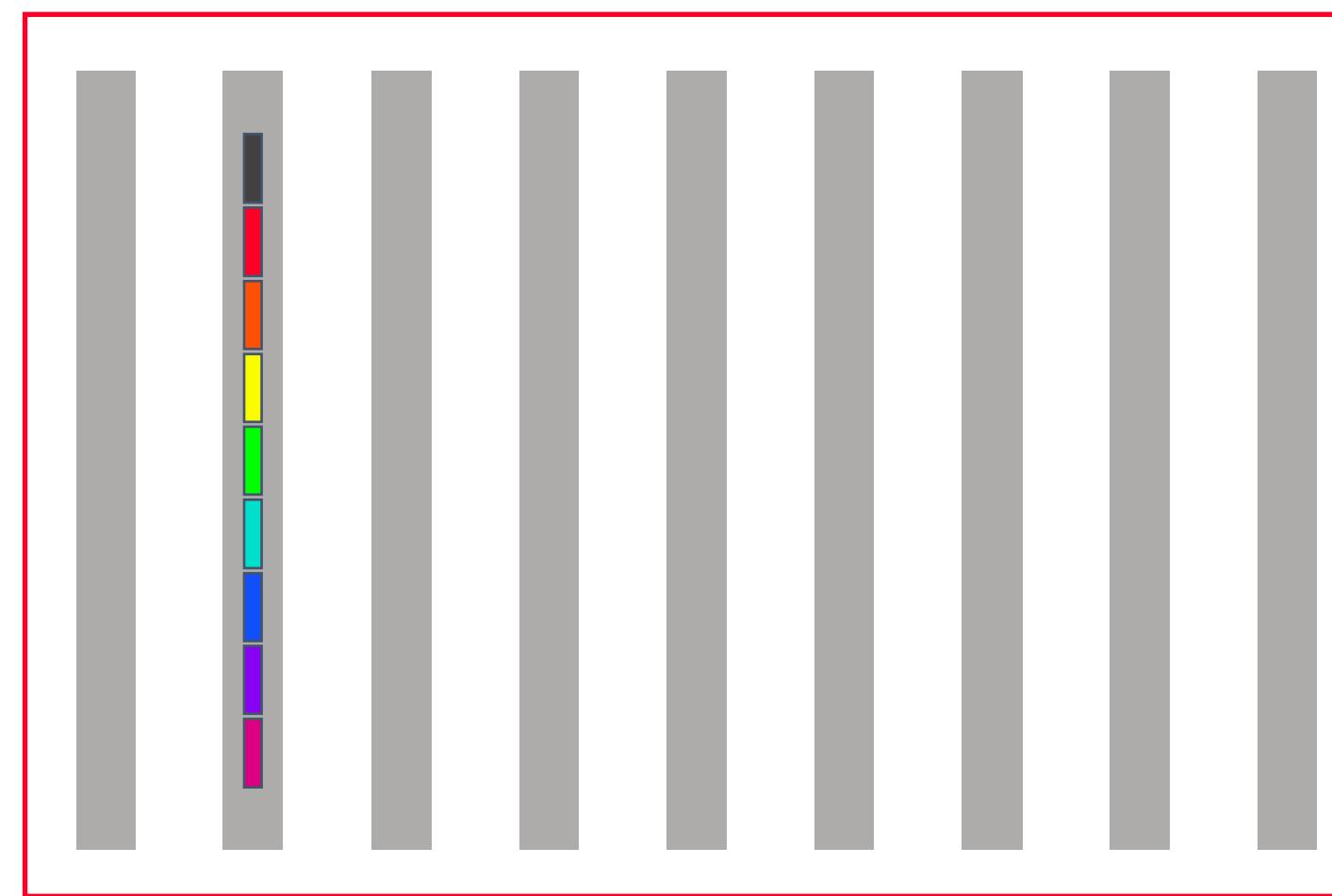
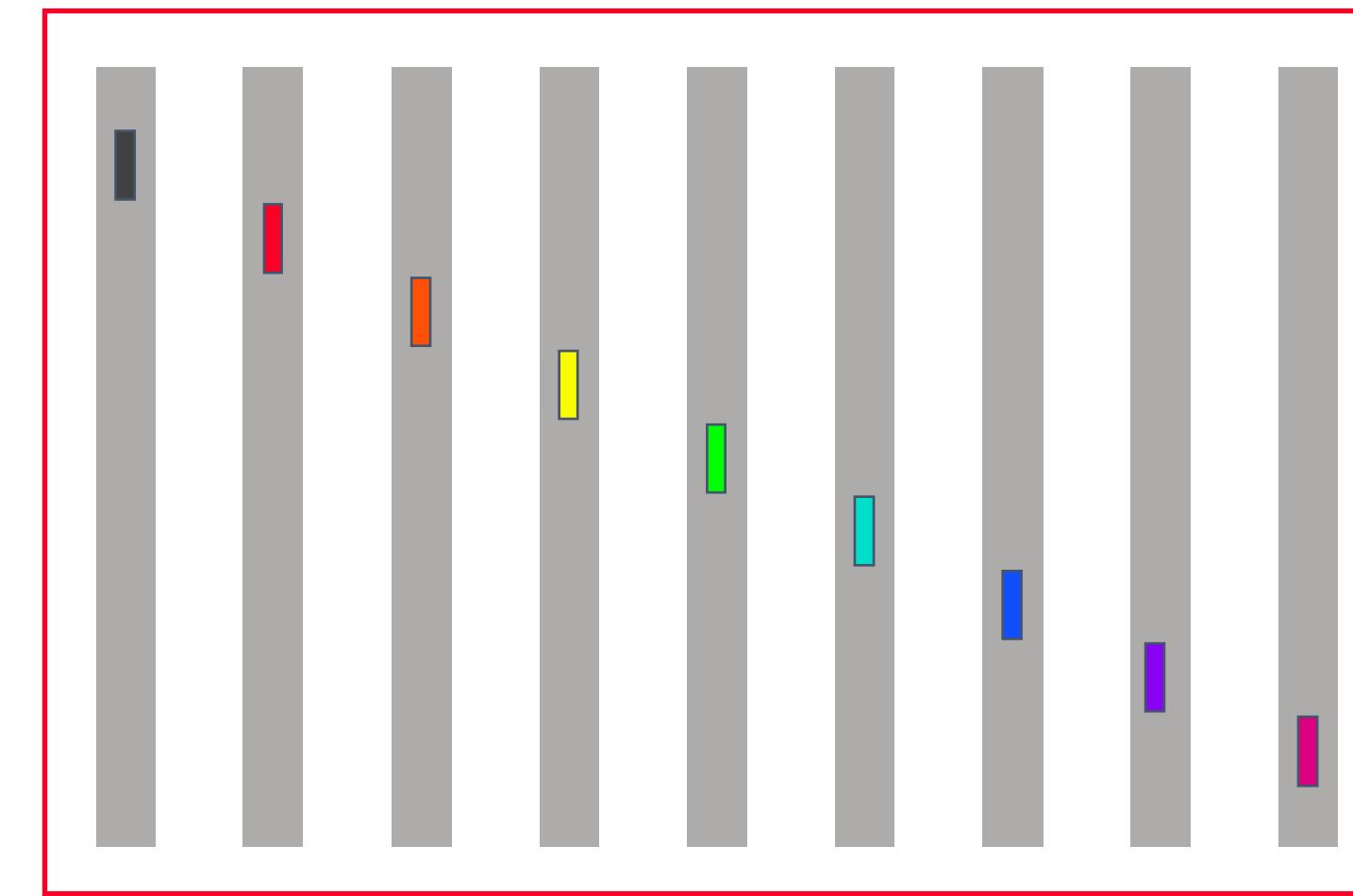
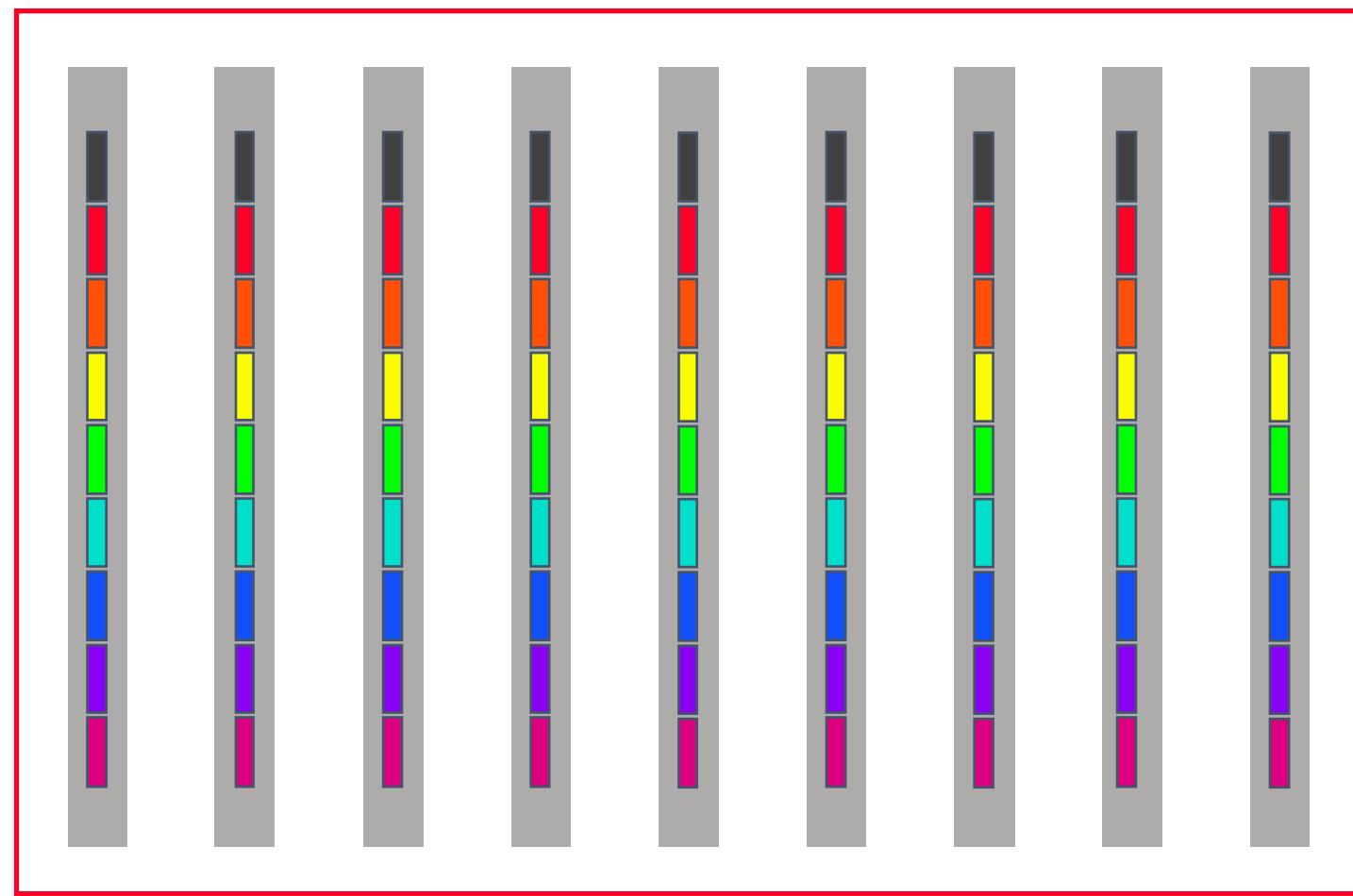
Reduce-scatter (small message)



Reduce-scatter (short vector)



Reduce-scatter (short vector)



Scatter

Cost of Reduce(-to-one)/scatter Reduce-scatter

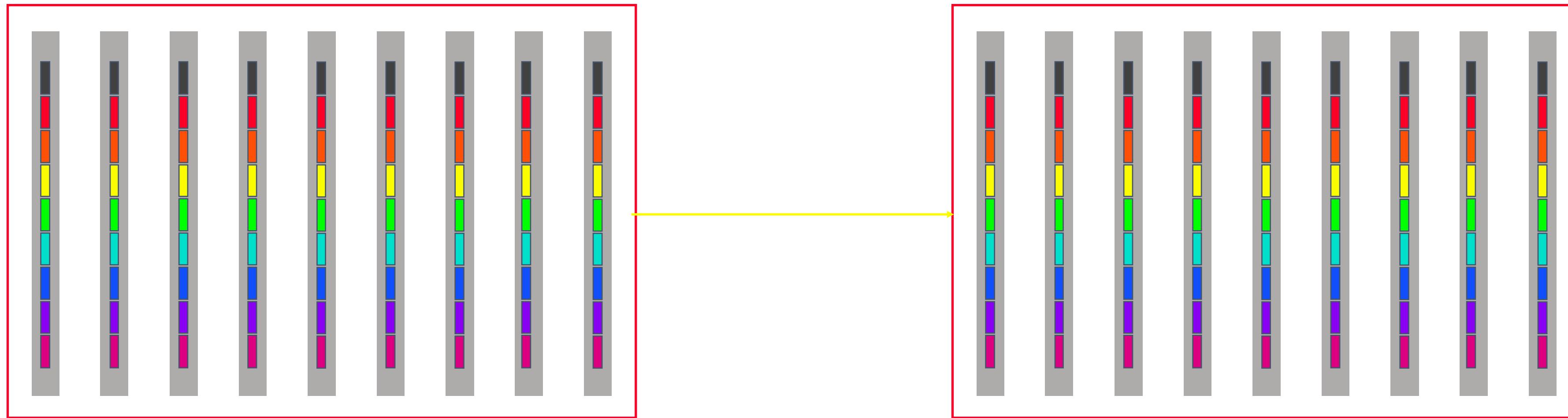
- Assumption: power of two number of nodes

$$\text{Reduce(-to-one)} \quad \log(p)(\alpha + n\beta + n\gamma)$$

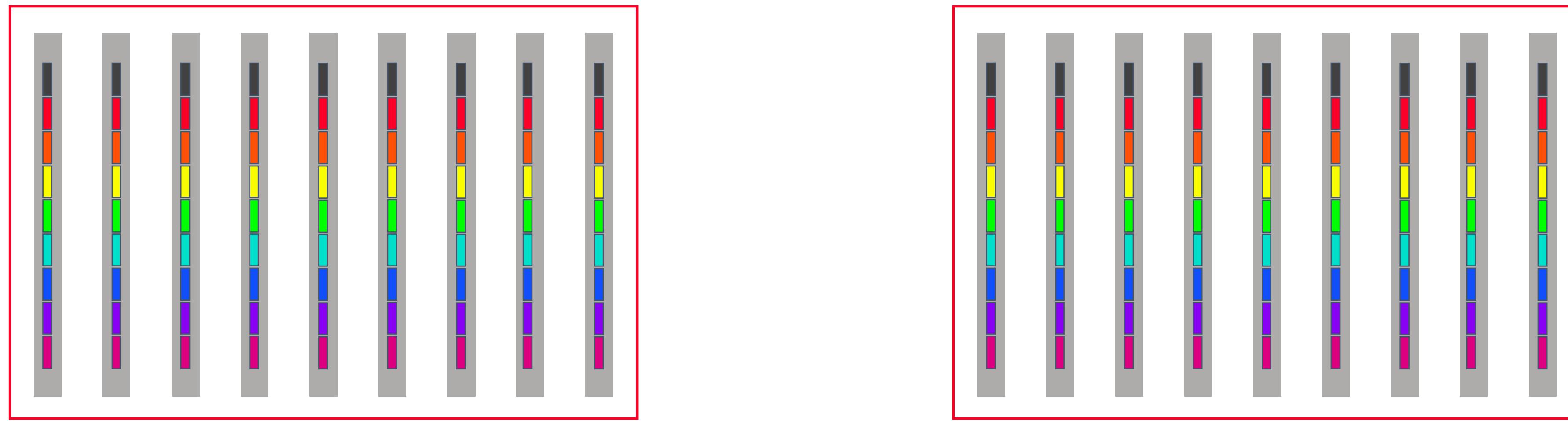
$$\text{scatter} \quad \log(p)\alpha + \frac{p-1}{p}n\beta$$

$$\frac{2\log(p)\alpha + \left(\frac{p-1}{p} + \log(p)\right)n\beta + \log(p)n\gamma}{\log(p)(\alpha + n\beta + n\gamma)}$$

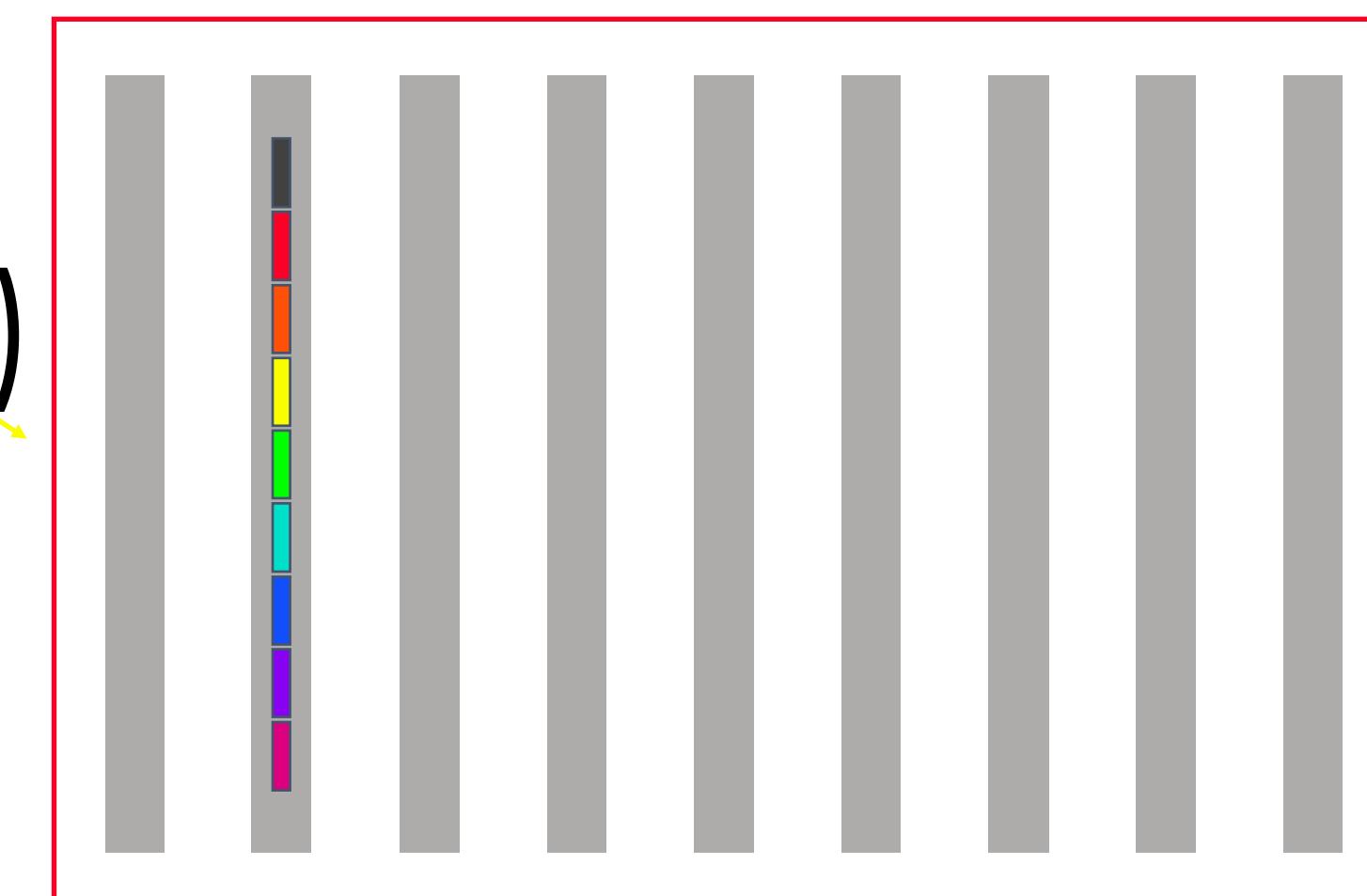
Allreduce (Latency-optimized)



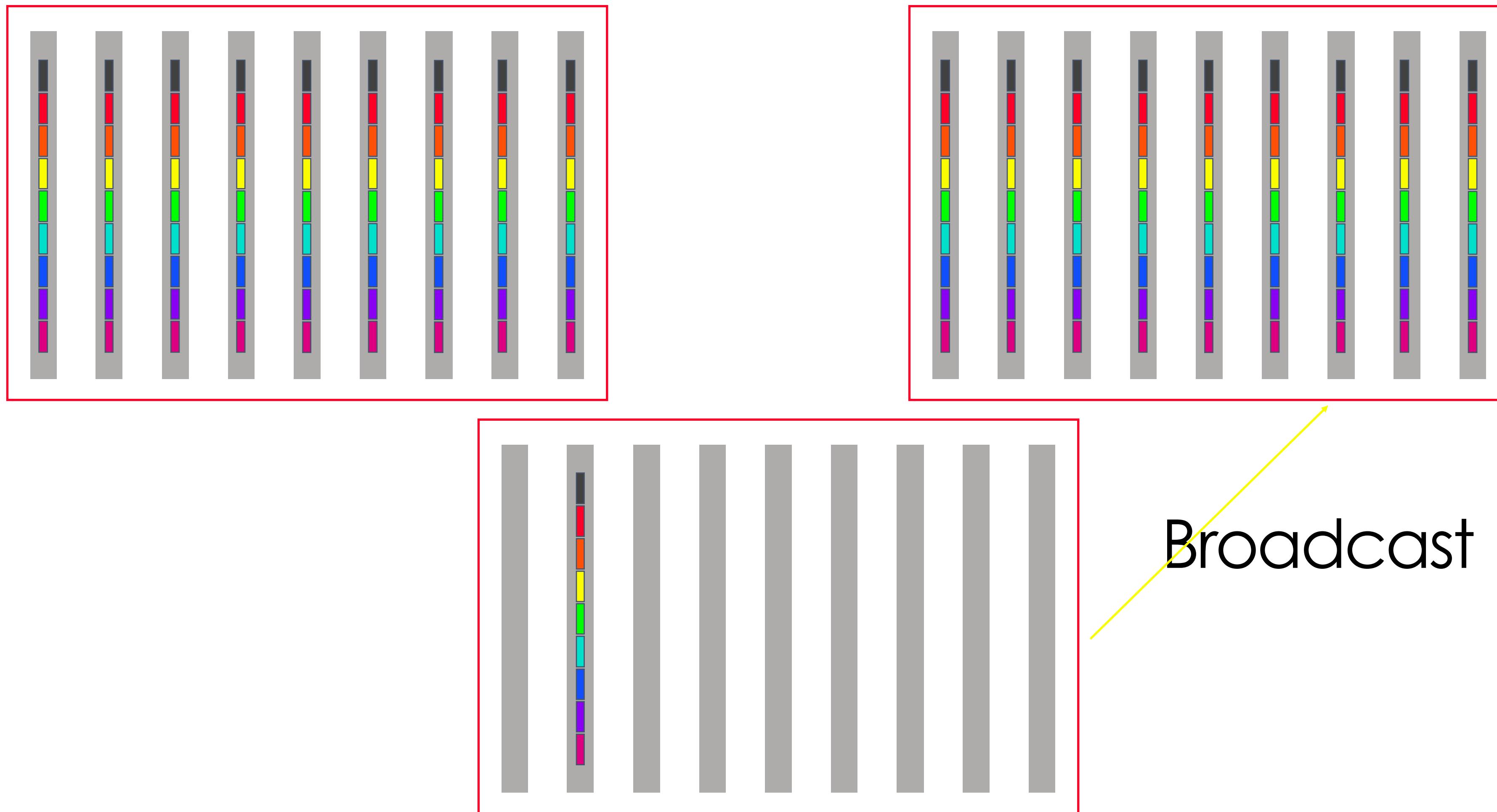
Allreduce (Latency-optimized)



Reduce(-to-one)



Allreduce (short vector)



Cost of reduce(-to-one)/broadcast Allreduce

- Assumption: power of two number of nodes

Reduce(-to-one) $\log(p)(\alpha + n\beta + n\gamma)$

broadcast $\log(p)(\alpha + n\beta)$

$$2\log(p)\alpha + 2\log(p)n\beta + \log(p)n\gamma$$

Recap

Reduce(-to-one)

$$\log(p)(\alpha + n\beta + n\gamma)$$

Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Broadcast

$$\log(p)(\alpha + n\beta)$$

Reduce-scatter

Allreduce

Allgather

Recap

Reduce(-to-one)

$$\log(p)(\alpha + n\beta + n\gamma)$$

Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Broadcast

$$\log(p)(\alpha + n\beta)$$

Reduce-scatter

$$2\log(p)\alpha + \log(p)n(\beta + \gamma) + \frac{p-1}{p}n\beta$$

Allreduce

Allgather

Recap

Reduce(-to-one)

$$\log(p)(\alpha + n\beta + n\gamma)$$

Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Broadcast

$$\log(p)(\alpha + n\beta)$$

Reduce-scatter

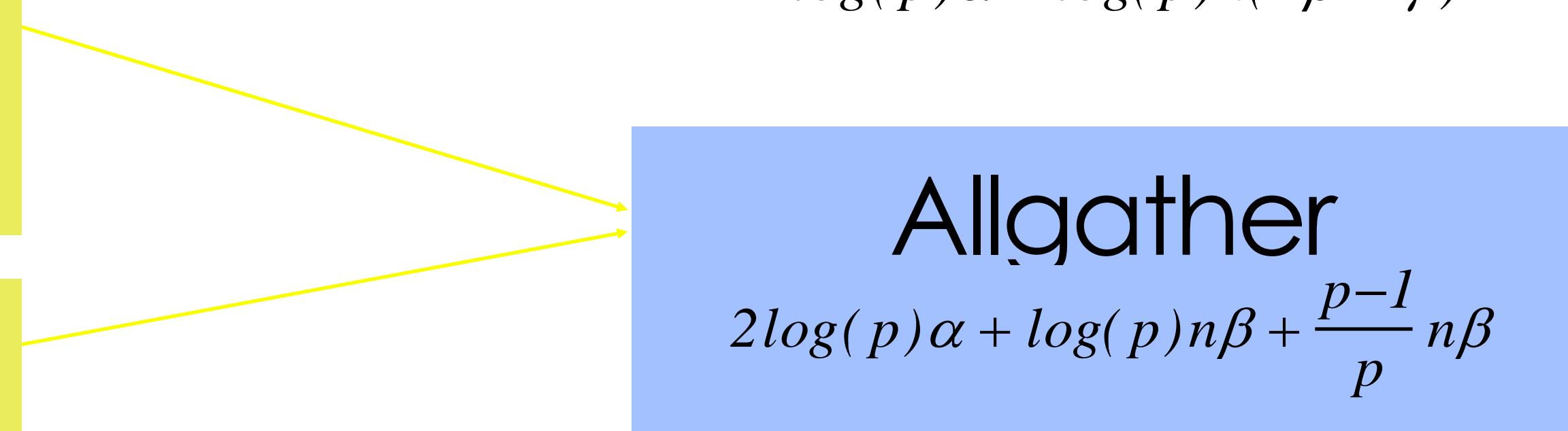
$$2\log(p)\alpha + \log(p)n(\beta + \gamma) + \frac{p-1}{p}n\beta$$

Allreduce

$$2\log(p)\alpha + \log(p)n(2\beta + \gamma)$$

Allgather

$$2\log(p)\alpha + \log(p)n\beta + \frac{p-1}{p}n\beta$$



Recap

Reduce(-to-one)

$$\log(p)(\alpha + n\beta + n\gamma)$$

Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Broadcast

$$\log(p)(\alpha + n\beta)$$

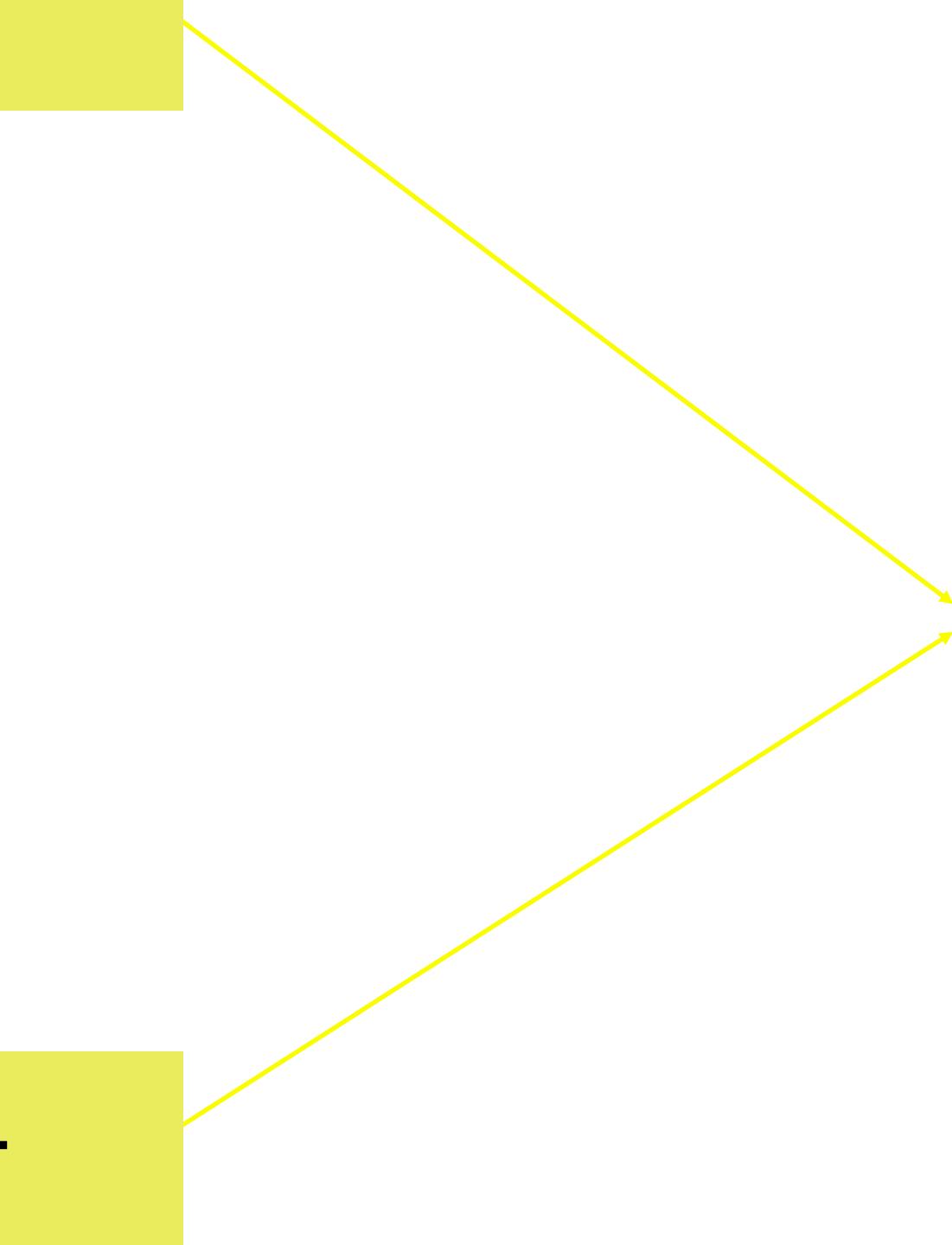
Reduce-scatter

$$2\log(p)\alpha + \log(p)n(\beta + \gamma) + \frac{p-1}{p}n\beta$$

Allreduce

$$2\log(p)\alpha + \log(p)n(2\beta + \gamma)$$

Allgather



Recap

Reduce(-to-one)

$$\log(p)(\alpha + n\beta + n\gamma)$$

Scatter

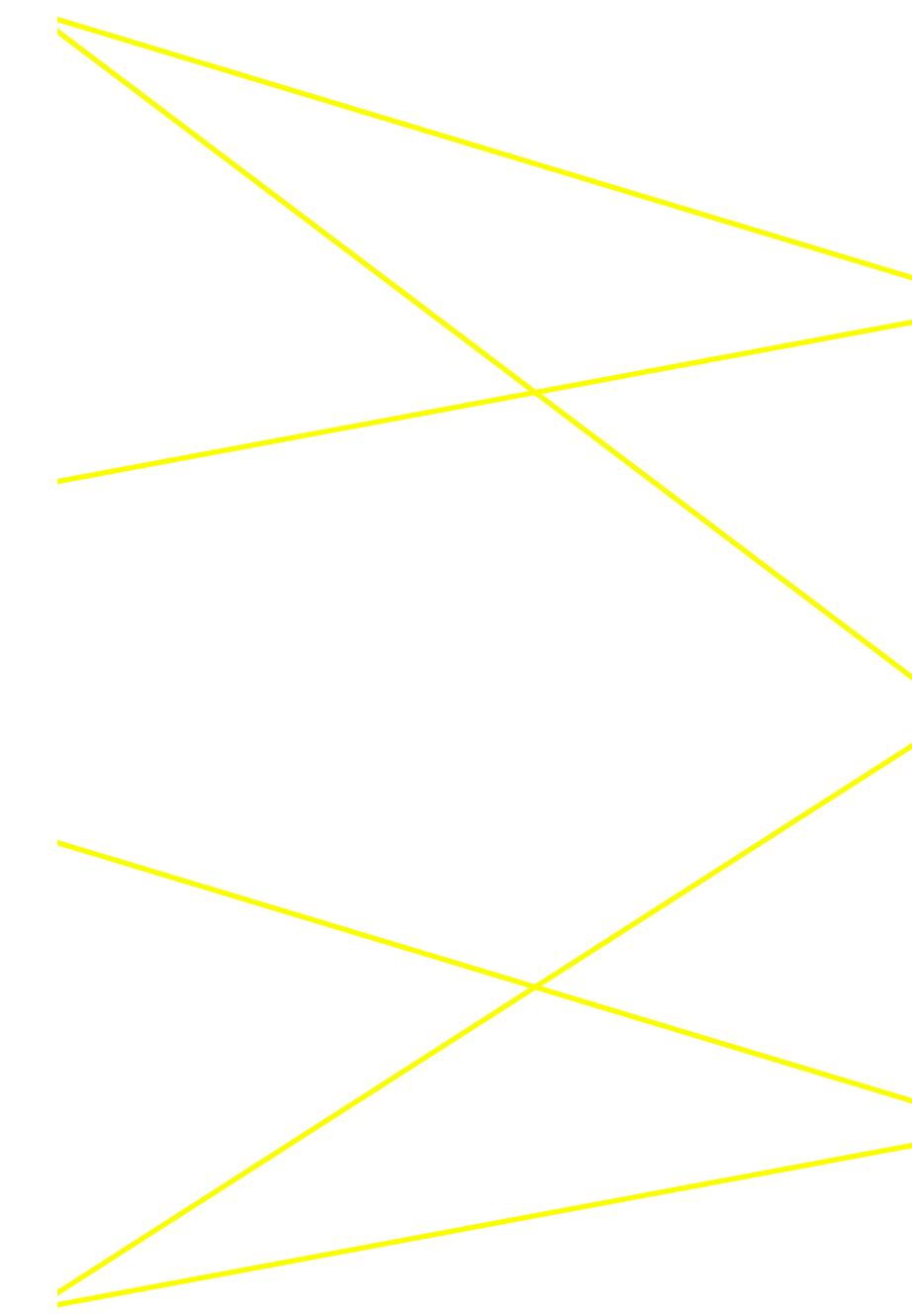
$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Broadcast

$$\log(p)(\alpha + n\beta)$$



Reduce-scatter

$$2\log(p)\alpha + \log(p)n(\beta + \gamma) + \frac{p-1}{p}n\beta$$

Allreduce

$$2\log(p)\alpha + \log(p)n(2\beta + \gamma)$$

Allgather

$$2\log(p)\alpha + \log(p)n\beta + \frac{p-1}{p}n\beta$$

Summary of MST algorithms

- Small message: Minimum Spanning Tree algorithm
 - Emphasize **low latency**
- **Can we do better**
- Problem of Minimum Spanning Tree Algorithm?
 - It prioritize latency rather than bandwidth
 - Hence: Some links are idle
- Next class: Large message size algorithm